

AUTOMATIC COLLECTION OF GROUP COMMUTATORS AND EXPRESSIONS

Walter J. Lamberth

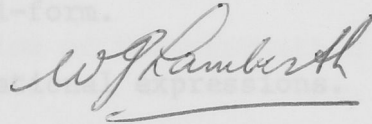
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PREFACE

Section 1 INTRODUCTION

The work reported in this thesis is my own. Where it is the development or adaptation of the work of others acknowledgment is given.



I thank Dr M.F. Newman for his supervision. I appreciated his guidance in the difficult selection of a suitable research topic in the hybrid field of group theory and symbol manipulation and his ready availability for critical discussion.

I also thank Dr M.A. Ward and Mr J.M. Campbell for several stimulating discussions.

Computations were performed on the A.N.U. S/360 computer. The thesis was patiently typed by Mrs F. Wickland.

Section 8	COMPUTATIONAL TREE	58
	8.1 Use of SLIP and FL1.	Page
Section 1	INTRODUCTION	1
Section 2	FORMAL GRAMMAR OF GROUP EXPRESSIONS	8
	2.1 Introduction.	8
	2.2 The Backus-normal-form.	8
	2.3 The grammar of rational expressions.	9
	2.4 The grammar of group expressions.	10
	2.5 Parsing of expressions and ambiguity.	12
Section 3	THE PARENTHESISING PROCESS	20
	3.1 Introduction.	20
	3.2 The partitioning functions.	21
	3.3 The parenthesising function.	29
Section 4	CANONIC EXPRESSIONS	33
	4.1 Introduction.	33
	4.2 Preliminary functions.	35
	4.3 Canonic rational expressions.	37
	4.4 Canonic group expressions.	40
Section 5	ROLE OF THE MODULUS	44
Section 6	TRANSFORMATION RULES	48
	6.1 Introduction.	48
	6.2 Group expression transformation.	50
	6.3 Rational expression transformation.	54
Section 7	USE OF THE TRANSFORMATION RULES	57
Section 13	PROCEDURES OF THE COLLECTING PROCESS	

Section 8	COMPUTATIONAL TECHNIQUES	58
	8.1 Use of SLIP and PL1.	58
	8.2 Representation of mathematical structures.	59
	8.3 Multiple structure reference.	62
	8.4 Structure transformation.	65
	8.5 Transformation representation.	65
	8.6 Pattern recognition or transformation selection.	67
	8.7 Structure scanning.	68
	8.8 Ordering of products.	69
	8.9 Input.	71
	8.10 Output.	72
Section 9	EFFECTIVENESS OF THE IMPLEMENTATION	74
Section 10	TERMS SYMBOLS AND NOTATIONS USED IN THE TEXT	77
	10.1 Commonly used symbols and notations.	77
	10.2 Symbols and notations introduced in the text.	78
	10.3 Functions and terms introduced in the text.	81
	10.4 Context-free grammars and languages.	86
	10.5 Ambiguity of formal grammars.	87
	10.6 Operator grammars.	88
	10.7 Precedence grammars.	88
	10.8 Conditional expressions.	89
Section 11	PROCEDURES OF THE SYMMETRIC LIST PROCESSOR - SYMPLIP	
Section 12	PRIMITIVES OF THE LIST PROCESSOR	
Section 13	PROCEDURES OF THE COLLECTING PROCESS	

Section 14 SAMPLE CALCULATION

Section 15 REFERENCES

The work described here arose from an interest in automatic symbol manipulation in pure mathematics. The commutator calculus of group theory was chosen as the specific area of study because of its apparent suitability and its interest to local mathematicians.

One of the main problems is that of converting expressions for group elements into a canonic form. This is normally called a collecting process.

The first stage of the work consisted in producing an intuitively based set of computer procedures for executing the collecting process. This was followed by an investigation, using a recursive formalism, of the theoretical foundations for the procedures. The relationship between the formal description and the computer procedures is close. However the former sometimes avoids, in the interests of clarity, details of the computer procedures which are aimed at increased efficiency.

While gaining its inspiration and motivation from group theory this work is not on group theory but rather about the meta-problems which are accentuated by an attempt to automate the collecting process.

The theory of the commutator calculus has been given by Magnus, Warram and Solitar [MAGNW660][†] and Ward [WARDM650] but detailed knowledge

[†] The code numbers appearing in brackets refer to entries in the list of references in Section 15. The first four letters of the code are the first four letters of the leading author's last name, the fifth letter of the code is the first letter of his first name. The first two digits indicate the year of publication. The final digit is used to distinguish multiple publications in the same year. This convention is that used in the Annotated Descriptor Based Bibliography on the Use of Computers for Non-numerical Mathematics [SMMJ661].

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of this is not necessary to appreciate the resulting problem in symbol manipulation.

The relevant theorem of commutator theory is the Basis theorem of P. Hall [HALLP340]. This states:

In a free group $F(r)$ on r generators there exists an infinite sequence of commutators C_v ($v = 1, 2, 3, \dots$) of non-decreasing weights in the generators such that for $n = 1, 2, 3, \dots$ and for every element W of $F(r)$,

$$W = C_1^{e_1} C_2^{e_2} \dots C_{k(n)}^{e_{k(n)}} V_{n+1}$$

where $V_{n+1} \in F_{n+1}$ (the $(n+1)^{\text{th}}$ group of the lower central series of $F(r)$) and where the integers $e_1, \dots, e_{k(n)}$ and V_{n+1} are uniquely determined by W . The number $k(n)$ is the number of C_v of weight $\leq n$.

The collecting process is the algorithm for transforming the element W to its unique form.

While W is referred to as an element of a group it is more useful for purposes of manipulation to regard it as an expression which is interpretable to a group element. Then the expressions are elements of a free universal algebra in the sense of Cohn [COHNP650]. This is essentially the approach taken here but it has been found desirable in part to take the abstraction one level further and regard the expressions as strings (or equivalently words, sentences or phrases) of a formal language with the single operation of concatenation and with a set of rules which select or generate the valid strings from the set of all strings on the alphabet.

The alphabet denoted by S' of this formal language is the set $\{a, \dots, z, 0, \dots, 9, +, -, *, |, /, :, (,), \vee\}$ and the symbol $' , '$. With this set as generators and the operation of concatenation the semigroup S'^* is formed. The language is a certain subset of S'^* . It is defined using another subset J' of S'^* - the set of rational expressions.

The subset J' is defined formally in Section 2. Briefly and without mention of the use of parentheses they may be described as follows:

- (a) ℓ, \dots, z and the rational numbers are rational expressions,
- (b) if λ_1 and λ_2 are rational expressions then $\lambda_1 + \lambda_2$, $\lambda_1 - \lambda_2$, $\lambda_1 * \lambda_2$ and $\lambda_1 | \lambda_2$ are rational expressions.

Of these operations the first three are the common rational operations of addition, subtraction and multiplication. The symbol $'|'$ is the combination operator and $\lambda_1 | \lambda_2$ is the binomial coefficient and is equivalent to the more common

$$\lambda_1 C_{\lambda_2} \quad \text{or} \quad \begin{pmatrix} \lambda_1 \\ \lambda_2 \end{pmatrix}$$

which are unsuited to the restricted expression style of present day line printers. While unary negation is included in the computer procedures it is excluded from the formal description of the grammar.

The subset of rational expressions is used to define the (disjoint) subset of S'^* referred to above. This is the formal language of group expressions and is denoted by G' . It is defined formally in Section 2 and briefly as follows:

- (a) a, \dots, k, \vee are group expressions,
- (b) if γ_1 and γ_2 are group expressions and λ_1 is a rational

expression then $\gamma_1\gamma_2$, γ_1,γ_2 and $\gamma_1:\lambda_1$ are group expressions.

Here juxtaposition of two group expressions denotes the group product operation on the two expressions.

The operation ':' is the exponentiation operator and the expression $\gamma_1:\lambda_1$ is equivalent to $\gamma_1\gamma_1\cdots\gamma_1$ (λ_1 times) and $\gamma_1:-1$ denotes the inverse of the expression γ_1 . The symbol v is the group identity. Commutation is denoted by the operator ',' and γ_1,γ_2 is called a commutator and is defined to be $\gamma_1^{-1}\gamma_2^{-1}\gamma_1\gamma_2$. This notation departs from the conventional practice which, regardless of precedence requirements, surrounds commutators with parentheses or brackets. In this text parentheses will be used only when it is necessary to override the defined precedences.

For more complex commutators the left norming convention is used so that $(\gamma_1,\gamma_2),\gamma_3$ may be written $\gamma_1,\gamma_2,\gamma_3$.

A common convention for avoiding use of parentheses in $\gamma_1,(\gamma_2,\gamma_3)$ is to use the semicolon operator. Then $\gamma_1,(\gamma_2,\gamma_3)$ becomes $\gamma_1;\gamma_2,\gamma_3$. In this text however the semicolon is replaced by two commas. This has the advantage that it is easily generalised to more complex cases and that the additional separation of the operands assists the eye in assessing the correct structure of the expression.

For theoretical and manipulative convenience S' is augmented by the symbols '.', '∩' and ',₁', ',₂', ..., ',_α' where α is an arbitrarily large but finite integer. The symbol '.' becomes an

explicit group product operation, the symbol ' \wedge ' an explicit concatenation operator between digits in the formation of integers and the symbol ' ${}_{,n}$ ', for example, is a replacement symbol for a string of n juxtaposed commas. The augmented alphabet is denoted by S and the semi-group generated by S is denoted by S^* .

Associated with the augmented semi-group S^* is a set J of modified rational expressions and a set G of modified group expressions. These sets have the property that they are unambiguous context-free phrase structured languages and as such the expressions may be generated and decomposed unambiguously. This is pursued in Section 2 where the language is defined by a grammar in the computationally useful formalism called the Backus-normal-form. The grammar is useful for determining if any string of symbols on the alphabet is valid or well formed and also for providing a basis for the unique parsing or parenthesising of a grammatical expression.

The conversion of group expressions to the canonic form proceeds in two stages. In the first stage the expression is transformed to a fully parenthesised form which is then converted to the canonic form. This is described in Section 3.

The process of fully parenthesising an expression is essentially the same as the process of analysing an expression. It is a useful process in that it provides an intermediate form which is structurally and grammatically simpler and which is the basis of many proofs and definitions in the commutator calculus.

The process is described in Section 3 in a formalism due to

McCarthy [MCCAJ630]. This formalism is called a recursive conditional expression and in its own area of application, namely that of computability it follows the trend towards more expressive computer languages. It was proposed by McCarthy as a more convenient, but still theoretically well founded, algorithmic formalism than that of either Turing machines [HERMH650, DAVIM580] or Markov Normal Algorithms [MARKA620]. The means for manipulating recursive conditional expressions are available as the programming languages LISP [MCCAJ650] and CPL [FOXLA660] and can be simulated to some extent in PL1.

The parenthesising process transforms the sets J and G to their fully parenthesised equivalents J^f and G^f .

A further pair of subsets of J^f and G^f are then defined in Section 4 where selection processes for determining membership of canonic subsets are given. The canonic subset of G^f is referred to by Ward as the set of basic expressions. In the case where only group products, inverses and commutators are admitted, Ward has proved that a group expression may be converted to the canonic or basic form by a finite application of the transformations. It seems reasonable to expect that a similar result would hold in the situation discussed here but this has not been proved.

Sections 5, 6 and 7 define the extended collecting process. The consequences of the extension are that the number of necessary transformations is approximately doubled and transformations for $(\gamma_1 \gamma_2) : \lambda$, $\gamma_1 : \lambda, \gamma_2$ and $\gamma_1, \gamma_2 : \lambda$ where λ is a rational expression are required. These present the difficulty that it does not appear possible to obtain the required transformation without a specific upper limit on weight. While it is believed that the inductive techniques used for

deriving these transformations could be specified algorithmically for any given upper limit on weight this has not been done. Rather, transformations which are valid for a modulus ≤ 7 have been calculated by hand and inserted into the procedures.

The PL1 procedures written to implement the collecting process bear a very close resemblance to the recursive conditional expressions and deviate significantly only in the interests of efficiency. The most significant difference is that in the computer implementation the role of the fully parenthesised expression is taken by an equivalent tree structure. This leads to rapid access to parts of an expression and very efficient modification of expressions. These procedures are discussed in Section 8 and are listed in Section 13.

It is not possible to make a generally valid statement about the effectiveness of the computer procedures. In fact the time required to collect the expression $(ab)^4$ is very much less than the time to collect the square of the collected value of $(ab)^2$. However the results leave much to be desired for while there are some computations which can be performed very much faster by machine than by hand there are other computations where the reverse is true.

To a great extent the reasons for this are understood and solutions are available and are discussed in Section 9. At the same time there is the ever present and imprecisely defined combinatorial expansion problem which will always make the heaviest demand on the ingenuity of the algorithm.

2. FORMAL GRAMMAR OF GROUP EXPRESSIONS

2.1 Introduction.

In this section the sets of rational expressions, group expressions and expressions will be defined as languages or subsets of words on an alphabet. The formal description of language is given in Section 10.4. The definitions of the various types of expressions will be given in terms of grammars which are also described formally in Section 10.4 but may be regarded informally as the information necessary either to generate all words or expressions of the language or to decide if some word on the alphabet is an element of the language.

It will be established that the group expressions form a context-free phrase structured language and that the grammar of this language is unambiguous.

2.2 The Backus-normal-form.

The formalism used to describe the alphabet and grammatical rules of the language is the Backus-normal-form otherwise known as the Backus-Naur-form or simply BNF. This notation was popularised by its use in the definition of the syntax of ALGOL-60 [NAURP600]. It gives special meaning to the following four meta-symbols:

< > ::= |

The first two are sometimes called diamond brackets and when they enclose a sequence of characters the sequence including the brackets is called a meta-variable.

The symbol '::=' separates the left and right parts of a rule. When parsing a sentence an occurrence of the left part may be replaced by the right part. The left and right parts of the rule are the components of the ordered pair described in Section 10.4.

The symbol '|' is a meta-connective and separates the right part of a rule into alternatives.

If any other symbol appears with these symbols but not enclosed in diamond brackets then that symbol stands for itself. Thus the rule

$$\langle \text{term} \rangle ::= \langle \text{factor} \rangle \mid \langle \text{term} \rangle * \langle \text{factor} \rangle$$

means that a $\langle \text{term} \rangle$ may be a $\langle \text{factor} \rangle$ or a $\langle \text{term} \rangle$ followed by the symbol '*' followed by a $\langle \text{factor} \rangle$.

2.3 The grammar of rational expressions.

The expressions to be formalised consist of two disjoint parts. These consist on the one hand of expressions made up of the rational numbers and certain variables representing these numbers. On the other hand there are expressions whose elements or operands are group expression generators. The former expressions may appear only on the right hand side of the exponentiation symbol ':' while the latter may appear only on the left hand side unless surrounded by parentheses.

Corresponding to the first mentioned part there is the grammar of rational expressions.

1. $\langle \text{rational expression generator} \rangle ::= \ell \mid m \mid n \mid o \mid p \mid q \mid r \mid s \mid t \mid$
 $u \mid v \mid w \mid x \mid y \mid z$

2. $\langle \text{digit} \rangle ::= 0|1|2|3|4|5|6|7|8|9$
3. $\langle \text{integer} \rangle ::= \langle \text{digit} \rangle \mid \langle \text{integer} \rangle \langle \text{digit} \rangle$
4. $\langle \text{rational number} \rangle ::= \langle \text{integer} \rangle \mid \langle \text{integer} \rangle / \langle \text{integer} \rangle$
5. $\langle \text{primary} \rangle ::= \langle \text{rational expression generator} \rangle \mid \langle \text{rational number} \rangle$
 $\mid (\langle \text{rational expression} \rangle)$
6. $\langle \text{factor} \rangle ::= \langle \text{primary} \rangle \mid \langle \text{factor} \rangle \langle \text{the symbol } '|\rangle \langle \text{integer} \rangle$
7. $\langle \text{term} \rangle ::= \langle \text{factor} \rangle \mid \langle \text{term} \rangle * \langle \text{factor} \rangle$
8. $\langle \text{rational expression} \rangle ::= \langle \text{term} \rangle \mid \langle \text{rational expression} \rangle + \langle \text{term} \rangle$
 $\mid \langle \text{rational expression} \rangle - \langle \text{term} \rangle$

This grammar defines the language or set of rational expressions which is denoted by the symbol J' .

It may be noted that because of the dual use of the symbol $'|'$ as an operator and also as a meta-symbol it has been described in its role as an operator informally inside the meta-brackets. It will also be noted that the number and period which precedes each rule is not part of the rule but merely serves to identify it.

2.4 The grammar of group expressions.

The grammar of group expressions is now given and it can be seen that in rule 11 the grammar of rational expressions is incorporated to the right of the symbol $':'$.

9. $\langle \text{group expression generator} \rangle ::= a|b|c|d|e|f|g|h|i|j|k|v$
10. $\langle \text{g-primary} \rangle ::= \langle \text{group expression generator} \rangle \mid (\langle \text{group expression} \rangle)$
11. $\langle \text{g-factor} \rangle ::= \langle \text{g-primary} \rangle \mid \langle \text{g-factor} \rangle : \langle \text{rational expression} \rangle$
12. $\langle \text{g-term} \rangle ::= \langle \text{g-factor} \rangle \mid \langle \text{g-term} \rangle \langle \text{g-factor} \rangle$

13. $\langle c_0 \rangle ::= \langle g\text{-term} \rangle$
14. $\langle c_\alpha \rangle ::= \bigcup_{i=0}^{i=\alpha-1} \{ \{ \langle c_\alpha \rangle, {}_\alpha \langle c_i \rangle \} \cup \{ \langle c_i \rangle, {}_\alpha \langle c_{\alpha-1} \rangle \} \}$ (see below)
15. $\langle \text{group expression} \rangle ::= \langle g\text{-term} \rangle \mid \langle c_\alpha \rangle$ for $\alpha > 0$.

This grammar denoted by G' defines the language or set of group expressions which is denoted by the symbol G' . It closely resembles the grammar actually used by the computer implementation for communication with the mathematician. In particular it may be noted that the generators are single symbols and in the definition of $\langle g\text{-term} \rangle$ the product operator is implicit.

Rule 14 of the grammar which defines the structure of commutators is unusual in that it uses the non-standard meta-symbols

$$\bigcup_{i=0}^{i=\alpha-1} \{ \quad \} ,_\alpha$$

In group expressions as used by the mathematician the symbol $' ,_\alpha '$ is actually a string of α juxtaposed commas.

It should also be noted that in more conventional BNF notation the rule for $\langle c_\alpha \rangle$, the commutator of depth α , is really an infinite set of rules beginning:

$$\begin{aligned} \langle c_1 \rangle &::= \langle c_1 \rangle, \langle c_0 \rangle \mid \langle c_0 \rangle, \langle c_0 \rangle \\ \langle c_2 \rangle &::= \langle c_2 \rangle, , \langle c_0 \rangle \mid \langle c_0 \rangle, , \langle c_1 \rangle \mid \\ &\quad \langle c_2 \rangle, , \langle c_1 \rangle \mid \langle c_1 \rangle, , \langle c_1 \rangle \\ \langle c_3 \rangle &::= \langle c_3 \rangle, , , \langle c_0 \rangle \mid \langle c_0 \rangle, , , \langle c_2 \rangle \mid \\ &\quad \langle c_3 \rangle, , , \langle c_1 \rangle \mid \langle c_1 \rangle, , , \langle c_2 \rangle \mid \\ &\quad \langle c_3 \rangle, , , \langle c_2 \rangle \mid \langle c_2 \rangle, , , \langle c_2 \rangle \end{aligned}$$

For this reason the grammar is not strictly of the context-free type. In practice a finite limit can be placed on α which is adequate for the finite expressions encountered in any specific case.

The mathematical notation defined by the rule for $\langle c_\alpha \rangle$ is, as mentioned in the introduction, also slightly unconventional. It embodies the common left norming convention which states that when precedence of operations is not defined by either the precedence of operators or the use of parentheses then the leftmost operator has precedence. So, for example a,b,c,d is to be read $((a,b),c),d$. However the conventional ';' operator has been replaced by ',,' so that the expression $(a,b),(c,d)$ is normally written $a,b;c,d$ while here it is written $a,b,,c,d$. This has the advantage that it can be easily generalised to arbitrarily many commas in more complex cases and also that it assists the eye in its parsing of an expression.

It may be noted in the examples just given that while commutators are conventionally surrounded by either parentheses or brackets this is not necessary here unless the grammatically defined precedences have to be overridden. In fact the grammar specifies a precedence for all operators so that the visual obscurity associated with full parenthesisation may be avoided.

2.5 Parsing of expressions and ambiguity.

For the purposes of manipulation the expressions in $J' \cup G'$ are inconvenient because the formal processes of mathematical manipulation and the associated proofs are usually given only for a subset of $J' \cup G'$ -

the subset of fully parenthesised expressions. The process of fully parenthesising, which is the subject of the next section, is closely related to that of parsing of a sentence in a language or transforming the sentence into an explicitly structured form.

It is important to have a grammar for which there is a unique parsing process so that for any element of the language there is a unique fully parenthesised form. This unique parsing property is not achieved without care in formulation of the grammar. Suppose that the previously given definition of $\langle c_{\alpha} \rangle$ is replaced by the following BNF rules.

```

<commutator operator> ::= , | <commutator operator>,
<commutator> ::= <g-term> | <commutator>
                  <commutator operator><g-term> |
                  <commutator><commutator operator><commutator>

```

Using the obvious abbreviations these rules become

```

<co> ::= , | <co>,
<c> ::= <gt> | <c><co><gt> | <c><co><c>

```

This grammar will always generate valid commutators although they may possibly be supplied with a redundant number of commas. However as the example below shows it is not possible to parse or fully parenthesise a commutator unambiguously using this grammar. See Section 10.5 for a summary of the relevant theory. Ginsburg [GINSS660] gives a full discussion.

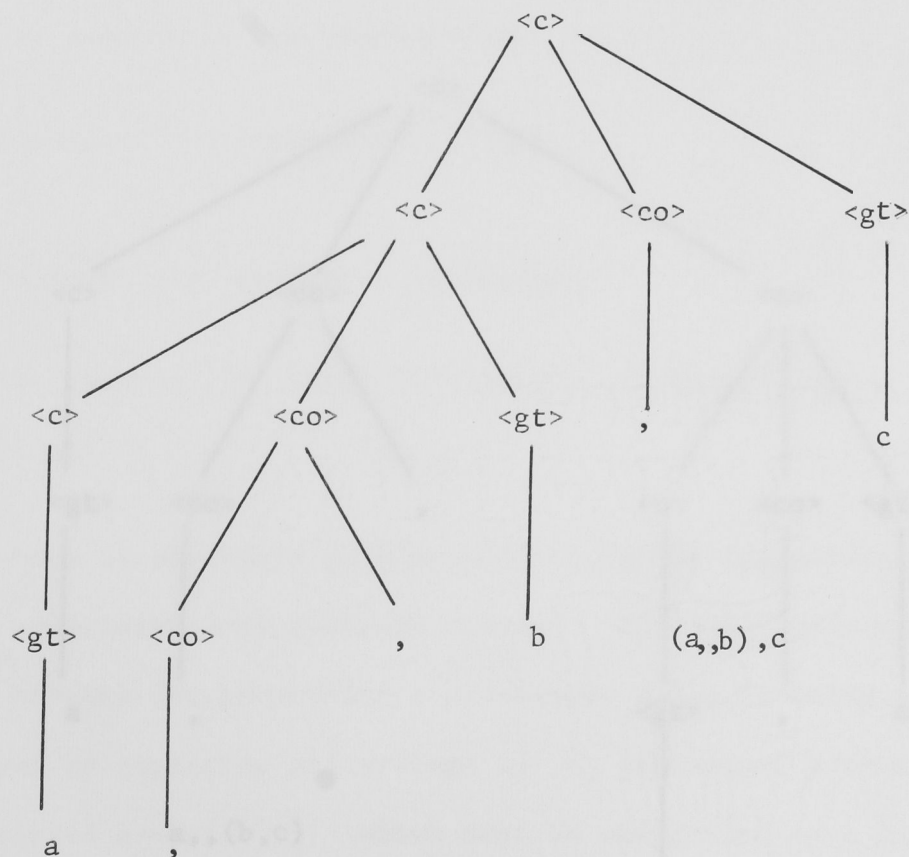
Applying this new grammar to the commutator a, b, c yields two left-most derivations. (The use of ' \rightarrow ' is in accord with Ginsburg's usage.)

The first derivation is as follows:

```

<c> → <c><co><gt>
      → <c><co><gt><co><gt>
      → <gt><co><gt><co><gt>
      → a<co><gt><co><gt>
      → a<co>,<gt><co><gt>
      → a,,<gt><co><gt>
      → a,,b<co><gt>
      → a,,b,<gt>
      → a,,b,c
  
```

The corresponding tree structure and fully parenthesised expression are given below



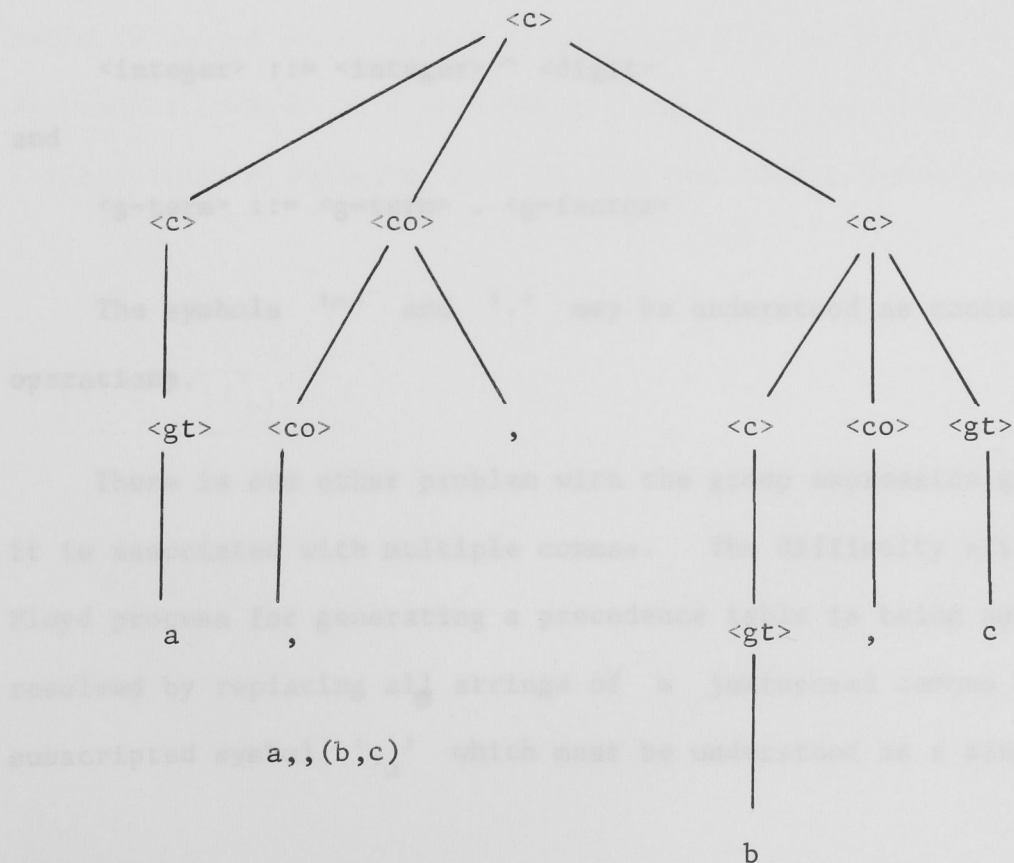
The second derivation is as follows:

```

<c> → <c><co><c>
      → <gt><co><c>
      → a<co><c>
      → a<co>,<c>
      → a,,<c>
      → a,,<c><co><gt>
      → a,,<gt><co><gt>
      → a,,b<co><gt>
      → a,,b,<gt>
      → a,,b,c

```

The corresponding tree structure and fully parenthesised expression are given below.



The ambiguity of the grammar just given is caused by the presence of juxtaposed meta-symbols in one of its rules. The existence of a unique parse structure for certain types of grammars has been proved by Floyd [FLOYR620, FLOYR621, FLOYR630] who has also established criteria for detecting ambiguity in grammars and has described algorithms for performing the parsing. A summary of these grammars is given in Section 10.

Inspection of the group expression grammar shows that there are two rules which do not conform to the restrictions of operator grammars because they contain juxtaposed meta-symbols. These are

3. $\langle \text{integer} \rangle ::= \langle \text{integer} \rangle \langle \text{digit} \rangle$

and

12. $\langle \text{g-term} \rangle ::= \langle \text{g-term} \rangle \langle \text{g-factor} \rangle$

The following simple modifications can be made to produce rules which do conform to the restrictions.

$\langle \text{integer} \rangle ::= \langle \text{integer} \rangle \wedge \langle \text{digit} \rangle$

and

$\langle \text{g-term} \rangle ::= \langle \text{g-term} \rangle . \langle \text{g-factor} \rangle$

The symbols ' \wedge ' and ' $.$ ' may be understood as concatenation operations.

There is one other problem with the group expression grammar and it is associated with multiple commas. The difficulty arises when the Floyd process for generating a precedence table is being applied and is resolved by replacing all strings of α juxtaposed commas by the subscripted symbol ' $,_{\alpha}$ ' which must be understood as a single symbol.

It is evident that a simple transformation exists for conversion from one form of group expression to the other.

Formally these processes of modification of the input/output grammar G' to the internally useful grammar which will be denoted by G consist of extending the set S' of generators of the semigroup S'^* with the symbols $'.'$, $'\cap'$ and $'_1'$, $'_2'$, ..., $'_n'$ where n is finite but arbitrarily large. The new set of generators is denoted by S and the semigroup generated by S is denoted S^* . Corresponding to this change in the grammar the set J' is transformed to a set denoted by J and the set G' is transformed to a set denoted by G .

As the set $J \cup G$ is frequently referred to in the text it will be denoted by the symbol E . The set $J' \cup G'$ is denoted E' .

With these difficulties removed it is possible to build a precedence table in accord with Floyd's algorithm and to verify that the group expression grammar is a precedence grammar and may be uniquely parsed. The resulting precedence table follows but proofs are omitted.

	(a	:	•	' ₁	' ₂	' ₃	ℓ	0	\sim	/		*	+	-)
)			•>	•>	•>	•>	•>					•>	•>	•>	•>	•>
a			•>	•>	•>	•>	•>									•>
:	<•		•>	•>	•>	•>	•>	<•	<•	<•	<•	<•	<•	<•	<•	•>
•	<•	<•	<•	•>	•>	•>	•>									•>
' ₁	<•	<•	<•	<•	•>	•>	•>									•>
' ₂	<•	<•	<•	<•	<•	•>	•>									•>
' ₃	<•	<•	<•	<•	<•	<•	•>									•>
ℓ												•>	•>	•>	•>	•>
0										•>	•>	•>	•>	•>	•>	•>
\sim									<•	•>	•>	•>	•>	•>	•>	•>
/									<•	<•		•>	•>	•>	•>	•>
									<•	<•	<•	•>	•>	•>	•>	•>
*	<•							<•	<•	<•	<•	<•	•>	•>	•>	•>
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Precedence table for group expression grammar.

It should be noted in connection with this table that some rows and columns are only representative of a class of symbols. In particular:

- a is a representative row or column for the letters a, \dots, k, v
- ℓ is a representative row or column for the letters ℓ, \dots, z
- 0 is a representative row or column for the digits $0, \dots, 9$
- '₁ '₂ and '₃ are the first three rows or columns of an arbitrarily large set.

3. THE PARENTHEISING PROCESS.

3.1 Introduction.

The group expressions defined in the previous section are in a form which is unsuited to efficient automatic manipulation. In this section functions are defined which enable the grammatically specified structure of an expression to be exposed.

The link between the grammar and the functions of this section is established by a function dlt ('dot less than'). With the use of this link three partitioning functions are defined and it is shown that the original expression may be reconstituted by appropriate concatenation of the values of the three functions.

Following this a parenthesising process is defined. It cannot be proved that this process does fully parenthesise except by reference to some other definition which must of necessity be equally precise. The given definition may, however, be checked for conformity with the intuitive notion. It is shown that the process can be performed in a finite number of steps and that it parses the expression in accord with the grammar.

The symbols μ and ϕ are used extensively in this section. The domain of μ is the set N of natural numbers and the domain of ϕ , possibly subscripted, is the set \hat{E} . \hat{E} is the set of all expressions in E together with all words derived from expressions by deletion of one or more symbols from either end of the expression together with the set $\{\Lambda\}$. More precisely \hat{E} is defined as follows:

$$\phi \in \hat{E} \iff \phi \in E \cup \{\Lambda\} \vee (\exists \psi: \psi \in \hat{E} \wedge (\phi = \text{lr}[\psi] \vee \phi = \text{rr}[\psi]))$$

where lr and rr are defined below.

The null word Λ has the role of an identity element in \hat{E} with the operation of concatenation and so

$$\phi\Lambda = \Lambda\phi \quad \text{and} \quad \Lambda\Lambda = \Lambda$$

It will always be assumed that expressions submitted for processing do not contain the null word. Either it is unnecessary in an expression or the expression itself is null in which case no useful computation can be performed on it. The usefulness of the null word is internal to the functions to be defined. The convention that the leftmost or rightmost symbol of the null word is the null word will be adopted.

3.2 The partitioning functions.

Four primitive partitioning functions are defined to take the values of the single symbols at either end of a word ϕ or to take the values of the remaining words after the end symbols have been deleted. A brief mnemonic is used as the name of each function but an explanatory name is associated with the definition and appears on the right hand side of the page. As their domain and range these functions have the set \hat{E} .

$ls[\phi] =$ - left symbol

the leftmost symbol of ϕ

$rs[\phi] =$ - right symbol

the rightmost symbol of ϕ

$lr[\phi] =$ - left remainder

ϕ with the rightmost symbol deleted

$rr[\phi] =$ - right remainder

ϕ with the leftmost symbol deleted

Of the four functions, rs and lr are most frequently used and it is convenient to have the following shorter notation.

$$\text{lr}[\phi] \equiv \phi^{\ell}$$

$$\text{rs}[\phi] \equiv \phi^r$$

A length function $\underline{\ell}$ from \hat{E} to N is defined by:

$$\ell[\phi] = \quad \quad \quad - \text{length}$$

$$\begin{bmatrix} \phi = \Lambda \rightarrow 0; \\ \ell[\phi^{\ell}] + 1 \end{bmatrix}$$

This definition has been given in the recursive conditional expression notation of McCarthy [MCCAJ630] which is summarised in Section 10.8.

The definition states that if $\phi = \Lambda$ then $\ell[\phi]$ has the value zero. Otherwise it takes the value $\ell[\phi^{\ell}] + 1$.

The level of parenthesisation of a word in \hat{E} is the number of pairs of parentheses surrounding that word. In the word $(a.b).(c.(c.b)).d$ the level of parenthesisation of d is zero while ab is at level 1 and cb at level 2. The element μ is used to indicate level of parenthesisation in the functions which follow. In these it is sometimes necessary when processing a word to ignore those parts which are enclosed in parentheses. This corresponds in Floyd's algorithm to the replacement of prime phrases by non-terminal symbols. To facilitate this type of operation the function skl is defined with domain $\hat{E} \times N$ and range \hat{E} . This function repeatedly replaces rightmost symbols by Λ while incrementing and decrementing μ in accord with changes in the level of parenthesisation. The process stops when μ has the value 0 and then the function takes the value of the word remaining.

Of the four functions, rs and lr are most frequently used and it is convenient to have the following shorter notation.

$$\text{lr}[\phi] \equiv \phi^{\ell}$$

$$\text{rs}[\phi] \equiv \phi^{\text{r}}$$

A length function $\underline{\ell}$ from \hat{E} to N is defined by:

$$\ell[\phi] = \begin{array}{l} \text{-- length} \end{array}$$

$$\left[\begin{array}{l} \phi = \Lambda \rightarrow 0; \\ \ell[\phi^{\ell}] + 1 \end{array} \right]$$

This definition has been given in the recursive conditional expression notation of McCarthy [MCCAJ630] which is summarised in Section 10.8.

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The level of parenthesisation of a word in \hat{E} is the number of pairs of parentheses surrounding that word. In the word $(a.b).(c.(c.b)).d$ the level of parenthesisation of d is zero while ab is at level 1 and cb at level 2. The element μ is used to indicate level of parenthesisation in the functions which follow. In these it is sometimes necessary when processing a word to ignore those parts which are enclosed in parentheses. This corresponds in Floyd's algorithm to the replacement of prime phrases by non-terminal symbols. To facilitate this type of operation the function skl is defined with domain $\hat{E} \times N$ and range \hat{E} . This function repeatedly replaces rightmost symbols by Λ while incrementing and decrementing μ in accord with changes in the level of parenthesisation. The process stops when μ has the value 0 and then the function takes the value of the word remaining.

$\text{skl}[\phi; \mu] =$

- skip level

$$\left[\begin{array}{l} \phi = \Lambda \rightarrow \Lambda; \\ \mu = 0 \rightarrow \phi; \\ \phi^r =) \rightarrow \text{skl}[\phi^\ell; \mu+1]; \\ \phi^r = (\rightarrow \text{skl}[\phi^\ell; \mu-1]; \\ \text{skl}[\phi^\ell; \mu] \end{array} \right]$$

The operation of this function may be demonstrated by applying it to the element $a, (b, (c, d))$ of \hat{E} . If this element were encountered during the analysis of a group expression, μ would have the value 1 corresponding to the presence of an unmatched left parenthesis.

$$\begin{aligned} & \text{skl}[a, (b, (c, d)); 1] \\ &= \text{skl}[a, (b, (c, d); 2] \\ &= \text{skl}[a, (b, (c,); 2] \\ &= \text{skl}[a, (b, (c; 2] \\ &= \text{skl}[a, (b, (; 2] \\ &= \text{skl}[a, (b, ; 1] \\ &= \text{skl}[a, (b; 1] \\ &= \text{skl}[a, (; 1] \\ &= \text{skl}[a, ; 0] \\ &= a, \end{aligned}$$

The function dlt (dot less than) is now defined. This has domain $\hat{E} \times \hat{E}$ and its range is the set T of truth values. In the definition the grammatical relation ' $<\cdot$ ' is used and is to be understood as follows. If ϕ_1 and $\phi_2 \in S$ then $\phi_1 <\cdot \phi_2$ if and only if the entry in the precedence table in Section 2 corresponding to the row indexed by ϕ_1 and the column

indexed by ϕ_2 contains the symbol ' \cdot '.

It may be noted that $\Lambda \prec \phi_2$ for all $\phi_2 \in S$. For strings of arbitrary length dlt takes the value T if there exists a symbol in ϕ_1 at zero level of parenthesisation which \prec some symbol of ϕ_2 at zero level of parenthesisation.

$$\text{dlt}[\phi_1; \phi_2] = \begin{array}{l} \left[\begin{array}{l} \phi_1 = \Lambda \rightarrow T; \quad \phi_2 = \Lambda \rightarrow F; \\ \phi_1 \in S \wedge \phi_2 \in S \wedge \phi_1 \prec \phi_2 \rightarrow T; \\ \phi_1 \notin S \wedge \phi_1^r =) \rightarrow \text{dlt}[\text{skl}[\phi_1^\ell; 1]; \phi_2]; \\ \phi_2 \notin S \wedge \phi_2^r =) \rightarrow \text{dlt}[\phi_1; \text{skl}[\phi_2^\ell; 1]]; \\ \phi_1 \notin S \wedge (\text{dlt}[\phi_1^\ell; \phi_2] \vee \text{dlt}[\phi_1^r; \phi_2]) \rightarrow T; \\ \phi_2 \notin S \wedge (\text{dlt}[\phi_1; \phi_2^\ell] \vee \text{dlt}[\phi_1; \phi_2^r]) \rightarrow T; \\ F \end{array} \right] \end{array} \quad \text{- dot less than}$$

Provided $\phi \in E$ the next three functions op' , lp' , and rp' define three unique subwords of the word ϕ such that

$$\text{lp}'[\phi] \text{op}'[\phi] \text{rp}'[\phi; 0] = \phi.$$

While these functions are intended to partition members of E their domains are \hat{E} or $\hat{E} \times N$ because in the process of evaluation they may recursively reference themselves with words which are not in E .

It should be noted that $\text{lp}'[(\phi)] = \Lambda$, $\text{rp}'[(\phi); 0] = (\phi)$

and $\text{op}'[(\phi)] = \Lambda$. In other words the functions lp' and rp' do not remove redundant external parentheses.

with right to left processing.

$\text{op}'[\phi] =$

- operator'

$$\left[\begin{array}{l} \phi = \Lambda \rightarrow \Lambda; \quad \phi \in S \rightarrow \Lambda; \\ \phi^r =) \rightarrow \text{op}'[\text{skl}[\phi^\ell; 1]]; \\ \text{dlt}[\phi^\ell; \phi^r] \rightarrow \text{op}'[\phi^\ell]; \\ \phi^r \end{array} \right]$$

 $\text{lp}'[\phi] =$

- left part'

$$\left[\begin{array}{l} \phi = \Lambda \rightarrow \Lambda; \quad \phi \in S \rightarrow \phi; \\ \phi^r =) \rightarrow \text{lp}'[\text{skl}[\phi^\ell; 1]]; \\ \text{dlt}[\phi^\ell; \phi^r] \rightarrow \text{lp}'[\phi^\ell]; \\ \phi^\ell \end{array} \right]$$

 $\text{rp}'[\phi; \mu] =$

- right part'

$$\left[\begin{array}{l} \phi = \Lambda \rightarrow \Lambda; \quad \phi \in S \rightarrow \Lambda; \\ \phi^r =) \rightarrow \text{rp}'[\phi^\ell; \mu+1]); \\ \phi^r = (\rightarrow \text{rp}'[\phi^\ell; \mu-1](; \\ \mu > 0 \vee \text{dlt}[\phi^\ell; \phi^r] \rightarrow \text{rp}'[\phi^\ell; \mu] \phi^r \\ \Lambda \end{array} \right]$$

It should be emphasised again that the words to be partitioned must be members of E . In this case all parentheses will be matched and invalid function references such as $\text{rp}'[(;0]$ cannot occur since when counting parentheses from the right there can never be an excess of opening parentheses over closing parentheses.

The uniqueness of the partitioning and its conformity to the structure defined by the grammar result from the use of the function dlt together with right to left processing.

An example of the action of \underline{rp}' on the expression $a, {}_2b, c.(a, b)$ is given.

$$\begin{aligned}
 & rp'[a, {}_2b, c.(a, b); 0] \\
 &= rp'[a, {}_2b, c.(a, b; 1)] \\
 &= rp'[a, {}_2b, c.(a, ; 1)b] \\
 &= rp'[a, {}_2b, c.(a ; 1), b] \\
 &= rp'[a, {}_2b, c.(; 1)a, b] \\
 &= rp'[a, {}_2b, c. ; 0](a, b) \\
 &= rp'[a, {}_2b, c ; 0].(a, b) \\
 &= rp'[a, {}_2b, ; 0]c.(a, b) \\
 &= rp'[a, {}_2b ; 0], c.(a, b) \\
 &= rp'[a, {}_2 ; 0]b, c.(a, b) \\
 &= \wedge b, c.(a, b) \\
 &= b, c.(a, b)
 \end{aligned}$$

Theorem:

$$lp'[skl[\phi; \mu]]op'[skl[\phi; \mu]]rp'[\phi; \mu] = \phi$$

for all $\phi \in \hat{E}$ and all $\mu \in N$.

Proof:

The theorem is proved by induction on the length of ϕ , that is $\ell[\phi]$. When $\ell[\phi] = 0$ the theorem is obviously true. Let n be greater than 0 and suppose that it is true for all $\psi \in \hat{E}$ such that $\ell[\psi] < n$. It is to be shown that it is true for all $\phi \in \hat{E}$ such that $\ell[\phi] = n$. The conditions stated for the various cases below are not necessarily mutually exclusive and it is assumed that case i implies not case j where $j < i$. This conforms with the conditional expression definitions.

Case 1. $\mu = 0 \wedge \phi \in S$

$$\begin{aligned}
 & 1p'[\text{skl}[\phi;\mu]]\text{op}'[\text{skl}[\phi;\mu]]\text{rp}'[\phi;\mu] \\
 &= 1p'[\phi]\text{op}'[\phi]\text{rp}'[\phi;0] \\
 &= \phi \wedge \wedge \\
 &= \phi
 \end{aligned}$$

Case 2. $\mu = 0 \wedge \phi^r =)$

$$\begin{aligned}
 & 1p'[\text{skl}[\phi;\mu]]\text{op}'[\text{skl}[\phi;\mu]]\text{rp}'[\phi;\mu] \\
 &= 1p'[\text{skl}[\phi^\ell;1]]\text{op}'[\text{skl}[\phi^\ell;1]]\text{rp}'[\phi^\ell;1]) \\
 &= \phi^\ell) \\
 &= \phi
 \end{aligned}$$

Case 3. $\mu = 0 \wedge \phi^r = ($

This cannot occur.

Case 4. $\mu = 0 \wedge \text{dlt}[\phi^\ell; \phi^r]$

$$\begin{aligned}
 & 1p'[\text{skl}[\phi;\mu]]\text{op}'[\text{skl}[\phi;\mu]]\text{rp}'[\phi;\mu] \\
 &= 1p'[\phi]\text{op}'[\phi]\text{rp}'[\phi;0] \\
 &= 1p'[\phi^\ell]\text{op}'[\phi^\ell]\text{rp}'[\phi^\ell;0]\phi^r \\
 &= \phi^\ell \phi^r \\
 &= \phi
 \end{aligned}$$

Case 5. $\mu > 0 \wedge \phi^r =)$

$$\begin{aligned}
 & 1p'[\text{skl}[\phi;\mu]]\text{op}'[\text{skl}[\phi;\mu]]\text{rp}'[\phi;\mu] \\
 &= 1p'[\text{skl}[\phi^\ell;\mu+1]]\text{op}'[\text{skl}[\phi^\ell;\mu+1]]\text{rp}'[\phi^\ell;\mu+1]) \\
 &= \phi^\ell) \\
 &= \phi
 \end{aligned}$$

Case 6. $\mu > 0 \wedge \phi^r = ($

$$\begin{aligned}
 & \text{lp}'[\text{skl}[\phi; \mu]] \text{op}'[\text{skl}[\phi; \mu]] \text{rp}'[\phi; \mu] \\
 &= \text{lp}'[\text{skl}[\phi^\ell; \mu-1]] \text{op}'[\text{skl}[\phi^\ell; \mu-1]] \text{rp}'[\phi^\ell; \mu-1] (\\
 &= \phi^\ell (\\
 &= \phi
 \end{aligned}$$

Case 7. $\mu > 0 \wedge \text{dlt}[\phi^\ell; \phi^r]$

$$\begin{aligned}
 & \text{lp}'[\text{skl}[\phi; \mu]] \text{op}'[\text{skl}[\phi; \mu]] \text{rp}'[\phi; \mu] \\
 &= \text{lp}'[\text{skl}[\phi^\ell; \mu]] \text{op}'[\text{skl}[\phi^\ell; \mu]] \text{rp}'[\phi^\ell; \mu] \phi^r \\
 &= \phi^\ell \phi^r \\
 &= \phi
 \end{aligned}$$

Case 8.

$$\begin{aligned}
 & \text{lp}'[\text{skl}[\phi; \mu]] \text{op}'[\text{skl}[\phi; \mu]] \text{rp}'[\phi; \mu] \\
 &= \text{lp}'[\text{skl}[\phi^\ell; \mu]] \text{op}'[\text{skl}[\phi^\ell; \mu]] \text{rp}'[\phi^\ell; \mu] \phi^r \\
 &= \phi^\ell \phi^r \\
 &= \phi
 \end{aligned}$$

Corollary:

$$\text{lp}'[\phi] \text{op}'[\phi] \text{rp}'[\phi; 0] = \phi \quad \text{if } \phi \in E$$

Proof:

$$\begin{aligned}
 & \text{lp}'[\phi] \text{op}'[\phi] \text{rp}'[\phi; 0] \\
 &= \text{lp}'[\text{skl}[\phi; 0]] \text{op}'[\text{skl}[\phi; 0]] \text{rp}'[\phi; 0] && \text{- by definitions} \\
 &= \phi && \text{- by the theorem}
 \end{aligned}$$

Note that the restriction on ϕ is necessary to ensure matched parentheses for skl and rp' and also to ensure a grammatical structure for dlt in case 4 of the theorem.

3.3 The parenthesising function.

A function is now given for fully parenthesising, in a non-redundant fashion, any expression in E . This function \underline{f} has as its domain the set $\hat{E} \times \hat{E} \times N$. When \underline{f} is first referenced ϕ_1 must be a member of E , $\phi_2 = \Lambda$ and $\mu = 0$. The range of \underline{f} is denoted by E^f and is the union of two disjoint subsets J^f and G^f which are respectively the fully parenthesised rational and group expressions.

$$\underline{f}[\phi_1; \phi_2; \mu] =$$

- fully parenthesise

$$\left[\begin{array}{l} \phi_1 = \Lambda \wedge \phi_2 = \Lambda \rightarrow \Lambda; \\ \phi_1 \in S \wedge \phi_2 = \Lambda \rightarrow \phi_1; \\ \phi_1 = \Lambda \wedge \phi_2^r =) \wedge 1s[\phi_2] = (\rightarrow \underline{f}[\phi_2^l; \Lambda; 0]; \\ \phi_1^r =) \rightarrow \underline{f}[\phi_1^l; \phi_2; \mu+1]; \\ \phi_1^r = (\rightarrow \underline{f}[\phi_1^l; \phi_2; \mu-1]; \\ \mu > 0 \rightarrow \underline{f}[\phi_1^l; \phi_1^r \phi_2; \mu]; \\ d1t[\phi_1^l; \phi_1^r] \rightarrow \underline{f}[\phi_1^l; \phi_1^r \phi_2; 0]; \\ \phi_1^l \in S \wedge \phi_2 \in S \rightarrow \phi_1^l \phi_1^r \phi_2; \\ \phi_1^l \in S \wedge \phi_2 \notin S \rightarrow \phi_1^l \phi_1^r (\underline{f}[\phi_2; \Lambda; 0]); \\ \phi_1^l \notin S \wedge \phi_2 \in S \rightarrow (\underline{f}[\phi_1^l; \Lambda; 0]) \phi_1^r \phi_2; \\ (\underline{f}[\phi_1^l; \Lambda; 0]) \phi_1^r (\underline{f}[\phi_2; \Lambda; 0]) \end{array} \right]$$

Note that because of the action of lines 3,4 and 5 of the definition the fully parenthesised expressions are free of redundant parentheses. The fully parenthesised form of $((a))$ is a .

As mentioned in Section 2 the multiple commas of E' are replaced in E by a subscripted comma which is interpreted as a single symbol.

In the computer procedures the process of fully parenthesising is followed by a function which deletes the subscripts from the commas.

To assist in the description of certain properties of \underline{f} the words computation and terminate are defined.

A computation of a function is a sequence of steps in the evaluation of the function. Each step consists of the replacement of one function reference by the appropriate value obtained from the definition. A computation is said to terminate if there exists a step for which there is no further function to be evaluated. The function is said to terminate finitely if the number of steps in the computation is finite. The techniques and language of computability theory [DAVIM580, HERMH650] are appropriate here but have not been used as their range of application in the text is too limited to justify the extensive theoretical modifications necessary. Changes would appear desirable since the domains and ranges most frequently encountered are subsets of \hat{E} rather than the set of positive integers used in conventional computability theory. McCarthy [MCCAJ630] has given an informal presentation of a suitable approach.

That a computation of the parenthesising process when applied to an expression in E does terminate in a finite number of steps may be seen by observing that if a function reference $f[\phi_1; \phi_2; \mu_1]$ is evaluated in any one step of a computation as $\psi_3 f[\psi_1; \psi_2; \mu_2] \psi_4$ where ψ_3 and ψ_4 are any elements of S^* then $2\ell[\psi_1] + \ell[\psi_2]$ is always less than $2\ell[\phi_1] + \ell[\phi_2]$. $\underline{\ell}$ is the previously defined length function.

Apart from showing that \underline{f} terminates it is necessary to show

that it does parenthesise in accord with the grammar. This will be demonstrated for the operation of exponentiation. Other cases can be treated similarly.

Suppose that a $\langle g\text{-factor} \rangle$, a $'\cdot'$ and a $\langle \text{rational expression} \rangle$ are concatenated as specified by rule 11 of the grammar:

$\langle g\text{-factor} \rangle \cdot \langle \text{rational expression} \rangle$

and let this resulting expression be denoted by ϕ . It will be shown that the parenthesising function partitions ϕ about the (rightmost) colon as required by the grammar.

First note that \underline{f} inspects symbols in a right to left scan skipping over parenthesised expressions and when $\mu = 0$ (i.e. when at zero level of parenthesisation) seeking out the first (and therefore rightmost) symbol which is less than (i.e. $<\cdot$) all symbols to the left of it at zero level of parenthesisation.

Now it will be seen by reference to the precedence table in Section 2 that $'\cdot'$ is dot less than any of the symbols $\ell, \dots, z, 0, \dots, 9, +, -, *, |, \wedge, /$ which, apart from parentheses, are the only symbols appearing in rational expressions. It will also be noted that \underline{f} does scan past parentheses when determining the manner of partitioning at zero level of parenthesisation.

Finally, it will be seen that all the symbols which can appear in a $\langle g\text{-factor} \rangle$ at zero level of parenthesisation, namely $a, \dots, k, v, \cdot, :$, are not dot less than $'\cdot'$.

Therefore the rightmost colon will be taken by \underline{f} as the point for partitioning of ϕ or as the 'outermost' operator of ϕ and this is the type of structure generated by the grammatical rule:

$$\langle \text{g-factor} \rangle ::= \langle \text{g-factor} \rangle \mid \langle \text{g-factor} \rangle : \langle \text{rational expression} \rangle$$

The following sample computations illustrate the effect of \underline{f}

$$\begin{aligned}
 \text{(a)} \quad & f[a, {}_2b, c; \Lambda; 0] \\
 &= f[a, {}_2b, ; c; 0] \\
 &= f[a, b; , c; 0] \\
 &= f[a, {}_2; b, c; 0] \\
 &= a, {}_2(f[b, c; \Lambda; 0]) \\
 &= a, {}_2(f[b, ; c; 0]) \\
 &= a, {}_2(b, c)
 \end{aligned}$$

$$\begin{aligned}
 \text{(b)} \quad & f[a, b, {}_2c; \Lambda; 0] \\
 &= f[a, b, {}_2; c; 0] \\
 &= (f[a, b; \Lambda; 0]), {}_2c \\
 &= (f[a, ; b; 0]), {}_2c \\
 &= (a, b), {}_2c
 \end{aligned}$$

4. CANONIC EXPRESSIONS.

4.1 Introduction.

In Section 3 the method for converting expressions in E into the fully parenthesised form E^f was described. In this section and those to follow the set E^f will be of major interest. The functions of this section serve to define a subset of E^f the expressions of which will be called canonic.

There are two distinguishable aspects to the definition. The first is associated with the structure or relationship of operators of expressions while the second is associated with the relationship of operands to one other. Structurally, canonic expressions take the form of products of powers of commutators. That is, the innermost or first executed operators must be commutator operators and the outermost operators must be product operators. A similar type of structure is required of the exponents themselves. The second aspect of the definition is concerned with the ordering of expressions and cannot be simply described.

The following extracts from Section 10 are relevant here.

- (a) '[' and ']' are meta-brackets
- (b) Greek lower case letters, possibly subscripted, are used as variables.
- (c) Roman upper case script letters denote sets
- (d) ';' is the meta-delimiter
- (e) When referred to in text non alpha-numeric symbols are surrounded by single quotes.

The set notation $\{<x>\}$ used here is intended to denote the set of all derivations of the meta-variable $<x>$.

<u>Set Name</u>	<u>Description</u>
N'	$\{<\text{integer}>\} \cup \{\infty\}$
Q	$\{<\text{rational number}>\}$
B	$\{<\text{rational expression generator}>\}$
R	$\{B \cup Q\}$
J	$\{<\text{rational expression}>\}$
J^f	fully parenthesised subset of J
L	$\{<\text{group expression generator}>\}$
G	$\{<\text{group expression}>\}$
G^f	fully parenthesised subset of G
E	$\{G \cup J\}$
E^f	fully parenthesised subset of E
O	$\{::/; ;*;+;-;.,;,\wedge\}$
T	$\{T;F\}$

The domains of the variables (possibly subscripted) to be used are:

<u>Variable</u>	<u>Domain</u>
X	O
ρ	G^f
σ	J^f
τ	E^f

The following additional conventions and notations are used.

- (f) $'<'$, $'>'$ and $'='$ are relations on the integers or rational members and have the normal meaning.

- (g) '=' is assumed to have an obvious meaning in comparing any expressions.
- (h) ' $\overset{A}{\lt}$ ' and ' $\overset{A}{\gt}$ ' are standard alphabetic orders.

4.2 Preliminary functions.

In many cases the functions to be defined are mutually recursive and for this reason it is not always possible to avoid reference to a particular function before it is defined. The definitions begin with three functions of basic importance. These are the operator function and the left and right part functions. These functions have for their domain the set of fully parenthesised expressions and they divide such expressions into the three obvious parts. They are defined in terms of the previously given functions op', lp' and rp' and differ from them in that the unprimed functions remove the pair of external parentheses if they are present. These functions are defined in the notation of recursive conditional expressions described in Section 10.8.

$op[\tau] = op'[\tau]$ - operator

$lp[\tau] =$ - left part

$$\left[\begin{array}{l} ls[lp'[\tau]] = (\wedge rs[lp'[\tau]] =) \rightarrow lr[rr[lp'[\tau]]]; \\ lp'[\tau] \end{array} \right]$$

$rp[\tau] =$ - right part

$$\left[\begin{array}{l} ls[rp'[\tau;0]] = (\wedge rs[rp'[\tau;0]] =) \rightarrow lr[rr[rp'[\tau;0]]]; \\ rp'[\tau] \end{array} \right]$$

length[σ] = - length

$$\left[\begin{array}{l} \sigma \in R \rightarrow 1; \\ \text{length}[lp[\sigma]] + \text{length}[rp[\sigma]] \end{array} \right]$$

The domain of this function is the set J^f and its range is the set N . It determines the number of operands in a rational expression. The function prd (product) has range T and takes the value T if and only if the expression σ is free of operators other than '*' or '^'. Function rat (rational) also has range T and takes the value T if and only if the argument is a rational number or has imbedded within it a rational number.

prd[σ] = - product

$$\left[\begin{array}{l} \sigma \in R \rightarrow T; \\ (\text{op}[\sigma] = * \vee \text{op}[\sigma] = \wedge) \wedge \text{prd}[lp[\sigma]] \wedge \text{prd}[rp[\sigma]] \rightarrow T; \\ F \end{array} \right]$$

rat[σ] = - rational

$$\left[\begin{array}{l} \sigma \in Q \rightarrow T; \\ \text{rat}[lp[\sigma]] \vee \text{rat}[rp[\sigma]] \end{array} \right]$$

Functions lo (left operand) and ro (right operand) have for their domain the set $E^f \times O$ and their range is the set E^f . They select respectively the leftmost operand and the rightmost operand of an expression with respect to an operator. If, for example, $\tau = ((a,b).(c,d)).((e,f).g)$ then $lo[\tau;.] = a,b$ and $ro[\tau;.] = g$.

$lo[\tau;X] =$ - left operand

$$\begin{bmatrix} \tau \in RUL \rightarrow \tau; \\ op[\tau] \neq X \rightarrow \tau; \\ lo[lp[\tau];X] \end{bmatrix}$$

$ro[\tau;X] =$ - right operand

$$\begin{bmatrix} \tau \in RUL \rightarrow \tau; \\ op[\tau] \neq X \rightarrow \tau; \\ ro[rp[\tau];X] \end{bmatrix}$$

The next two functions to be given are complementary to lo and ro. Thus if τ has the value given in the example above then $le[\tau;.] = ((a,b).(c,d)).(e,f)$. These functions differ from all previous functions in that they produce the desired value by composition or concatenation while the previous functions have used decomposition only.

$le[\tau;X] =$ - left expression

$$\begin{bmatrix} op[\tau] \neq X \rightarrow \tau; \\ rp[\tau] \in RUL \vee op[rp[\tau]] \neq X \rightarrow lp[\tau]; \\ lp'[\tau]X(le[rp[\tau];X]) \end{bmatrix}$$

$re[\tau;X] =$ - right expression

$$\begin{bmatrix} op[\tau] \neq X \rightarrow \tau; \\ lp[\tau] \in RUL \vee op[lp[\tau]] \neq X \rightarrow rp[\tau]; \\ (re[lp[\tau];X])Xrp'[\tau;0] \end{bmatrix}$$

4.3 Canonic rational expressions.

With the auxiliary functions now defined the five functions

cp (canonic product), ltp (less than product), lpr (left product), rpr (right product) and ce (canonic expression) can be given.

Together these constitute the definition of a canonic exponent. For simplicity only three operators are permitted in a canonic exponent. These are the sum, difference and product operators. Structurally the canonic exponent is in the form of sums of products and since exponents themselves may not have exponents the products may consist of repetitions of the same rational expression generator.

Within products ordinary alphabetic ordering is used for rational expression generators. The rational number coefficient, if any, appears at the right hand end of the product. The ordering of operands of a summation is done first by length of the operands and for operands of the same length alphabetic ordering is used. The unary minus operation is not permitted.

$cp[\sigma] =$ - canonic product

$$\left[\begin{array}{l} \sigma \in R \rightarrow T; \\ \sim \text{prd}[\sigma] \rightarrow F; \\ \text{rat}[\text{le}[\sigma; *]] \rightarrow F; \\ \sim [\text{cp}[\text{lp}[\sigma]] \wedge \text{cp}[\text{rp}[\sigma]]] \rightarrow F; \\ \text{ro}[\text{lp}[\sigma]; *] \stackrel{A}{\leq} \text{lo}[\text{rp}[\sigma]; *] \rightarrow T; \\ F \end{array} \right]$$

$\text{ltpr}[\sigma_1; \sigma_2] =$

- less than product

$$\left[\begin{array}{l} \sigma_1 \in Q \wedge \sigma_2 \in Q \rightarrow F; \\ \sigma_1 \in B \wedge \sigma_2 \in R \rightarrow T; \\ \text{ro}[\sigma_1; *] \in Q \rightarrow \text{ltpr}[\text{le}[\sigma_1; *]; \sigma_2]; \\ \text{ro}[\sigma_2; *] \in Q \rightarrow \text{ltpr}[\sigma_1; \text{le}[\sigma_2; *]]; \\ \text{length}[\sigma_1] < \text{length}[\sigma_2] \rightarrow T; \\ \text{lo}[\sigma_1; *] \stackrel{A}{<} \text{lo}[\sigma_2; *] \rightarrow T; \\ \text{ltpr}[\text{re}[\sigma_1; *]; \text{re}[\sigma_2; *]] \end{array} \right]$$

 $\text{lpr}[\sigma] =$

- left product

$$\left[\begin{array}{l} \sigma \in R \rightarrow \sigma; \\ \text{op}[\sigma] = * \rightarrow \sigma; \\ \text{op}[\sigma] = + \rightarrow \text{lpr}[\text{lo}[\sigma; +]]; \\ \text{op}[\sigma] = - \rightarrow \text{lpr}[\text{lo}[\sigma; -]] \end{array} \right]$$

 $\text{rpr}[\sigma] =$

- right product

$$\left[\begin{array}{l} \sigma \in R \rightarrow \sigma; \\ \text{op}[\sigma] = * \rightarrow \sigma; \\ \text{op}[\sigma] = + \rightarrow \text{rpr}[\text{ro}[\sigma; +]]; \\ \text{op}[\sigma] = - \rightarrow \text{rpr}[\text{ro}[\sigma; -]] \end{array} \right]$$

$ce[\sigma] =$ - canonic exponent

$$\left[\begin{array}{l} \sigma \in R \rightarrow T; \\ \text{prd}[\sigma] \wedge \text{cp}[\sigma] \rightarrow T; \\ \text{nce}[\text{lp}[\sigma]] \rightarrow F; \\ \text{nce}[\text{rp}[\sigma]] \rightarrow F; \\ \text{ltpr}[\text{rpr}[\text{lp}[\sigma]]; \text{lpr}[\text{rp}[\sigma]]] \rightarrow T; \\ F \end{array} \right]$$

4.4 Canonic group expressions

It is now convenient to define the weight of a group expression by the function wt which maps the group expressions G^f to the set N' of integers and the symbol ' ∞ '. It is to be understood that in all arithmetic involving the set N' the following rules apply:

$$(a) \quad i + \infty = \infty$$

$$(b) \quad i - \infty = 0$$

$$(c) \quad \infty \pm i = \infty$$

$$(d) \quad \infty + \infty = \infty$$

$$(e) \quad \infty - \infty = 0$$

$wt[\rho] =$ - weight

$$\left[\begin{array}{l} \rho = v \rightarrow \infty; \quad \rho \in L \rightarrow 1; \\ \text{op}[\rho] = : \rightarrow \text{wt}[\text{lp}[\rho]]; \\ \text{op}[\rho] = . \rightarrow \\ \quad \left[\begin{array}{l} \text{wt}[\text{lp}[\rho]] < \text{wt}[\text{rp}[\rho]] \rightarrow \text{wt}[\text{lp}[\rho]]; \\ \text{wt}[\text{rp}[\rho]] \end{array} \right]; \\ \text{op}[\rho] = , \rightarrow \text{wt}[\text{lp}[\rho]] + \text{wt}[\text{rp}[\rho]] \end{array} \right]$$

Note that in this function there is an imbedded or nested conditional expression. It is evaluated in the same way as any of the conditional expressions which in other cases have been used to define functions. Function wt acts as a coarse order on group expressions.

The binary function lt (less than) uses the function wt and two auxiliary functions to define an order on group expressions. The domain of lt is the set $G^f \times G^f$ and its range is the set T . The auxiliary functions are ld and tr.

$ld[\rho] =$ - leading part

$$[\underline{lt}[lp[\rho];rp[\rho]] \rightarrow rp[\rho];lp[\rho]]$$

$tr[\rho] =$ - trailing part

$$[\underline{lt}[lp[\rho];rp[\rho]] \rightarrow lp[\rho];rp[\rho]]$$

$lt[\rho_1; \rho_2] =$

- less than

$$\left[\begin{array}{l} \rho_1 \in L \wedge \rho_2 \in L \wedge \rho_1 \overset{A}{<} \rho_2 \rightarrow T; \\ wt[\rho_1] < wt[\rho_2] \rightarrow T; \\ wt[\rho_1] > wt[\rho_2] \rightarrow F; \\ op[\rho_1] = : \rightarrow lt[lp[\rho_1]; \rho_2]; \\ op[\rho_2] = : \rightarrow lt[\rho_1; lp[\rho_2]]; \\ op[\rho_1] = . \wedge lt[lp[\rho_1]; \rho_2] \wedge lt[rp[\rho_1]; \rho_2] \rightarrow T; \\ op[\rho_2] = . \wedge lt[\rho_1; lp[\rho_2]] \wedge lt[\rho_1; rp[\rho_2]] \rightarrow T; \\ lt[ld[\rho_1]; ld[\rho_2]] \rightarrow T; \\ ld[\rho_1] = ld[\rho_2] \wedge lt[tr[\rho_1]; tr[\rho_2]] \rightarrow T; \\ ld[\rho_1] = ld[\rho_2] \wedge tr[\rho_1] = tr[\rho_2] \\ \quad \wedge lt[lp[\rho_2]; lp[\rho_1]] \rightarrow T; F \end{array} \right]$$

Function com takes the value T if the argument whose domain is the set G^f has no operator other than ', '. Such an expression is a commutator.

 $com[\rho] =$

- commutator

$$\left[\begin{array}{l} \rho \in L \rightarrow T; \\ op[\rho] = , \wedge com[lp[\rho]] \wedge com[rp[\rho]] \rightarrow T; \\ F \end{array} \right]$$

The definition of the function canonic can now be given. It has domain G^f and range T. This function takes the value T for arguments from a certain subset of G^f . Such expressions are said to be canonic. As remarked earlier the significance and uniqueness of a similar canonic form are dealt with by Ward.

canonic[ρ] =

- canonic

$$\begin{array}{l}
 \rho \in L \rightarrow T; \\
 \sim \text{canonic}[\text{lp}[\rho]] \rightarrow F; \\
 \text{op}[\rho] = : \wedge \text{com}[\text{lp}[\rho]] \wedge \text{ce}[\text{rp}[\rho]] \rightarrow T; \\
 \sim \text{canonic}[\text{rp}[\rho]] \rightarrow F; \\
 \text{op}[\rho] = . \wedge \text{lt}[\text{lp}[\rho]; \text{rp}[\rho]] \rightarrow T; \\
 \text{com}[\rho] \rightarrow \\
 \quad \begin{array}{l}
 \text{lt}[\text{rp}[\rho]; \text{lp}[\rho]] \rightarrow T; \\
 \text{wt}[\text{lp}[\rho]] > 1 \wedge [\text{lt}[\text{rp}[\text{lp}[\rho]]; \text{rp}[\rho]] \\
 \quad \vee \text{rp}[\text{lp}[\rho]] = \text{rp}[\rho]] \\
 \rightarrow T; \\
 \text{F}
 \end{array} \\
 \text{F}
 \end{array}$$

5. ROLE OF THE MODULUS.

In Sections 6 and 7 the collecting process proper will be described. It will be shown to consist of a number of transformation rules together with a description of the manner in which they are to be applied. There is one rule, however, which is so central to the process that it will be discussed separately. This rule is concerned with an integer called the modulus which is associated with any group expression that is to be collected or transformed to the canonic form. The value of the integer is supplied by the mathematician together with the group expression to be collected. The mathematical significance of this modulus is not discussed here but may be found in P. Hall [HALLP340] and Magnus, Karrass and Solitar [MAGNW660].

The computational significance of the modulus is that all commutators of weight greater than or equal to the modulus 'reduce' to the identity element or 'are equivalent' to that element. This process is analogous to modular arithmetic. It requires careful application for two reasons. The first is that timely application will effect a vast improvement in speed of calculation and the other is that delayed application may result in a non-terminating process. The Jacobi transformation (see Section 6) would be susceptible to this latter problem.

Ideally then all transformations should be regarded as applicable modulo a certain weight (called the modulus) and in fact terms which would later reduce to the identity should not be generated in the

first place. While this is the approach of the computer procedures the description here delays the reduction process in the interests of clarity.

There is a further aspect of modular transformations. To discuss this it is convenient to use the word subexpression. This is a part of an expression which is itself an expression. Now while certain subexpressions have weight less than the modulus supplied for the main expression they may eventually contribute sufficient weight to a subexpression containing them to cause the generation of expressions whose weight is not less than the modulus of the main expression. An example of this is given. A hand collection of ab,c,d modulo weight 4 might take the form:

$$\begin{aligned}
 ab,c,d &= (a,b)(a,c,b)(b,c),d \\
 &= (a,b,d)(a,b,d,(a,c,b)(b,c))((a,c,b)(b,c),d) \\
 &= (a,b,d) \vee ((a,c,b)(b,c),d) \\
 &= (a,b,d)((a,c,b)(b,c),d) \\
 &= (a,b,d)(a,c,b,d)(a,c,b,d,,b,c)(b,c,d) \\
 &= (a,b,d) \vee \vee (b,c,d) \\
 &= (a,b,d) \vee (b,c,d) \\
 &= (a,b,d)(b,c,d)
 \end{aligned}$$

A more efficient method is to recognise in the second line that the subexpression a,c,b in its context is already 'too heavy' and may be replaced by the identity even though its weight is only 3. Then the calculation is as follows:

$$\begin{aligned}
ab,c,d &= (a,b)(a,c,b)(b,c),d \\
&= (a,b)\vee(b,c),d \\
&= (a,b)(b,c),d \\
&= (a,b,d)(a,b,d,,b,c)(b,c,d) \\
&= (a,b,d)\vee(b,c,d) \\
&= (a,b,d)(b,c,d)
\end{aligned}$$

While this example illustrates the principle it does not demonstrate the immense practical importance of this method. In more complex expressions practical computation would not be possible unless this method were used.

In the following functions the reduction process is formalised. Function modl enables the calculation of a submodulus for the left part of an expression whose modulus is known. Function modr operates similarly for the right part. The domain of these functions is the set $G^f \times N'$ and their range is the set N' . The process of replacing expressions whose weight is not less than the modulus of that expression by the identity is performed by function red.

$$\text{modl}[\tau;\mu] = \quad \text{- left modulus}$$

$$\left[\begin{array}{l}
\text{op}[\tau] = . \rightarrow \mu; \\
\text{op}[\tau] = : \rightarrow \mu; \\
\text{op}[\tau] = , \rightarrow \mu - \text{wt}[\text{rp}[\tau]]
\end{array} \right]$$

$$\text{modr}[\tau;\mu] = \quad \text{- right modulus}$$

$$\left[\begin{array}{l}
\text{op}[\tau] = . \rightarrow \mu; \\
\text{op}[\tau] = : \rightarrow \infty; \\
\text{op}[\tau] = , \rightarrow \mu - \text{wt}[\text{lp}[\tau]]
\end{array} \right]$$

$\text{red}[\tau; \mu] =$ 5. TRANSFORMATION RULES. - reduce

$$\left[\begin{array}{l} \tau \in L \rightarrow \tau; \\ \tau \in J^f \rightarrow \tau; \\ \text{wt}[\tau] \geq \mu \rightarrow \nu; \\ \text{red}[lp'[\tau]; \mu] \text{op}[\tau] \text{red}[rp'[\tau; 0]; \mu] \end{array} \right]$$

6. TRANSFORMATION RULES.

6.1 Introduction.

In this section the transformation rules to be used for converting an expression of E^f into the canonic form are given. These rules are classified broadly on the basis of the lowest precedence operation appearing in an expression before it is transformed. So the function trp is concerned with product transformations. The classification is arbitrary from a description viewpoint but it does lead to more efficient manipulation.

The predicates appearing in the functions take the value T if the expression on the left of the equality has the form of the right hand side of the equality. For example the predicate $\rho = \rho_1 :- \text{!}\rho_2$ has the value T if and only if there exist expressions ρ_1 and ρ_2 such that $\rho_1 :- \text{!}\rho_2$ is the same expression as ρ . If such a predicate has the value T the function takes as its value the expression to the right of the arrow. Each such component of the function can be regarded as a transformation rule.

In many cases the function scan appears within the transformation rule. This function will be defined in the next section. It should suffice to state here that this function maps an element from the domain $E^f \times N'$ into the subset of E^f consisting of canonic expressions. It is applied within the transformation function to ensure that all subexpressions produced by the transformation are converted to canonic form.

In the transformation functions the domain of the variable ρ , possibly subscripted, is the set of fully parenthesised group expressions G^f and the domain of σ is the set of fully parenthesised rational expressions J^f .

While it is required that the arguments of the functions be fully parenthesised the transformation function is permitted to generate non fully parenthesised expressions. This has been done to maintain clarity and simplicity. The reason for specifying the restricted domain for the arguments is also that of simplicity. If non parenthesised expressions were permitted the transformations involving commutation would have a complexity similar to that of the relevant grammar. The middle course chosen is convenient for another reason - to the mathematician the non-parenthesised expression or non fully parenthesised expression is more easily interpreted while, within the computer, the fully parenthesised expression, particularly as a list tree, is more easily manipulated. Thus the expressions specifying the transformations must be fully parenthesised in a manner similar to that described in Section 3 for the input expressions themselves. In practice while the parenthesising process is applied automatically to the input expressions the equivalent operation on the transformations has been done once and for all by hand.

It is a necessary assumption in what follows that the left and right parts of the expressions to be transformed are in canonic form. Section 7 will show how this condition is ensured.

For simplicity the unary minus operation has not been included in the grammar of group expressions or in the function \underline{f} . When unary minus is used in the transformations in the form $'-\sigma'$ it should be understood as $0-\sigma$. Unary operations are permitted in the computer procedures.

6.2 Group expression transformation.

The group product transformation \underline{trp} has domain $G^f \times N'$.

$\underline{trp}[\rho; \mu] =$ - transform product

$$\left[\begin{array}{l} \rho = \rho_1 : -1\rho_1 \rightarrow v; \quad \rho = \rho_1\rho_1 : -1 \rightarrow v; \\ \rho = \rho_1 v \rightarrow \rho_1; \quad \rho = v\rho_1 \rightarrow \rho_1; \\ \rho = \rho_1\rho_2 \rightarrow \rho_1:2; \\ \rho = \rho_1:\sigma\rho_1 \rightarrow \text{scan}[\rho_1:(\sigma+1);\mu]; \\ \rho = \rho_1\rho_1:\sigma \rightarrow \text{scan}[\rho_1:(\sigma+1);\mu]; \\ \rho = \rho_1:\sigma_1\rho_1:\sigma_2 \rightarrow \text{scan}[\rho_1:(\sigma_1+\sigma_2);\mu]; \\ \rho = \rho_1\rho_2 \wedge \text{lt}[\text{lo}[\rho_2;.\];\text{ro}[\rho_1;.]] \\ \quad \rightarrow \text{scan}[\text{le}[\rho_1;.]\text{lo}[\rho_2;.]\text{ro}[\rho_1;.] \\ \quad \quad (\text{ro}[\rho_1;.]\text{lo}[\rho_2;.])\text{re}[\rho_2;.];\mu]; \\ \rho \end{array} \right]$$

Some explanation of the eighth rule may assist understanding.

If ρ_1 and ρ_2 are irreducible to the product operation and if $\text{lt}[\rho_2;\rho_1]$ then \underline{trp} is to take the value $\text{scan}[\rho_2\rho_1(\rho_1,\rho_2);\mu]$. On the other hand if ρ_1 is reducible and ρ_2 is irreducible ρ_1 may be written as $\rho_3\rho_4$ where ρ_4 is irreducible. In this case the function

takes the value $\text{scan}[\rho_3 \rho_2 \rho_4(\rho_4, \rho_2); \mu]$ if $\text{lt}[\rho_2; \rho_4]$. The other two possibilities are treated in a similar manner.

The commutation transformations are grouped in the function trc which has a similar domain and range to trp.

trc $[\rho; \mu] =$ - transform commutator

$$\begin{aligned} \rho &= \rho_1, \rho_1 \rightarrow v; \quad \rho = \rho_1, v \rightarrow v; \quad \rho = v, \rho_1 \rightarrow v; \\ \rho &= \rho_1:-1, \rho_2 \rightarrow \text{scan}[(\rho_1, \rho_2):-1(\rho_2, \rho_1, \rho_1:-1); \mu]; \\ \rho &= \rho_1, \rho_2:-1 \rightarrow \text{scan}[(\rho_1, \rho_2):-1(\rho_2, \rho_1, \rho_2:-1); \mu]; \\ \rho &= \rho_1 \rho_2, \rho_3 \rightarrow \text{scan}[(\rho_1, \rho_3)(\rho_1, \rho_3, \rho_2)(\rho_2, \rho_3); \mu]; \\ \rho &= \rho_1, \rho_2 \rho_3 \rightarrow \text{scan}[(\rho_1, \rho_3)(\rho_1, \rho_2)(\rho_1, \rho_2, \rho_3); \mu]; \\ \rho &= \rho_1, \rho_2 \wedge \text{lt}[\rho_1; \rho_2] \wedge \text{com}[\rho_1] \wedge \text{com}[\rho_2] \rightarrow (\rho_2, \rho_1):-1; \\ \rho &= \rho_3, \rho_2, \rho_1 \wedge \text{lt}[\rho_1; \rho_2] \wedge \text{lt}[\rho_2; \rho_3] \wedge \text{com}[\rho_1] \\ &\quad \wedge \text{com}[\rho_2] \wedge \text{com}[\rho_3] \\ &\quad \rightarrow \text{scan}[(\rho_1, \rho_3, \rho_3, \rho_2)(\rho_3, \rho_2, \rho_2, \rho_1, \rho_3) \\ &\quad (\rho_2, \rho_1, \rho_3):-1(\rho_3, \rho_2, \rho_2, \rho_1)(\rho_2, \rho_1, \rho_1, \rho_3, \rho_2) \\ &\quad (\rho_1, \rho_3, (\rho_3, \rho_1, \rho_2):-1)(\rho_3, \rho_1, \rho_2) \\ &\quad (\rho_2, \rho_1, \rho_1, \rho_3)(\rho_1, \rho_3, \rho_3, \rho_2, \rho_1); \mu]; \\ \rho &= \rho_1, \rho_2:\sigma \rightarrow \text{scan}[(\rho_2, \rho_1):-\sigma(\rho_2, \rho_1, \rho_2):-\sigma|2) \\ &\quad (\rho_2, \rho_1, \rho_2, \rho_2):-\sigma|3)(\rho_2, \rho_1, \rho_2, \rho_2, \rho_2):-\sigma|4) \\ &\quad (\rho_2, \rho_1, \rho_2, \rho_2, \rho_2, \rho_2):-\sigma|5) \\ &\quad (\rho_2, \rho_1, \rho_2, \rho_2, \rho_1):-\sigma+1|3) \\ &\quad (\rho_2, \rho_1, \rho_2, \rho_2, \rho_2, \rho_1):-\sigma+1|4); \mu]; \\ \rho &= \rho_2:\sigma, \rho_1 \rightarrow \text{scan}[(\rho_2, \rho_1):\sigma(\rho_2, \rho_1, \rho_2):(\sigma|2)(\rho_2, \rho_1, \rho_2, \rho_2):(\sigma|3) \\ &\quad (\rho_2, \rho_1, \rho_2, \rho_2, \rho_2):(\sigma|4)(\rho_2, \rho_1, \rho_2, \rho_2, \rho_2, \rho_2):(\sigma|5) \\ &\quad (\rho_2, \rho_1, \rho_2, \rho_2, \rho_1):(\sigma|3+\sigma+1|3) \\ &\quad (\rho_2, \rho_1, \rho_2, \rho_2, \rho_2, \rho_1):(\sigma|4+2*(\sigma+1)|4); \mu]; \end{aligned}$$

The ninth transformation here is known as the Jacobi identity. The next two transformations are valid only when $\mu \leq 7$. This is not a serious limitation as it is possible to provide a transformation which is valid up to any desired value of μ . Krause [KRAUE640] has described a similar calculation for $(\rho_1 \rho_2) : \sigma$. It is also possible to provide a transformation which is valid for arbitrarily large values of μ provided $\sigma \in N$.

It may be noted that in this as in other transformation functions if a match is not achieved the function takes the value of the expression supplied as an argument.

The function for transforming exponent expressions has no new features and is given without further comment.

$\text{tre}[\rho; \mu] =$

- transform exponent

$$\begin{aligned}
 \rho &= \rho_1:0 \rightarrow v; \rho = \rho_1:1 \rightarrow \rho_1; \rho = v:\sigma \rightarrow v; \\
 \rho &= \rho_1:\sigma_1:\sigma_2 \rightarrow \text{scan}[\rho_1:(\sigma_1*\sigma_2);\mu]; \\
 \rho &= (\rho_1\rho_2):-1 \rightarrow \text{scan}[\rho_2:-1\rho_1:-1;\mu]; \\
 \rho &= (\rho_1,\rho_2):\sigma \rightarrow \text{scan}[\rho_1:\sigma\rho_2:(\rho_2,\rho_1):(\sigma|2)(\rho_2,\rho_1,\rho_1):(\sigma|3) \\
 &\quad (\rho_2,\rho_1,\rho_2):(\sigma|2+2*\sigma|3)(\rho_2,\rho_1,\rho_1,\rho_1):(\sigma|4) \\
 &\quad (\rho_2,\rho_1,\rho_1,\rho_2):(2*\sigma|3+3*\sigma|4) \\
 &\quad (\rho_2,\rho_1,\rho_2,\rho_2):(2*\sigma|3+3*\sigma|4) \\
 &\quad (\rho_2,\rho_1,\rho_1,\rho_1):(\sigma|5)(\rho_2,\rho_1,\rho_1,\rho_1,\rho_2):(3*\sigma|4+4*\sigma|5) \\
 &\quad (\rho_2,\rho_1,\rho_1,\rho_2,\rho_2):(\sigma|3+6*\sigma|4+6*\sigma|5) \\
 &\quad (\rho_2,\rho_1,\rho_2,\rho_2,\rho_2):(3*\sigma|4+4*\sigma|5) \\
 &\quad (\rho_2,\rho_1,\rho_1,,\rho_2,\rho_1):(\sigma|3+7*\sigma|4+6*\sigma|5) \\
 &\quad (\rho_2,\rho_1,\rho_2,,\rho_2,\rho_1):(6*\sigma|3+18*\sigma|4+12*\sigma|5) \\
 &\quad (\rho_2,\rho_1,\rho_1,\rho_1,\rho_1,\rho_1):(\sigma|6) \\
 &\quad (\rho_2,\rho_1,\rho_1,\rho_1,\rho_1,\rho_2):(4*\sigma|5+5*\sigma|6) \\
 &\quad (\rho_2,\rho_1,\rho_1,\rho_1,\rho_2,\rho_2):(3*\sigma|4+12*\sigma|5+10*\sigma|6) \\
 &\quad (\rho_2,\rho_1,\rho_1,\rho_2,\rho_2,\rho_2):(3*\sigma|4+12*\sigma|5+10*\sigma|6) \\
 &\quad (\rho_2,\rho_1,\rho_2,\rho_2,\rho_2,\rho_2):(4*\sigma|5+5*\sigma|6) \\
 &\quad (\rho_2,\rho_1,\rho_1,\rho_1,,\rho_2,\rho_1):(3*\sigma|4+13*\sigma|5+10*\sigma|6) \\
 &\quad (\rho_2,\rho_1,\rho_1,\rho_2,,\rho_2,\rho_1): \\
 &\quad \quad (2*\sigma|3+24*\sigma|4+52*\sigma|5+30*\sigma|6) \\
 &\quad (\rho_2,\rho_1,\rho_2,\rho_2,,\rho_2,\rho_1): \\
 &\quad \quad (3*\sigma|3+27*\sigma|4+54*\sigma|5+30*\sigma|6) \\
 &\quad (\rho_2,\rho_1,\rho_2,,\rho_2,\rho_1,\rho_1): \\
 &\quad \quad (4*\sigma|3+21*\sigma|4+36*\sigma|5+20*\sigma|6);\mu];
 \end{aligned}$$

6.3 Rational expression transformations

The remaining transformation functions have as their domain the set J^f . They use the functions prod, sum, dif and invfac which have as their domain either the set Q or $Q \times Q$. Their range is the set Q . These are regarded as basic functions and detailed definition is not given. Details of such functions may be found in Davis [DAVIM580] or Markov [MARKA620]. In simple terms prod multiplies rational numbers, sum adds them and dif subtracts them. Function invfac evaluates the factorial of the argument and then inverts it.

Function trm transforms multiplicative expressions, trs transforms sum expressions and trd transforms difference expressions.

trm[σ] =

- transform multiply

$$\begin{aligned}
 & \left[\begin{array}{l}
 \sigma = \sigma_1 * 1 \rightarrow \sigma_1; \quad \sigma = 1 * \sigma_1 \rightarrow \sigma_1; \\
 \sigma = \sigma_1 * 0 \rightarrow 0; \quad \sigma = 0 * \sigma_1 \rightarrow 0; \\
 \sigma = -\sigma_1 * -\sigma_2 \rightarrow \text{scan}[\sigma_1 * \sigma_2; 0]; \\
 \sigma = -\sigma_1 * \sigma_2 \rightarrow \text{scan}[-(\sigma_1 * \sigma_2); 0]; \\
 \sigma = \sigma_1 * -\sigma_2 \rightarrow \text{scan}[-(\sigma_1 * \sigma_2); 0]; \\
 \sigma = \sigma_1 * (\sigma_2 + \sigma_3) \rightarrow \text{scan}[\sigma_1 * \sigma_2 + \sigma_1 * \sigma_3; 0]; \\
 \sigma = \sigma_1 * (\sigma_2 - \sigma_3) \rightarrow \text{scan}[\sigma_1 * \sigma_2 - \sigma_1 * \sigma_3; 0]; \\
 \sigma = (\sigma_1 + \sigma_2) * \sigma_3 \rightarrow \text{scan}[\sigma_1 * \sigma_3 + \sigma_2 * \sigma_3; 0]; \\
 \sigma = (\sigma_1 - \sigma_2) * \sigma_3 \rightarrow \text{scan}[\sigma_1 * \sigma_3 - \sigma_2 * \sigma_3; 0]; \\
 \sigma = \sigma_1 * \sigma_2 \wedge \sigma_1 \in Q \wedge \sigma_2 \in Q \rightarrow \text{prod}[\sigma_1; \sigma_2]; \\
 \sigma = \sigma_1 * \sigma_2 \wedge \text{ro}[\sigma_1; *]^A \text{lo}[\sigma_2; *] \\
 \quad \rightarrow \text{scan}[\text{le}[\sigma_1; *] \text{lo}[\sigma_2; *] \text{ro}[\sigma_1; *] \text{re}[\sigma_2; *]; 0]; \\
 \sigma
 \end{array} \right]
 \end{aligned}$$

`trs[σ] =` transform the transformation functions `defl` - transform sum

$$\left[\begin{array}{l} \sigma = \sigma_1 + 0 \rightarrow \sigma_1; \quad \sigma = 0 + \sigma_1 \rightarrow \sigma_1; \\ \sigma = \sigma_1 + \sigma_2 \wedge \sigma_1 \in \mathcal{Q} \wedge \sigma_2 \in \mathcal{Q} \rightarrow \text{sum}[\sigma_1; \sigma_2]; \\ \sigma = \sigma_1 + -\sigma_2 \rightarrow \text{scan}[\sigma_1 - \sigma_2; 0]; \\ \sigma = -\sigma_1 + \sigma_1 \rightarrow 0; \\ \sigma = \sigma_1 + \sigma_1 \rightarrow \sigma_1 * 2; \\ \sigma = \sigma_1 + \sigma_2 \wedge \text{lhs}[\sigma_2; \sigma_1] \rightarrow \sigma_2 + \sigma_1; \\ \sigma \end{array} \right]$$

`trd[σ] =` - transform difference

$$\left[\begin{array}{l} \sigma = \sigma_1 - \sigma_1 \rightarrow 0; \\ \sigma = \sigma_1 - 0 \rightarrow \sigma_1; \quad \sigma = 0 - \sigma_1 \rightarrow -\sigma_1; \\ \sigma = \sigma_1 - \sigma_2 \wedge \sigma_1 \in \mathcal{Q} \wedge \sigma_2 \in \mathcal{Q} \rightarrow \text{dif}[\sigma_1; \sigma_2]; \\ \sigma = \sigma_1 - \sigma_2 \wedge \text{lhs}[\sigma_2; \sigma_1] \rightarrow -\sigma_2 + \sigma_1; \\ \sigma \end{array} \right]$$

Combinatorial expressions are transformed by function trb and the function syp which generates a symbolic product.

`syp[σ; μ1; μ2] =` - symbolic product

$$\left[\begin{array}{l} \mu_2 = \mu_1 \rightarrow \sigma; \\ \sigma * \text{syp}[(\sigma - \mu_2); \mu_1; \text{sum}[\mu_2; 1]] \end{array} \right]$$

`trb[σ] =` - transform combination

$$[\sigma = \sigma_1 | \sigma_2 \rightarrow \text{syp}[\sigma_1; \sigma_2; 1] * \text{invfac}[\sigma_2]]$$

For convenience the transformation functions defined above are combined into a single function tf.

The transformation function tf is applied to expressions by the tf[τ ; μ] = an. This has for its domain the set of expressions and its range the set of canonical expressions. - transform

$$\left[\begin{array}{l} \text{op}[\tau] = . \rightarrow \text{trp}[\tau; \mu]; \quad \text{op}[\tau] = : \rightarrow \text{tre}[\tau; \mu]; \\ \text{op}[\tau] = , \rightarrow \text{trc}[\tau; \mu]; \quad \text{op}[\tau] = * \rightarrow \text{trm}[\tau]; \\ \text{op}[\tau] = + \rightarrow \text{trs}[\tau]; \quad \text{op}[\tau] = - \rightarrow \text{trd}[\tau]; \\ \text{op}[\tau] = | \rightarrow \text{trb}[\tau] \end{array} \right]$$

This function scans the left and right parts of the expression and then applies the function tf. This action is necessary since the transformations have been organized with the assumption that both the left and right parts of any expression to be transformed are to be scanned. Such an assumption is not theoretically necessary but it does improve greatly the efficiency of the process by avoiding multiple scanning of the same expression.

It may be also observed that function red is applied to the left and right parts of the expression and that scan uses the functions scanl and scanr to generate models for subexpressions.

This section completes the description of the collecting process.

7. USE OF THE TRANSFORMATION RULES.

8.1 The transformation function tf is applied to expressions by the function scan. This has for its domain the set $E \times N'$ and for its range the set of canonic expressions.

$\text{scan}[\tau; \mu] =$ - scan

$$\left[\begin{array}{l} \tau \in RVL \rightarrow \tau; \\ \text{tf}[\text{scan}[\text{red}[\text{lp}'[\tau]; \text{modl}[\tau; \mu]]; \text{modl}[\tau; \mu]] \text{op}[\tau] \\ \text{scan}[\text{red}[\text{rp}'[\tau; 0]; \text{modr}[\tau; \mu]]; \text{modr}[\tau; \mu]] \end{array} \right]$$

This function scans the left and right parts of the expression and then applies the function tf. This method is necessary since the transformations have been organised with the assumption that both the left and right parts of any expression to be transformed must be canonic. Such an assumption is not theoretically necessary but it does improve greatly the efficiency of the process by avoiding multiple scanning of the same expression.

It may be also observed that function red is applied to the left and right parts of the expression and that scan uses the functions modl and modr to generate moduli for scanning of the subexpressions.

This section completes the description of the collecting process.

8. COMPUTATIONAL TECHNIQUES.

8.1 Use of SLIP and PL1.

The computational techniques used belong to the fields of symbol manipulation and list processing. Background information is obtainable in a survey paper by Sammet [SAMMJ660] and a full description of the list processor is contained in Weizenbaum [WEIZJ630]. A version of Weizenbaum's symmetric list processor SLIP was implemented in the programming languages PL1 and Assembler.

The version has been called SYMPLIP to emphasise the differences. At the time that this work was done the list processing features of PL1 were not implemented. For this reason and because of the desire to minimise debugging time, the routines of SLIP were essentially transliterated into PL1 with deletion of only the most obviously obsolete features such as AND and LANORM. As well as this, the SLIP recursive features were deleted and where recursion was required as in LSTEQL, the recursive features of PL1 were used. The primitives were written in Assembler.

The major difficulty encountered in the implementation was due to the short word length of 32 bits in the S/360 and the fact that all addressing is with respect to bytes. Since the structure of SLIP would have been seriously affected by a decision to use three words per cell, a scheme was devised using a LAVS relocation constant, relative addressing within LAVS and storage of word addresses rather

than byte addresses. This approach retains the two word cell and permits the use of a 128 kilo-byte LAVS in a 256 kilo-byte store.

8.2 Representation of mathematical structures.

The terminal symbols (group expression generators or rational expression generators) are represented by two cell lists consisting of the header and a datum containing the symbol and type identification. Figure 1 shows the list representation for the group element a . The datum field is regarded as consisting of four characters. The first is unused. The second defines the type of datum.

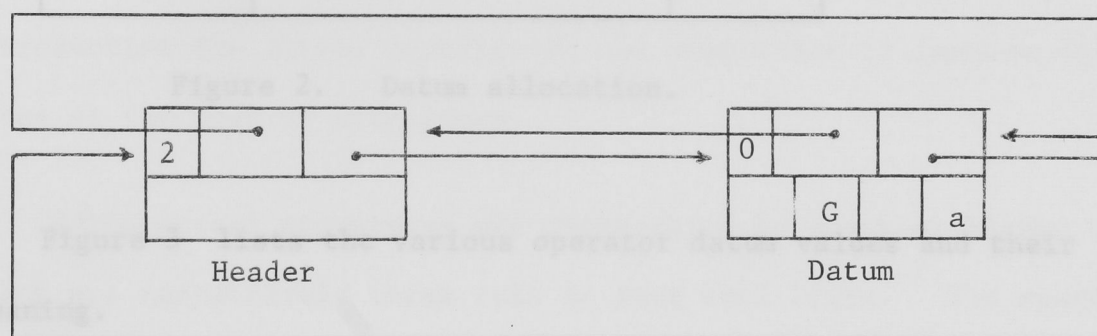


Figure 1. Representation of the generator a .

Figure 2 shows how the various character fields of the datum have been allocated.

Datum word character positions:

1	2	3	4
---	---	---	---

Character	Use	Value
1	Not used	
2	Operator designator	0
	Group expression generator	G
	Group identity	U
	Integer	I
	Rational expression generator	L
3	Special uses	
4	Datum value	

Figure 2. Datum allocation.

Figure 3 lists the various operator datum values and their meaning.

Datum values when character 2 = 0	Meaning
,	commutation
.	multiplication
:	exponentiation
'	group inversion
+	integer summation
/	integer division
	combination
-	binary integer negation
U	unary " "
*	integer multiplication

Figure 3.

The multiple comma operators are not represented in Figure 3 because there is only one operation of commutation and the use of multiple commas is merely a device for suppression of parentheses without loss of structure in the expression. Parentheses do not appear in the list structure simply because, while they help to define structure in a linear string of symbols, the tree or list contains the structure of an expression inherently.

It will be noted in the description above and in the following text that the formal description of earlier sections is not in one-to-one correspondence with the current description. As mentioned in the introduction the formal description was simplified to improve clarity, often at the cost of efficiency.

Mathematical structures are represented by unary or binary trees which are respectively three cell or four cell lists. The operation a, b has the tree and list representation as shown in Figure 4.

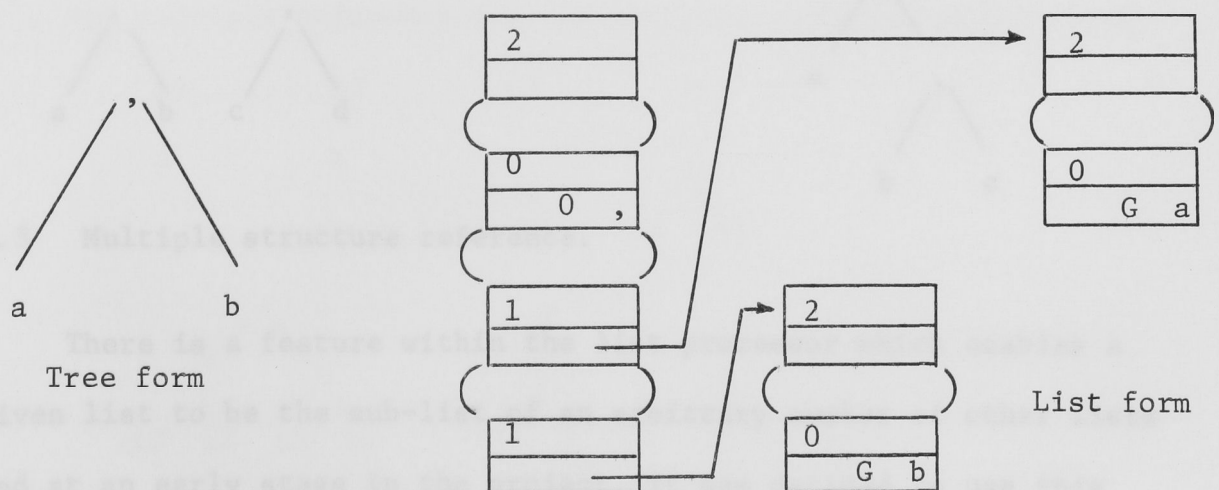
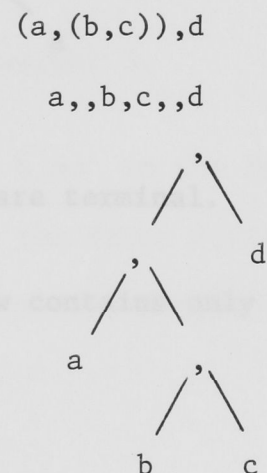
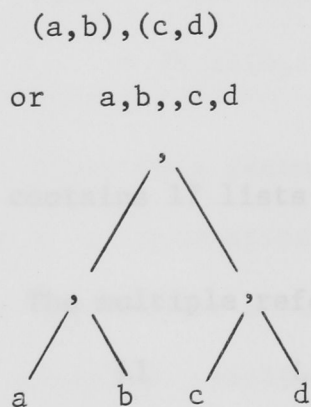


Figure 4.

This representation is general in that a or b themselves may be arbitrarily complex since the operation structure contains the names of the lists containing the operands.

While this choice of binary structure is perhaps the most obvious, it leads to prohibitive inefficiency when ordering products of commutators. After some reflection on this it is realised that since the associative law for multiplication holds, any binary tree representation of a product of more than two commutators will contain structure which is unrelated to the mathematical meaning of the expression. The means used to overcome this difficulty will be discussed later.

Consider the commutators $(a,b),(c,d)$ and $(a,(b,c)),d$. These have the following tree structures.



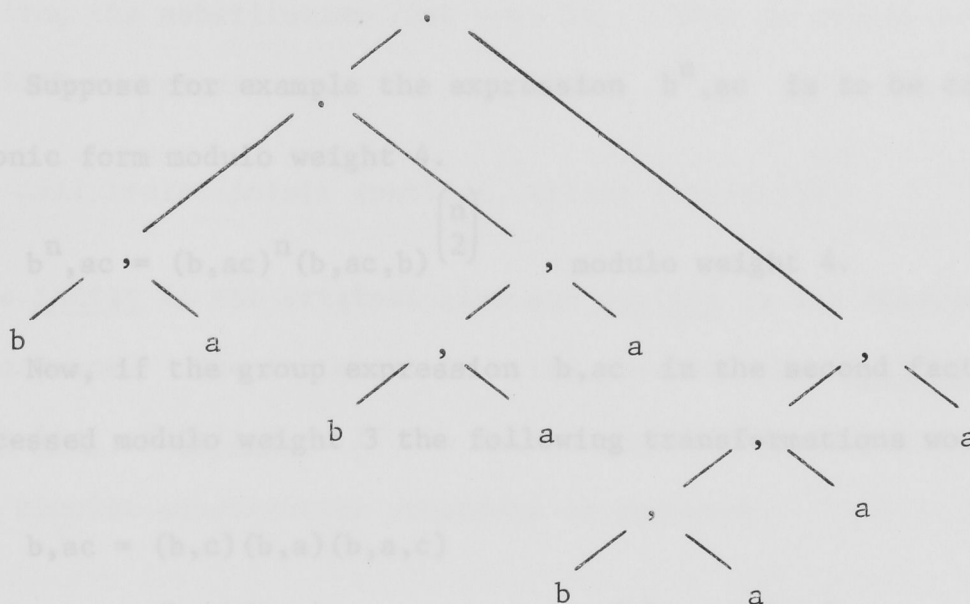
8.3 Multiple structure reference.

There is a feature within the list processor which enables a given list to be the sub-list of an arbitrary number of other lists and at an early stage in the project, it was decided to use this

feature to economise on the use of cells. Suppose, for example, that it is necessary to generate a list structure for the string

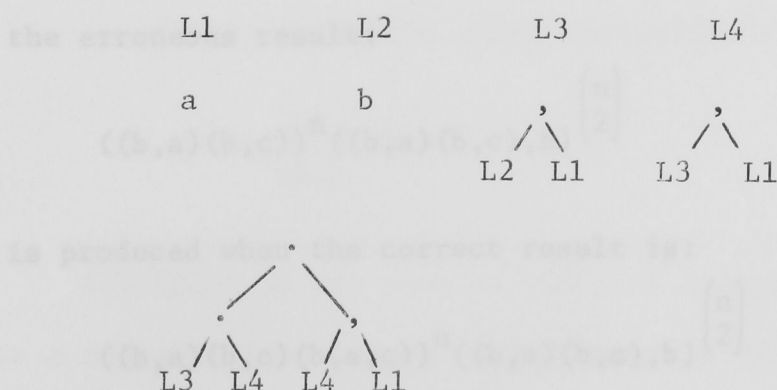
$(b,a)(b,a,a)(b,a,a,a)$

The single reference tree is as follows:



and contains 17 lists assuming a and b are terminal.

The multiple reference tree shown below contains only 7 lists



Multiple reference has a further advantage when converting expressions to canonic form since the conversion need be performed only once for any given subexpression provided the resultant canonic form is not dependent on context, i.e. the surrounding mathematical structure. This can, unfortunately, occur in many cases and special precautions are required.

Suppose for example the expression $b^{n,ac}$ is to be converted to canonic form modulo weight 4.

$$b^{n,ac} = (b,ac)^n (b,ac,b)^{\binom{n}{2}} \quad \text{modulo weight 4.} \quad \text{-- LEO}$$

Now, if the group expression b,ac in the second factor were processed modulo weight 3 the following transformations would occur:

$$\begin{aligned} b,ac &= (b,c)(b,a)(b,a,c) & \text{-- RPRDOUT} \\ &= (b,a)(b,c) & \text{modulo weight 3.} \end{aligned}$$

Since this second expression replaces b,ac in the factor b,ac,b it automatically replaces b,ac in the first factor when in fact this factor should have been processed modulo weight 4. Thus the erroneous result:

$$((b,a)(b,c))^n ((b,a)(b,c),b)^{\binom{n}{2}}$$

is produced when the correct result is:

$$((b,a)(b,c)(b,a,c))^n ((b,a)(b,c),b)^{\binom{n}{2}}$$

8.4 Structure transformation

The decision to permit multiple list references complicates the task of replacing one list structure by another because it is not easily known just what lists refer to a list to be substituted.

The difficulty is resolved by emptying the list to be replaced and building the substitution list onto it. This is simply performed by the SLIP statement:

```
call iralst(inlstr (outlist, mtlist (inlist)));
```

where inlist is the original list and outlist is the desired replacement. The statement is valid provided that outlist is not a sub-list of inlist. If the latter possibility exists a slightly more complex substitution procedure is required.

8.5 Transformation representation.

Each transformation is represented by a procedure usually of the form:

```
TRANS: PROCEDURE (STRUCTURE,MODULUS);
```

```
    . . .
```

```
END;
```

STRUCTURE is the name of the list structure which is known to be in a suitable form for the transformation TRANS. The parameter MODULUS provides where necessary the contextual information necessary to guide the transformation algorithm. On exit from the procedure the structure to which STRUCTURE points has been appropriately transformed.

Because of the problem associated with multiply referenced lists and also in the interest of speed of computation, it is normal practice to avoid the generation of non-canonic commutators within the transformation procedure.

When for example, the transformation rule

$ab,c = (a,c)(a,c,b)(b,c)$ is applied to ab,cd , modulo weight 4,

the first factor of the replacement expression generated is a,cd .

Since this is itself non-canonic, it is immediately subjected to a further transformation of the type

$a,bc = (a,c)(a,b)(a,b,c)$

This transforms a,cd to $(a,d)(a,c)(a,c,d)$. At this point the three commutators are still non-canonic since $a < d$ and $a < c$. After further manipulation:

$$a,cd = (c,a)^{-1}(d,a)^{-1}(c,a,d)^{-1}$$

The second term of the first transformation may now be processed.

It is obviously the first term commutated with b . That is

$$(c,a)^{-1}(d,a)^{-1}(c,a,d)^{-1},b$$

The first transformation is again called, effectively from within itself and eventually it realises that any commutator involving $(c,a,d)^{-1}$ and b will be of weight 4 and consequently equivalent to the identity. After further manipulation, this expression is reduced

to the canonic form $(c,a,b)^{-1}(d,a,b)^{-1}$. Note, of course, that only the main structure is replaced, the sub-structures, where applicable, remain attached to other super-structures.

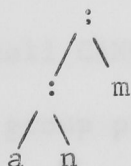
This is an efficient process because, on the assumption that both left and right parts of a structure are in canonic form, the depth of the non-canonic perturbation caused by the transformation is known in advance hence the canonic form can be restored more rapidly.

The manner in which this recursive behaviour is controlled will be explained below.

8.6 Pattern recognition or transformation selection.

Once a commutator expression is in the form of a list structure, the problem of pattern recognition and selection of transformation is trivial. The recognition process is centred in a procedure called CANONCL. Starting at the top of the tree this procedure makes a systematic attempt to classify the structure. The first basis for classification is the operator at the top of the tree. Observation of this operator permits immediate selection of one of the structure types described earlier in the text.

If, for example, CANONCL is presented with the structure $((a)^n)^m$ which in tree form is



the operator at the top of the tree is ':' and we know that the structure belongs to the set of exponential structures. The next symbol to be checked is the operator of the left part of the structure and if as in the example, this is also ':' then the transformation PRODEXP is called.

If a match does not occur, it is valid to assume that the structure submitted to CANONCL is, in fact, canonic and it is returned unaltered. The simplicity of this recognition process results from the generality of the transformation rules and the structural simplicity of the list representation.

8.7 Structure scanning.

When an arbitrarily non-basic commutator expression is submitted for conversion to canonic form a recursive scanning procedure traces a path down to the terminal symbols and then begins the ascent. As each binary operation is encountered on the ascent, the structure is submitted to CANONCL which returns the corresponding canonic expression. Because this process of forming basic expressions was begun at the deepest level, it is always true that the left and right parts of the structure submitted to CANONCL are canonic and this in turn leads to simplification of the transformation procedures.

There is an exceptional situation in which the scanning procedure does not necessarily call CANONCL. This is the case in which the current operator is a group product and the operator of the super-

structure is also a group product. The purpose of this is to delay the process of product ordering until the largest possible string of product terms is available.

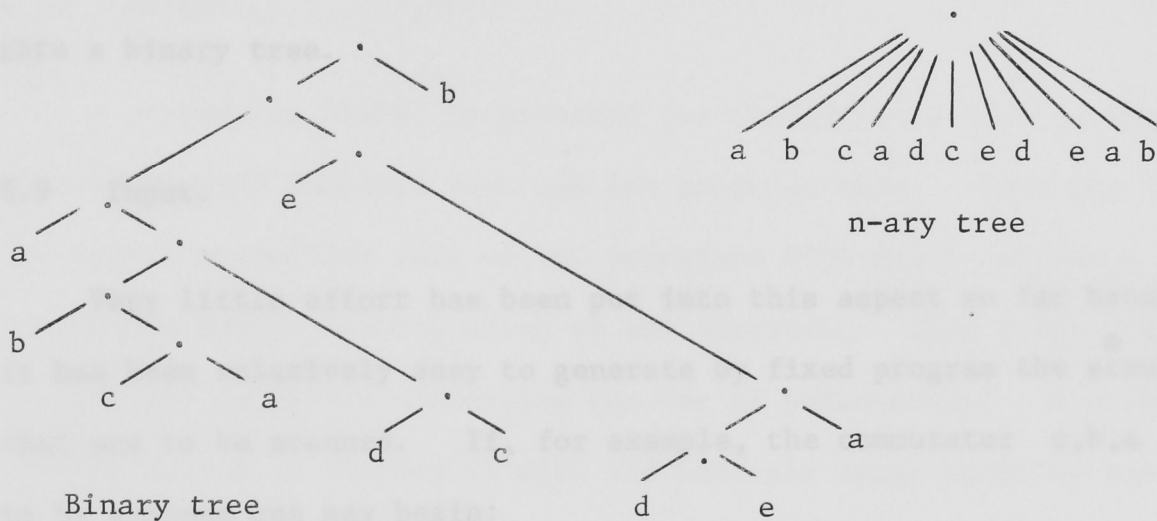
8.8 Ordering of products.

It will be recalled that the process of converting a commutator expression to canonic form consists of a sort of flotation process on the tree in which product operators are forced to assemble in a continuous block at the top of the structure followed by a thin layer of exponents (including inverses) followed by a block of commutation operators at the bottom. The upper layer of products will usually be very large and the time required for ordering them is some combinatorial power of the number of product operators in the string.

The original simple-minded approach to this product ordering problem was to treat product operations in exactly the same way as all other operations. In this way when the first binary product expression was detected by CANONCL the latter simply called the transformational procedure ORDPROD. This checked the order of right and left parts and if they were out of order it reordered them and introduced the additional commutator if necessary. While this is relatively simple in principle, it leads to a vast amount of back tracking and frequent repetitions of identical calculations.

To achieve tolerable efficiency in this vital aspect of the collecting process, it was necessary to assemble the largest possible

contiguous product string (i.e. no other imbedded operators) and transform it from an arbitrarily structured binary tree into a single simple n-ary tree. A diagram may help here:



Here the letters at the base of the tree may or may not be terminal but if they are non-terminal, they will be canonic and their uppermost operator is not a product.

The single node n-ary tree has the advantage that it can be scanned very rapidly and it is very easy to insert and delete components. When forming this simple list, various properties of the component expressions are evaluated and associated with the expressions so that they are evaluated only once during the ordering process. These quantities are the weight of the expression and its order number within weight. This latter number is so generated that for any two commutators a and b of equal weight if $a < b$ then the order number of a is less than the order number of b .

With this method, the speed of ordering was improved by a factor of about 20 in some cases and an even greater improvement would be observed for larger expressions.

When ordering is complete, the n -ary tree is transformed back into a binary tree.

8.9 Input.

Very little effort has been put into this aspect so far because it has been relatively easy to generate by fixed program the structures that are to be scanned. If, for example, the commutator c,b,a is to be scanned one may begin:

```
a = list(9); call nxlft(unspec('OG0a'),a);
b = list(9); call nxlft(unspec('OG0b'),b);
c = list(9); call nxlft(unspec('OG0c'),c);
l = lg(',',lg(',',c,b),a);
```

whereupon l contains the name of the commutator c,b,a .

Nevertheless, there are occasions where the input string is very long and the task of generating it in the above fashion is very laborious. For this reason, a procedure called STRNGET has been written for reading normal strings from cards and converting them into list structures. The routine is unsophisticated in that it assumes that the letters of the first half of the alphabet are group elements and those of the second half are literal representatives of integers.

Furthermore, although it does use multiple references to lists of the terminal symbols, it does not do this for non-terminal structures and the result can be a wasteful use of the list of available space.

8.10 Output.

A procedure, STRPN, is provided for transforming list structures to strings with standard notation and printing them. Both the input and output procedures call on the procedure PREC which returns a precedence number corresponding to any operator. This information is used for decisions concerning the use of parentheses. A further problem arises in connection with the multiple comma notation and it was found necessary to pre-scan the structure using the third character field of the structure datum word for recording the number of commas where applicable. This information is then readily accessible to the print routine which is a relatively simple recursive procedure.

While the datum word leaves adequate scope for a more extensive symbol range, input-output considerations lead to undesirable restrictions. Firstly, it is preferable not to use multi-character identifiers for letters since this makes expressions too long and therefore harder to read and also the convention of implicit multiplication must be dropped. Secondly, it has proved to be essential to have in excess of 30 or 40 symbols for group elements and exponents.

The ideal of course, is to be able to write expressions very similar to the following:

$$\left(\begin{matrix} \alpha_1 & \beta_1 & \gamma_1 \\ a_1 & b_1 & c_1 \end{matrix}, \begin{matrix} \alpha_2 & \beta_2 & \gamma_2 \\ a_2 & b_2 & c_2 \end{matrix} \right)$$

and graphic display with electrostatic printer will be necessary before this difficulty is overcome.

An assortment of computations on expressions of the form

$$\begin{matrix} \alpha_1 & \beta_1 & \gamma_1 & \alpha_2 & \beta_2 & \gamma_2 \\ a_1 & b_1 & c_1 & a_2 & b_2 & c_2 \end{matrix}$$

using literal arguments have been performed and the results checked by hand. In these cases the machine has proved to be only 3 or 10 times faster than the human calculator but far less subject to error. It is expected that the ratio of machine speed would increase with problem complexity because the machine has a more efficient technique for modifying expressions and in a relative sense the longer the expression the more effective the technique.

Much of the inefficiency of this automatic collecting process is due to the SLIP-FID combination. Even the simplest list processing operations are executed by a long chain of operations which in many cases could be replaced by a small number of in-line instructions, taking less space than the existing code.

However, there is still room for significant improvement of the algorithm. This is often realized when checking through the

9. EFFECTIVENESS OF THE IMPLEMENTATION.

An encouraging degree of success has been achieved with the automatic collecting process. One of the first expressions to be collected was $(ab)^4$ modulo weight 7 and this computation took 13 minutes of machine time. The computation, which was checked using the formulae of Krause [KRAUE640], has not been repeated by hand but it is expected that it would take many hours.

An assortment of computations on expressions of the form:

$$a^{\alpha_1} b^{\beta_1} (b, a)^{\gamma_1}, a^{\alpha_2} b^{\beta_2} (b, a)^{\gamma_2}$$

using literal exponents have been performed and the results checked by hand. In these cases the machine has proved to be only 5 or 10 times faster than the human calculator but far less subject to error. It is expected that the ratio of machine speed would increase with problem complexity because the machine has a more efficient technique for modifying expressions and in a relative sense the longer the expression the more effective the technique.

Much of the inefficiency of this automatic collecting process is due to the SLIP=PL1 combination. Even the simplest list processing operations are executed by a long chain of functions which in many cases could be replaced by a small number of in-line instructions, taking less space than the function call.

However, there is still room for significant improvement of the algorithm. This is often realised when checking through the

intermediate printout of a computation. In several cases, it can be observed that the machine achieves a result by a very devious path when the result is immediately obvious to the human calculator. In all such cases, the means for improving the machine algorithm have been quite simple but each improvement causes an increase in the size of the program and consequently less workspace.

An interesting example is that of evaluating

$$((b,a,b,,b,a)(b,a,b,b,b))^3 \quad \text{modulo weight 6}$$

To the human operator the result is immediate:

$$(b,a,b,,b,a)^3(b,a,b,b,b)^3$$

The computer procedures however, proceed as follows:

$$\begin{aligned} & ((b,a,b,,b,a)(b,a,b,b,b))^3 \quad \text{modulo weight 6} \\ = & (b,a,b,,b,a)(b,a,b,b,b)(b,a,b,,b,a)(b,a,b,b,b) \\ & (b,a,b,,b,a)(b,a,b,b,b) \\ = & (b,a,b,,b,a)^2(b,a,b,b,b)^2(b,a,b,,b,a)(b,a,b,b,b) \\ = & (b,a,b,,b,a)^3(b,a,b,b,b)^3 \end{aligned}$$

In cases where the value of the exponent is large, this is a very slow process.

There is another problem for the machine algorithm in its competition with the human algorithm. This is the problem of context dependent transformations which were discussed in Section 5. While it has been possible to build a very small number of these into the

machine algorithm the experienced human calculator has a large range of them at his disposal and can rapidly determine just when they should be applied.

For example, the commutator expression ab,cd modulo weight 3 is $(a,c)(a,d)(b,c)(b,d)$. This is a sort of macro transformation which is applicable when the expression is already of maximum weight.

This example may be partially generalised as shown in the following transformation:

$$\left(\prod_{i=1}^m a_i \right), \left(\prod_{j=1}^n b_j \right) = \prod_{i=1}^m \left(\prod_{j=1}^n (a_i, b_j) \right) \quad \text{modulo weight 3.}$$

Here the normally recursive methods for breaking down an expression are replaced by an iterative method, and instead of being able to select the required transformation on the basis of a couple of levels of operators it is necessary to consider an arbitrary number of levels.

The range of such transformations is limited only by the experience and memory of the human calculator. Machine algorithms of this type would in some of the calculations performed give rise to speed increases of 10 or more.

In this approach lies the greatest promise for an efficient algorithm. But the approach must be pursued fundamentally and systematically rather than heuristically. The transformations used in this text were originally chosen as the smallest and simplest set amenable to inductive proof and devoid of considerations of algorithmic efficiency. A completely new set of transformations should be developed with the latter aim in mind.

10. TERMS, SYMBOLS AND NOTATIONS USED IN THE TEXT.

10.1 Commonly used symbols and notations.

\wedge	propositional connective <u>and</u>
\vee	propositional connective <u>or</u>
\sim	propositional connective <u>not</u>
\Rightarrow	implication (note different usage in Section 10.4)
\Leftrightarrow	equivalence
\exists	there exists
$\{x P(x)\}$	the set of elements x with property P
$A \cup B$	the <u>union</u> of sets A and B
$A \cap B$	the <u>intersection</u> of sets A and B
$\bigcup_{i=m}^n A_i$	the union of a finite family of sets
$a \in A$	a is a member of the set A
$a \notin A$	a is not a member of the set A
$A \subseteq B$	A is a subset of B
$\binom{n}{r}$	$= n!/((n-r)!r!)$
x, y	the element $x^{-1} y^{-1} xy$ and known as a commutator
∞	infinity
Π	product function
function	a function is a triple (D, R, f) where D and R are non-empty sets and to each element x in D there corresponds an element denoted $f(x)$ in R . D is called the domain and R the range of the function.
predicate	a function whose range is the set $\{T, F\}$

10.2 Symbols and notations introduced in the text.

Symbol	Explanation	Page
Λ	the null word	21
[left meta-bracket	33, 89
]	right meta-bracket	33, 89
ϕ^{ℓ}	the function $lr[\phi]$	22
ϕ^r	the function $rs[\phi]$	22
T	the truth value <u>true</u>	24, 34, 89
F	the truth value <u>false</u>	24, 34, 89
<	left diamond bracket	8
>	right diamond bracket	8
::=	<u>becomes</u>	8
	separator for alternatives of a rule	8
<	integer or rational relation <u>less than</u>	34
>	integer or rational relation <u>greater than</u>	25, 34
=	integer or rational relation	34, 35
$\overset{A}{<}$	alphabetic order <u>less than</u>	35, 38
$\overset{A}{>}$	alphabetic order <u>greater than</u>	35
<•	relation in precedence grammar	18, 88
•>	relation in precedence grammar	18, 88
$\dot{=}$	relation in precedence grammar	18, 88
;	meta-delimiter	89
{<x>}	the set of all derivations of meta-variable <x>	34
a,...,k	group expression generators	3, 10
0,...,9	digits	3, 10
$\ell,...,z$	rational expression generators	3, 9

\vee	group identity	3, 10
$+$	rational summation operator	3, 10
$-$	rational difference operator	3, 10
$*$	rational product operator	3, 10
$ $	rational combination operator	3, 10
$/$	rational division operator	3, 10
$:$	group exponent operator	3, 10
$.$	group product operator	4, 16
$,$	group commutator operator	3, 4, 11
$,_{\alpha}$	subscripted commutator operator	4, 5, 11, 16
$($	left parenthesis	3, 10
$)$	right parenthesis	3, 10
\sim	digit concatenation operator	4, 16
variable with domain D		
B	{<rational expression generator>}	34, 39
E	$\{G \cup J\}$	17
\hat{E}	set of elements of E truncated on the left or right	20
E'	$\{G' \cup J'\}$	17
E^f	fully parenthesised subset of E	29
G	the internal grammar	17
G'	the input/output grammar	11, 17
G	{<group expression>}	5, 6, 17
G'	input/output equivalent of G	3, 11, 17
G^f	fully parenthesised subset of G	6, 29
J	{<rational expression>}	5, 17

J'	input/output equivalent of J	3, 10, 17
L	{<group expression generator>}	34, 37
N	{<integer>}	20, 22
N'	{<integer>} \cup $\{\infty\}$	34, 40
O	{::;/; ;+;-;.,;^;*}	34, 36
Q	{<rational number>}	34, 36
R	$\{B \cup Q\}$	34, 36
S	$\{R \cup L \cup O\}$ and symbols of the form $'_{\alpha}'$	5, 17
S'	$\{R \cup L \cup O\} \setminus \{.,;\wedge\}$	3, 4, 17
S'^*	the semigroup on S'	3, 17
S^*	the semigroup on S	5, 17
T	$\{T;F\}$	23, 34
γ	variable with domain G	3
λ	variable with domain J	3
ρ	variable with domain G^f	40
σ	variable with domain J^f	36
τ	variable with domain E^f	35
ϕ	variable with domain \hat{E} or E	20, 24
X	variable with domain O	37
ψ	variable with domain \hat{E} or E	21
μ	variable with domain N' or N	20, 22
α	variable with domain N	4, 11

10.3 Functions and terms introduced in the text.

<u>Function</u>	<u>Remark</u>	<u>Page</u>
<u>canonic</u>	canonic	43
<u>ce</u>	canonic exponent	40
<u>com</u>	commutator	42
<u>cp</u>	canonic product	38
<u>dif</u>	difference	54
<u>dlt</u>	dot less than	24
<u>f</u>	fully parenthesise	29
<u>invfac</u>	inverse of factorial	54
<u>ℓ</u>	length of word in	22
<u>ld</u>	leading part	41
<u>le</u>	left expression	37
<u>length</u>	length	36
<u>lo</u>	left operand	37
<u>lp</u>	left part	35
<u>lp'</u>	left part'	25
<u>lpr</u>	left product	39
<u>lr</u>	left remainder	21
<u>ls</u>	left symbol	21
<u>lt</u>	less than	42
<u>ltp</u>	less than product	39
<u>modl</u>	left modulus	46
<u>modr</u>	right modulus	46
<u>op</u>	operator	35
<u>op'</u>	operator'	25
<u>prd</u>	product	36

<u>prod</u>	product	54
<u>rat</u>	rational	36
<u>re</u>	right expression	37
<u>red</u>	reduce	47
<u>ro</u>	right operand	37
<u>rp</u>	right part	35
<u>rp'</u>	right part	25
<u>rpr</u>	right product	39
<u>rr</u>	right remainder	21
<u>rs</u>	right symbol	21
<u>scan</u>	scan expression	57
<u>skl</u>	skip level	23
<u>sum</u>	summation	54
<u>syp</u>	symbolic product	55
<u>tf</u>	transform	56
<u>tr</u>	trailing part	41
<u>trb</u>	transform combinatorial	55
<u>trc</u>	transform commutator	51
<u>trd</u>	transform difference	55
<u>tre</u>	transform exponent	53
<u>trm</u>	transform rational product	54
<u>trp</u>	transform group product	50
<u>trs</u>	transform sum	55
<u>wt</u>	weight	40

Term	Page
<c>	13
< c_α >	11
<co>	13
<commutator>	13
<commutator operator>	13
<digit>	10
<factor>	9
<g-factor>	10
<g-primary>	10
<group expression>	10
<group expression generator>	10
<g-term>	10
<gt>	13
<integer>	10
<primary>	10, 70, 86
<rational expression>	10
<rational expression generator>	9
<rational number>	10
<term>	10
alphabet	2, 8, 86
alternatives	9
backus-normal-form	8
canonic	33, 42

canonic form	1, 5
canonic group expression	40
collecting process	1
commutator	2
computation	30
concatenation	2
context-free phrase structured language	5, 86
formal language	2, 8
generators	2
grammar	3, 9
group expressions	3
level of parenthesisation	22
modulus	7, 44
meta-connective	9
meta-variable	8
non-terminal	22, 70, 86
null word	21
parsing	5
partitioning function	20, 21
phrase	2
precedence table	16, 18
rational expressions	3, 8, 9
rational numbers	3
recursive conditional expressions	6, 89
sentence	2
string	2
subexpression	45, 64

submodulus	our free grammars and languages.	46
subword		24
alphabet:	a finite non empty set	
terminal		59, 86
terminate		30
terminate finitely		30
transformation		48
weight		2, 40
word		2, 8

10.4 Context free grammars and languages.

alphabet:	a finite non empty set
sentence on	
an alphabet:	a finite sequence of elements of the alphabet
word:	a sentence
string:	a sentence
Λ	the sentence of length zero
S^*	the set of all words over (or on) an alphabet S
context-free grammar:	a 4-tuple $G = (V, S, P, s)$ where
	V is an alphabet
	$S \subseteq V$ is an alphabet
	P is a finite set of ordered pairs
	(u, v) with $u \in \{V-S\} - \{\Lambda\}$ and
	$v \in V^*$
	$s \in V-S$

The alphabet S is the set of symbols which will appear in words or sentences of the language (to be defined). These symbols are sometimes called terminal symbols.

The alphabet $V-S$ is a set of symbols which are used in defining the grammar of the language. These are sometimes called non-terminal symbols or meta-symbols.

P is a set of rules for transforming words on the meta-symbols into words on the full alphabet. Instead of being written in the form (u, v) these rules may be written $u \rightarrow v$.

s is the distinguished meta-symbol and usually will stand for 'a sentence' or 'a program'.

$w \Rightarrow y$ or w generates y for $w, y \in V^*$ if there exist z_1, z_2, u and v and $w = z_1 u z_2$ and $y = z_1 v z_2$ and $(u, v) \in P$.

$w \xRightarrow{*} y$ if either $w = y$ or there exist w_0, w_1, \dots, w_r such that $w_0 = w$, $w_r = y$ and $w_i \Rightarrow w_{i+1}$.
 $w \xRightarrow{*} y$ is called a generation or derivation.

context free language:

If G is a context free grammar then the subset $L(G)$ of S^* defined by $L(G) = \{x \in S^* \mid s \xRightarrow{*} x\}$ is called a context-free language

Informally this means that the language is made up of all those words in S^* which can be derived from the meta-variable s or the set of words which can be parsed in terms of the definition of a grammatical sentence.

10.5 Ambiguity of formal grammars.

A leftmost derivation is a derivation

$$w_0 \Rightarrow \dots \Rightarrow w_r$$

such that for each i , $w_i = u_i v_i y_i$ and $w_{i+1} = u_i z_i y_i$ with $(v_i, z_i) \in P$ and $u_i \in S^*$.

A grammar G is ambiguous if there is some word in $L(G)$ which can be generated by two different leftmost derivations.

10.6 Operator grammars.

An operator grammar is a context-free grammar such that no member of P takes the form (U, xU_1U_2y) where $U, U_1, U_2 \in V-S$ and x, y are arbitrary strings.

10.7 Precedence grammars.

A precedence grammar is an operator grammar for which no more than one of the three relations $<\cdot, \doteq$ or $\cdot>$ hold between ordered pairs of terminal symbols. These relations are defined below.

$T_1 \doteq T_2$ if $T_1, T_2 \in S$ and there is a rule (U, xT_1T_2y) or $(U, xT_1U_1T_2y)$

where $U_1 \in V-S$.

$T_1 \cdot > T_2$ if there is a rule (U, xU_1T_2y) and a derivation

$U_1 \Rightarrow z$ where $U_1 \in V-S$ and T_1 is the rightmost terminal symbol of z .

$T_1 < \cdot T_2$ if there is a rule (U, xT_1U_1y) and a derivation $U_1 \Rightarrow z$

where $U_1 \in V-S$ and T_2 is the leftmost terminal symbol of z .

These relations are called precedence relations.

If a context-free grammar is ambiguous it is not a precedence grammar. This follows from the two definitions and the parsing algorithm whose uniqueness is dependent on the uniqueness of the precedence relations.

10.8 Conditional expressions.

A conditional expression has the form

$$[p_1 \rightarrow e_1; p_2 \rightarrow e_2; \dots; p_n \rightarrow e_n] \text{ or}$$

$$[p_1 \rightarrow e_1; p_2 \rightarrow e_2; \dots; p_n \rightarrow e_n; e_{n+1}]$$

where the p_i are propositional expressions or predicates which have the value T or F (standing for true or false), and the e_i are expressions.

The predicates may themselves be conditional expressions as also may the expressions. In the text the expressions are either group expressions, rational expressions, integer expressions or predicates.

The conditional expressions above have the value of the first e_i such that p_i is the leftmost predicate having the value T. In the case of the second conditional expression if all p_i have the value F then it takes the value e_{n+1} .

If there is an undefined p_i and no true p_j to the left of it then the expression is undefined.

If p_i is the leftmost true predicate and e_i is undefined then the expression is undefined.

If in the case of the first conditional expression there is no true predicate then the conditional expression is undefined.

These conditional expressions are used to define functions which are called conditional functions.

```

BASIC:  PROC(COM)          BIT(1)    RECURSIVE;
        /* RETURNS '1'B IFF COM IS BASIC */
        DCL COM FIXED BIN(31,0);
        DCL BASIC ENTRY(FIXED BIN(31,0))RETURNS(BIT(1));
        DCL (
            WT,BOT,CONT
        )ENTRY(FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));
        DCL (
            LT,EQ
        )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
            RETURNS(BIT(1));
        DCL (
            L,R
        )FIXED BIN(31,0);
        DCL (
            MADNTP
        )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
            RETURNS(FIXED BIN(31,0));
        DCL NTERM ENTRY(FIXED BIN(31,0))RETURNS(BIT(1));
        IF ~ NTERM(COM) THEN RETURN('1'B);
        L=CONT(MADNTP(COM,2)+1); R=BOT(COM);
        IF NTERM(L) THEN RETURN (BASIC(L) & BASIC(R) & LT(R,L)
                                & (LT(BOT(L),R) | EQ(BOT(L),R)));
        ELSE RETURN (BASIC(L) & BASIC(R) & LT(R,L));
        END BASIC;

```

```

CANCEL  3
CANCEL  4
CANCEL  5
BASIC   0
BASIC   1
BASIC   2
BASIC   3
BASIC   4
BASIC   5
BASIC   6
BASIC   7
BASIC   8
BASIC   9
BASIC  10
BASIC  11
BASIC  12
BASIC  13
BASIC  14
BASIC  15
BASIC  16
BASIC  17
BASIC  18
BASIC  19
BASIC  20
BASIC  21
BASIC  22
BASIC  23
BASIC  24

```

```

CANCEL:  PROC(STR);
        /* A'A OR AA'-> 1 */
        DCL (

```

```

CANCEL  0
CANCEL  1
CANCEL  2

```

A=ADV CAL,STR,A	CANCEL	3
IF A=0 THEN FIXED BIN(31,0);	CANCEL	4
DCL (CANCEL	5
NXTLFT,INLSTR	CANCEL	6
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	CANCEL	7
RETURNS(FIXED BIN(31,0));	CANCEL	8
DCL LIST ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	CANCEL	9
DCL (CANCEL	10
MTLIST	CANCEL	11
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	CANCEL	12
A=LIST(9);	CANCEL	13
CAL=NXTLFT(UNSPEC('0000'),A);	CANCEL	14
CALL IRALST(INLSTR(A,MTLIST(STR)));	CANCEL	15
END CANCEL;	CANCEL	16

ADVLEL: PROC(LR,A)	ADVLEL	0
FIXED BIN(31,0);		
DCL (LR,A)	ADVLEL	1
FIXED BIN(31,0);		
DCL REED ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVLEL	2
DCL ADVLL ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	ADVLEL	3
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	ADVLEL	4
A=ADVLL(LR,0,0);	ADVLEL	5
IF A=0 THEN RETURN(REED(LR));	ADVLEL	6
RETURN (A);	ADVLEL	7
END ADVLEL;	ADVLEL	8

ADVLER: PROC(LR,A)	ADVLER	0
FIXED BIN(31,0);		
DCL (LR,A)	ADVLER	1
FIXED BIN(31,0);		
DCL REED ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVLER	2
DCL ADVLR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	ADVLER	3
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	ADVLER	4

```

A=ADVLR(LR,0,0);
IF A=0 THEN RETURN(REED(LR));
RETURN (A);
END ADVLR;

```

```

ADVLR 5
ADVLR 6
ADVLR 7
ADVLR 8

```

```

ADVLL: PROC(LR,J,K)          FIXED BIN(31,0);
DCL(LR,J,K)                 FIXED BIN(31,0);
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL SETDIR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
                 FIXED BIN(31,0),FIXED BIN(31,0))
                 RETURNS(FIXED BIN(31,0));
DCL STRIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
                 RETURNS(FIXED BIN(31,0));
DCL(CLR,LK,CAND,LLV) FIXED BIN(31,0) STATIC;
DCL CAL FIXED BIN(31,0) STATIC;
CLR=CONT(LR);
L5: LK=LNKL(CONT(LNKL(CLR)));
CAND=CONT(LK);
CALL SETDIR(-1,LK,-1,CLR);
IF ID(CAND)=2 THEN GOTO L2;
IF ID(CAND)≠J & ID(CAND)≠K THEN GOTO L5;
LLV=0; GOTO L6;
L2: LLV=-1;
L6: CAL= STRIND(CLR,LR);
RETURN(LLV);
END ADVLL;

```

```

ADVLL 0
ADVLL 1
ADVLL 2
ADVLL 3
ADVLL 4
ADVLL 5
ADVLL 6
ADVLL 7
ADVLL 8
ADVLL 9
ADVLL 10
ADVLL 11
ADVLL 12
ADVLL 13
ADVLL 14
ADVLL 15
ADVLL 16
ADVLL 17
ADVLL 18
ADVLL 19
ADVLL 20
ADVLL 21
ADVLL 22

```

```

ADVNL: PROC(LR,A)          FIXED BIN(31,0);

```

```

ADVNL 0

```



```

DCL (LR,A)          FIXED BIN(31,0);
DCL REED  ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL ADVLL ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));
A=ADVLL(LR,1,1);
IF A=0 THEN RETURN(REED(LR));
RETURN (A);
END ADVLNL;

```

```

ADVLNL 1
ADVLNL 2
ADVLNL 3
ADVLNL 4
ADVLNL 5
ADVLNL 6
ADVLNL 7
ADVLNL 8

```

```

ADVLNR: PROC(LR,A)          FIXED BIN(31,0);
DCL (LR,A)          FIXED BIN(31,0);
DCL REED  ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL ADVLR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));
A=ADVLNR(LR,1,1);
IF A=0 THEN RETURN(REED(LR));
RETURN (A);
END ADVLNR;

```

```

ADVLNR 0
ADVLNR 1
ADVLNR 2
ADVLNR 3
ADVLNR 4
ADVLNR 5
ADVLNR 6
ADVLNR 7
ADVLNR 8

```

```

ADVLR:  PROC(LR,J,K)          FIXED BIN(31,0);
DCL(LR,J,K)          FIXED BIN(31,0);
DCL CONT  ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL ID    ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL LNKL  ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL LNKR  ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL SETDIR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
FIXED BIN(31,0),FIXED BIN(31,0))
RETURNS(FIXED BIN(31,0));
DCL STRIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
RETURNS(FIXED BIN(31,0));

```

```

ADVLR 0
ADVLR 1
ADVLR 2
ADVLR 3
ADVLR 4
ADVLR 5
ADVLR 6
ADVLR 7
ADVLR 8
ADVLR 9
ADVLR 10

```

```
DCL(CLR,LK,CAND,LRV) FIXED BIN(31,0) STATIC;
```

```
DCL CAL FIXED BIN(31,0) STATIC;
```

```
CLR=CONT(LR);
```

```
L5: LK=LNKR(CONT(LNKL(CLR)));
```

```
CAND=CONT(LK);
```

```
CALL SETDIR(-1,LK,-1,CLR);
```

```
IF ID(CAND)=2 THEN GOTO L2;
```

```
IF ID(CAND)≠J & ID(CAND)≠K THEN GOTO L5;
```

```
LRV=0; GOTO L6; FIXED BIN(31,0) RETURN(FIXED BIN(31,0));
```

```
L2: LRV=-1;
```

```
L6: CAL= STRIND(CLR,LR); FIXED BIN(31,0) RETURN(FIXED BIN(31,0));
```

```
RETURN(LRV);
```

```
END ADVLR; RETURN(FIXED CLR);
```

```
ADVLR 11
```

```
ADVLR 12
```

```
ADVLR 13
```

```
ADVLR 14
```

```
ADVLR 15
```

```
ADVLR 16
```

```
ADVLR 17
```

```
ADVLR 18
```

```
ADVLR 19
```

```
ADVLR 20
```

```
ADVLR 21
```

```
ADVLR 22
```

```
ADVLR 23
```

```
ADVLR: PROC(LR,A) FIXED BIN(31,0);
```

```
DCL (LR,A) FIXED BIN(31,0);
```

```
DCL REED ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
```

```
DCL ADVLR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
```

```
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0)); FIXED BIN(31,0);
```

```
A=ADVLR(LR,1,0);
```

```
IF A=0 THEN RETURN(REED(LR));
```

```
RETURN (A);
```

```
END ADVLR;
```

```
ADVLR 0
```

```
ADVLR 1
```

```
ADVLR 2
```

```
ADVLR 3
```

```
ADVLR 4
```

```
ADVLR 5
```

```
ADVLR 6
```

```
ADVLR 7
```

```
ADVLR 8
```

```
ADVLR: PROC(LR,A) FIXED BIN(31,0);
```

```
DCL (LR,A) FIXED BIN(31,0);
```

```
DCL REED ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
```

```
DCL ADVLR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
```

```
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));
```

```
A=ADVLR(LR,1,0);
```

```
ADVLR 0
```

```
ADVLR 1
```

```
ADVLR 2
```

```
ADVLR 3
```

```
ADVLR 4
```

```
ADVLR 5
```

```

IF A=0 THEN RETURN(REED(LR));
RETURN (A);
END ADVLWR;

```

```

ADVLWR 6
ADVLWR 7
ADVLWR 8

```

```

ADVSEL: PROC(LR,A) RETURNS FIXED BIN(31,0);
DCL (LR,A) FIXED BIN(31,0);
DCL REED ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL ADVSL ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));
A=ADVSL(LR,0,0);
IF A=0 THEN RETURN(REED(LR));
RETURN (A);
END ADVSEL;

```

```

ADVSEL 0
ADVSEL 1
ADVSEL 2
ADVSEL 3
ADVSEL 4
ADVSEL 5
ADVSEL 6
ADVSEL 7
ADVSEL 8

```

```

ADVSR: PROC(LR,A) RETURNS FIXED BIN(31,0);
DCL (LR,A) FIXED BIN(31,0);
DCL REED ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL ADVSR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));
A=ADVSR(LR,0,0);
IF A=0 THEN RETURN(REED(LR));
RETURN (A);
END ADVSR;

```

```

ADVSR 0
ADVSR 1
ADVSR 2
ADVSR 3
ADVSR 4
ADVSR 5
ADVSR 6
ADVSR 7
ADVSR 8

```

```

ADVSL: PROC(L,J,K) RETURNS FIXED BIN(31,0);
DCL (L,J,K) FIXED BIN(31,0);
DCL (R,LCP,M,ADVSA,CAND,LK) FIXED BIN(31,0) STATIC;
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));

```

```

ADVSL 0
ADVSL 1
ADVSL 2
ADVSL 3

```

DCL INHALT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSL	4
DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSL	5
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSL	6
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSL	7
DCL SETDIR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	ADVSL	8
FIXED BIN(31,0),FIXED BIN(31,0))	ADVSL	9
RETURNS(FIXED BIN(31,0));	ADVSL	10
DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	ADVSL	11
FIXED BIN(31,0),FIXED BIN(31,0))	ADVSL	12
RETURNS(FIXED BIN(31,0));	ADVSL	13
DCL STRIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	ADVSL	14
RETURNS(FIXED BIN(31,0));	ADVSL	15
DCL NUCELL ENTRY RETURNS(FIXED BIN(31,0));	ADVSL	16
DCL LCNTR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSL	17
DCL RCELL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSL	18
DCL CAL FIXED BIN(31,0) STATIC;	ADVSL	19
R=CONT(L); CAND=CONT(LNKL(R));	ADVSL	20
IF ID(CAND)=1 THEN GOTO L6;	ADVSL	21
L1: LCP=LNKL(CAND); CALL SETDIR(-1,LCP,-1,R);	ADVSL	22
CAND=CONT(LCP);	ADVSL	23
IF ID(CAND)=2 THEN GOTO L4;	ADVSL	24
IF ID(CAND)=J ID(CAND)=K THEN GOTO L8;	ADVSL	25
IF ID(CAND)=-1 THEN GOTO L1;	ADVSL	26
L6: M=NUCELL; CAL= STRIND(R,M); CAL= STRIND(CONT(L+1),M+1);	ADVSL	27
CALL SETIND(-1,LNKR(CONT(LNKL(R)+1)),LCNTR(L)+1,L+1);	ADVSL	28
CALL SETDIR(-1,-1,M,R);	ADVSL	29
CAND=CONT(INHALT(LNKL(R)+1));	ADVSL	30
GOTO L1;	ADVSL	31
L4: IF LCNTR(L)=-0 THEN GOTO L9;	ADVSL	32
ADVSA=-1; GOTO L12;	ADVSL	33
L9: LK=LNKR(R); R=CONT(LK); CAL= STRIND(CONT(LK+1),L+1);	ADVSL	34
CAND=CONT(LNKL(R)); CALL RCELL(LK);	ADVSL	35
GOTO L1;	ADVSL	36
L8: ADVSA=0;	ADVSL	37

L12: CAL= STRIND(R,L);	ADVSL	38
RETURN(ADVSA);	ADVSL	39
END ADVSL;	ADVSL	40

ADVSNL: PROC(LR,A)	ADVSNL	0
DCL (LR,A)	ADVSNL	1
DCL REED ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSNL	2
DCL ADVSL ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	ADVSNL	3
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	ADVSNL	4
A=ADVSL(LR,1,1);	ADVSNL	5
IF A=0 THEN RETURN(REED(LR));	ADVSNL	6
RETURN (A);	ADVSNL	7
END ADVSNL;	ADVSNL	8

ADVSNR: PROC(LR,A)	ADVSNR	0
DCL (LR,A)	ADVSNR	1
DCL REED ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSNR	2
DCL ADVSR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	ADVSNR	3
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	ADVSNR	4
A=ADVSR(LR,1,1);	ADVSNR	5
IF A=0 THEN RETURN(REED(LR));	ADVSNR	6
RETURN (A);	ADVSNR	7
END ADVSNR;	ADVSNR	8

ADVSR: PROC(L,J,K)	ADVSR	0
DCL (L,J,K)	ADVSR	1
DCL (R,LCP,M,ADVSA,CAND,LK) FIXED BIN(31,0) STATIC;	ADVSR	2
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSR	3

DCL INHALT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSR	4
DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSR	5
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSR	6
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSR	7
DCL SETDIR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	ADVSR	8
FIXED BIN(31,0),FIXED BIN(31,0))	ADVSR	9
RETURNS(FIXED BIN(31,0));	ADVSR	10
DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	ADVSR	11
FIXED BIN(31,0),FIXED BIN(31,0))	ADVSR	12
RETURNS(FIXED BIN(31,0));	ADVSR	13
DCL STRIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	ADVSR	14
RETURNS(FIXED BIN(31,0));	ADVSR	15
DCL NUCELL ENTRY RETURNS(FIXED BIN(31,0));	ADVSR	16
DCL RCELL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSR	17
DCL LCNTR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSR	18
DCL CAL FIXED BIN(31,0) STATIC;	ADVSR	19
R=CONT(L); CAND=CONT(LNKL(R));	ADVSR	20
IF ID(CAND)=1 THEN GOTO L6;	ADVSR	21
L1: LCP=LNKR(CAND); CALL SETDIR(-1,LCP,-1,R);	ADVSR	22
CAND=CONT(LCP);	ADVSR	23
IF ID(CAND)=2 THEN GOTO L4;	ADVSR	24
IF ID(CAND)=J ID(CAND)=K THEN GOTO L8;	ADVSR	25
IF ID(CAND)=-1 THEN GOTO L1;	ADVSR	26
L6: M=NUCELL; CAL= STRIND(R,M); CAL= STRIND(CONT(L+1),M+1);	ADVSR	27
CALL SETIND(-1,LNKR(CONT(LNKL(R)+1)),LCNTR(L)+1,L+1);	ADVSR	28
CALL SETDIR(-1,-1,M,R);	ADVSR	29
CAND=CONT(INHALT(LNKL(R)+1));	ADVSR	30
GOTO L1;	ADVSR	31
L4: IF LCNTR(L)=-0 THEN GOTO L9;	ADVSR	32
ADVSA=-1; GOTO L12;	ADVSR	33
L9: LK=LNKR(R); R=CONT(LK); CAL= STRIND(CONT(LK+1),L+1);	ADVSR	34
CAND=CONT(LNKL(R)); CALL RCELL(LK);	ADVSR	35
GOTO L1;	ADVSR	36
L8: ADVSA=0;	ADVSR	37

L12: CAL= STRIND(R,L);	ADVSR	38
RETURN(ADVSA);	ADVSR	39
END ADVSR;	ADVSR	40
ADVSWL: PROC(LR,A)	ADVSWL	0
FIXED BIN(31,0);		
DCL (LR,A)	ADVSWL	1
FIXED BIN(31,0);		
DCL REED ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSWL	2
DCL ADVSL ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	ADVSWL	3
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	ADVSWL	4
A=ADVSL(LR,1,0);	ADVSWL	5
IF A=0 THEN RETURN(REED(LR));	ADVSWL	6
RETURN (A);	ADVSWL	7
END ADVSWL;	ADVSWL	8
ADVSWR: PROC(LR,A)	ADVSWR	0
FIXED BIN(31,0);		
DCL (LR,A)	ADVSWR	1
FIXED BIN(31,0);		
DCL REED ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	ADVSWR	2
DCL ADVSR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	ADVSWR	3
FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	ADVSWR	4
A=ADVSR(LR,1,0);	ADVSWR	5
IF A=0 THEN RETURN(REED(LR));	ADVSWR	6
RETURN (A);	ADVSWR	7
END ADVSWR;	ADVSWR	8
BOT:	BOT	0
B: PROC(P)	BOT	1
FIXED BIN(31,0);		
DCL (P)	BOT	2
FIXED BIN(31,0);		
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	BOT	3

DCL LNKL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	BOT	4
DCL LOCT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	BOT	5
RETURN	(CONT(LNKL(CONT(LOCT(P))))+1));	BOT	6
END BOT;		BOT	7

DELETE:	PROC(K)	FIXED BIN(31,0);	DELETE	0
	DCL (K)	FIXED BIN(31,0);	DELETE	1
	DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	DELETE	2
	DCL ID	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	DELETE	3
	DCL LNKL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	DELETE	4
	DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	DELETE	5
	DCL SETIND	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	DELETE	6
		FIXED BIN(31,0),FIXED BIN(31,0))	DELETE	7
		RETURNS(FIXED BIN(31,0));	DELETE	8
	DCL RCELL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	DELETE	9
	DCL (LL,LR)	FIXED BIN(31,0) STATIC;	DELETE	10
	IF ID(CONT(K))=2 THEN DC;		DELETE	11
	PUT LIST (' AN ATTEMPT HAS BEEN MADE TO DELETE A ',		DELETE	12
	'HEADER _ ZERO HAS BEEN DELIVERED AND THE ',		DELETE	13
	'PROGRAM CONTINUED');		DELETE	14
	RETURN(0);		DELETE	15
	END;		DELETE	16
	LL=LNKL(CONT(K));		DELETE	17
	LR=LNKR(CONT(K));		DELETE	18
	CALL RCELL(K);		DELETE	19
	CALL SETIND(-1,-1,LR,LL);		DELETE	20
	CALL SETIND(-1,LL,-1,LR);		DELETE	21
	RETURN(CONT(K+1));		DELETE	22
	END DELETE;		DELETE	23

INITAS: PROC(M,N);

INITAS 0

INITAS:	DCL SETDIR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	INITAS	1
	FIXED BIN(31,0),FIXED BIN(31,0))	INITAS	2
	RETURNS(FIXED BIN(31,0));	INITAS	3
	DCL MADOV ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INITAS	4
	DCL(M(N),N)FIXED BIN(31,0);	INITAS	5
	DCL(I,J,K) FIXED BIN(31,0);	INITAS	6
	DCL AVSL FIXED BIN(31,0) STATIC EXT;	INITAS	7
	DCL RELOC FIXED BIN(31,0) STATIC EXT;	INITAS	8
	DCL LAVSIZ FIXED BIN(31,0) STATIC EXT;	INITAS	9
	LAVSIZ=N;	INITAS	10
	RELOC=MADOV(M(1))-4;	INITAS	11
	DO I=1 TO N; M(I)=0; END;	INITAS	12
	K=N-2;	INITAS	13
	DO I=1 TO K BY 2;	INITAS	14
	CALL SETDIR(-1,-1,I+2,M(I));	INITAS	15
	END;	INITAS	16
	CALL SETDIR(0,N-1,1,AVSL);	INITAS	17
	END INITAS;	INITAS	18
INITRD:	PROC(K) FIXED BIN(31,0);	INITRD	0
	DCL (K) FIXED BIN(31,0);	INITRD	1
	DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INITRD	2
	DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INITRD	3
	DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	INITRD	4
	FIXED BIN(31,0),FIXED BIN(31,0))	INITRD	5
	RETURNS(FIXED BIN(31,0));	INITRD	6
	CALL SETIND(-1,LNKL(CONT(K+1)),-1,K);	INITRD	7
	RETURN(K);	INITRD	8
	END INITRD;	INITRD	9
INL STL:		INL STL	0

INL :	PROC(M,N)	FIXED BIN(31,0);	INLSTL	1
	DCL (M,N)	FIXED BIN(31,0);	INLSTL	2
	DCL (L,ITOP,IBOT,IPRE)	FIXED BIN(31,0) STATIC;	INLSTL	3
	DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INLSTL	4
	DCL LNKL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INLSTL	5
	DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INLSTL	6
	DCL SETIND	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	INLSTL	7
		FIXED BIN(31,0),FIXED BIN(31,0))	INLSTL	8
		RETURNS(FIXED BIN(31,0));	INLSTL	9
	DCL LOCT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INLSTL	10
	L=LOCT(M);		INLSTL	11
	ITOP=LNKR(CONT(L));		INLSTL	12
	IBOT=LNKL(CONT(L));		INLSTL	13
	CALL SETIND(-1,L,L,L);		INLSTL	14
	IPRE=LNKL(CONT(N));		INLSTL	15
	CALL SETIND(-1,IBOT,-1,N);		INLSTL	16
	CALL SETIND(-1,-1,ITOP,IPRE);		INLSTL	17
	CALL SETIND(-1,IPRE,-1,ITOP);		INLSTL	18
	CALL SETIND(-1,-1,N,IBOT);		INLSTL	19
	RETURN(L);		INLSTL	20
	END INLSTL;		INLSTL	21

INL STR :			INLSTR	0
INR :	PROC(M,N)	FIXED BIN(31,0);	INLSTR	1
	DCL (M,N)	FIXED BIN(31,0);	INLSTR	2
	DCL (L,ITOP,IBOT,ISLC)	FIXED BIN(31,0) STATIC;	INLSTR	3
	DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INLSTR	4
	DCL LNKL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INLSTR	5
	DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INLSTR	6
	DCL SETIND	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	INLSTR	7
		FIXED BIN(31,0),FIXED BIN(31,0))	INLSTR	8
		RETURNS(FIXED BIN(31,0));	INLSTR	9

DCL LOCT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INLSTR 10
L=LOCT(M);	INLSTR 11
ITOP=LNKR(CONT(L));	INLSTR 12
IBOT=LNKL(CONT(L));	INLSTR 13
CALL SETIND(-1,L,L,L);	INLSTR 14
ISUC=LNKR(CONT(N));	INLSTR 15
CALL SETIND(-1,-1,ITOP,N);	INLSTR 16
CALL SETIND(-1,IBOT,-1,ISUC);	INLSTR 17
CALL SETIND(-1,N,-1,ITOP);	INLSTR 18
CALL SETIND(-1,-1,ISUC,IBOT);	INLSTR 19
RETURN(L);	INLSTR 20
END INLSTR;	INLSTR 21

IRALST: PROC(P) FIXED BIN(31,0);	IRALST 0
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IRALST 1
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IRALST 2
DCL RCELL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IRALST 3
DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	IRALST 4
FIXED BIN(31,0),FIXED BIN(31,0))	IRALST 5
RETURNS(FIXED BIN(31,0));	IRALST 6
DCL LCNTR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IRALST 7
DCL LOCT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IRALST 8
DCL MTLIST ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IRALST 9
DCL (L,N,IRALSA) FIXED BIN(31,0) STATIC;	IRALST 10
DCL P FIXED BIN(31,0);	IRALST 11
DCL CAL FIXED BIN(31,0) STATIC;	IRALST 12
/* THE LIST P IS ERASED. IF THIS LIST IS THE	IRALST 13
SUBLIST OF ANOTHER LIST THE REFERENCE COUNTER	IRALST 14
IS COUNTED DOWNAND THE LIST IS NOT ERASED.	IRALST 15
THE VALUE OF THE FUNCTION IS THE NUMBER OF LISTS	IRALST 16
OF WHICH P IS A SUBLIST. IF THE VALUE IS ZERO	IRALST 17
THE LIST HAS BEEN ERASED */	IRALST 18

L=LUOT(P);	IRALST	19
CALL SETIND(-1,-1,LCNTR(L)-1,L+1);	IRALST	20
IRALSA=LCNTR(L);	IRALST	21
IF IRALSA \neq 0 THEN RETURN(IRALSA);	IRALST	22
CAL= MLIST(P);	IRALST	23
N=LNKL(CONT(L+1));	IRALST	24
IF N \neq 0 THEN DO;	IRALST	25
CALL SETIND(1,-1,-1,L);	IRALST	26
CALL SETIND(-1,N,N,L+1); END;	IRALST	27
CALL RCELL(L);	IRALST	28
RETURN(IRALSA);	IRALST	29
END IRALST;	IRALST	30

IRARDR: PROC(K) FIXED BIN(31,0);	IRARDR	0
DCL (K) FIXED BIN(31,0);	IRARDR	1
DCL (M,N) FIXED BIN(31,0) STATIC;	IRARDR	2
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IRARDR	3
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IRARDR	4
DCL LCNTR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IRARDR	5
DCL RCELL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IRARDR	6
M=K;	IRARDR	7
L3: N=LNKR(CONT(M)); CALL RCELL(M);	IRARDR	8
IF N=0 THEN RETURN(LCNTR(K));	IRARDR	9
M=N; GOTO L3;	IRARDR	10
END IRARDR;	IRARDR	11

LAVSOUT: PROC(LAVS,M);	LAVSOUT	0
DCL AVSL FIXED BIN(31,0) STATIC EXT;	LAVSOUT	1
DCL(LAVS(M),M)FIXED BIN(31,0);	LAVSOUT	2
DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LAVSOUT	3

DCL LNKR ENTRY(FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	LAVSOUT	4
DCL LNKL ENTRY(FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	LAVSOUT	5
DCL MADDV ENTRY(FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	LAVSOUT	6
DCL I FIXED BIN(31,0);	LAVSOUT	7
DCL RELOC FIXED BIN(31,0) STATIC EXT;	LAVSOUT	8
PUT SKIP(4);	LAVSOUT	9
PUT EDIT('AVSL')(A);	LAVSOUT	10
PUT EDIT(ID(AVSL),LNKL(AVSL),LNKR(AVSL))(F(10));	LAVSOUT	11
PUT SKIP(4);	LAVSOUT	12
PUT EDIT(' ADDR ID LNKL LNKR '	LAVSOUT	13
ID LNKL LNKR DATUM')(A);	LAVSOUT	14
PUT SKIP(4);	LAVSOUT	15
PUT EDIT((I,ID(LAVS(I)),LNKL(LAVS(I)),	LAVSOUT	16
LNKR(LAVS(I)),ID(LAVS(I+1)),LNKL(LAVS(I+1)),	LAVSOUT	17
LNKR(LAVS(I+1)),UNSPEC(LAVS(I+1))	LAVSOUT	18
DO I=1 TO M BY 2))	LAVSOUT	19
(4 F(8),X(10),3 F(8),X(5),B(32),SKIP);	LAVSOUT	20
END LAVSOUT;	LAVSOUT	21

LCNTR: PROC(K) FIXED BIN(31,0);	LCNTR	0
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LCNTR	1
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LCNTR	2
DCL K FIXED BIN(31,0);	LCNTR	3
RETURN(LNKR(CONT(K+1)));	LCNTR	4
END LCNTR;	LCNTR	5

LIST: PROC(K) FIXED BIN(31,0);	LIST	0
/* CREATES AN EMPTY LIST AND LEAVES ITS NAME AS ITS VALUE	LIST	1
AND IN THE CELL K WHICH THEREBY BECOMES AN ALIAS FOR THE LIS	LIST	2
*/	LIST	3

DCL NUCELL ENTRY RETURNS(FIXED BIN(31,0));	LIST	4
DCL SETDIR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	LIST	5
FIXED BIN(31,0),FIXED BIN(31,0))	LIST	6
RETURNS(FIXED BIN(31,0));	LIST	7
DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	LIST	8
FIXED BIN(31,0),FIXED BIN(31,0))	LIST	9
RETURNS(FIXED BIN(31,0));	LIST	10
DCL LISA FIXED BIN(31,0) STATIC;	LIST	11
DCL K FIXED BIN(31,0);	LIST	12
LISA=NUCELL;	LIST	13
CALL SETDIR(0,LISA,LISA,LISA);	LIST	14
CALL SETIND(2,LISA,LISA,LISA);	LIST	15
IF K=9 THEN DO; CALL SETIND(-1,-1,1,LISA+1); K=LISA; END;	LIST	16
RETURN(LISA);	LIST	17
END LIST;	LIST	18

LISTMT: PROC(P) FIXED BIN(31,0);	LISTMT	0
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LISTMT	1
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LISTMT	2
DCL LOCT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LISTMT	3
DCL P FIXED BIN(31,0);	LISTMT	4
DCL L FIXED BIN(31,0) STATIC;	LISTMT	5
L=CONT (LOCT(P));	LISTMT	6
IF L=CONT(LNKR(L)) THEN RETURN(0);	LISTMT	7
RETURN(-1);	LISTMT	8
END LISTMT;	LISTMT	9

LOCT: PROC(K) FIXED BIN(31,0);	LOCT	0
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LOCT	1
DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LOCT	2

DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LOCT	3
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LOCT	4
DCL NAMTST ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LOCT	5
DCL K FIXED BIN(31,0);	LOCT	6
IF NAMTST(K)=0 THEN RETURN(K);	LOCT	7
PUT SKIP(2) EDIT ('LIST WAS REQUIRED AS AN OPERAND BUT WAS',	LOCT	8
' NOT FOUND - PROGRAM REGRETFULLY TERMINATED')	LOCT	9
(A);	LOCT	10
PUT SKIP(2) LIST('LNKR=',LNKR(K),'LNKL=',LNKL(K));	LOCT	11
SIGNAL CONDITION(ALPHA);	LOCT	12
END LOCT;	LOCT	13

LOFRDR: PROC(K) FIXED BIN(31,0);	LOFRDR	0
DCL (K) FIXED BIN(31,0);	LOFRDR	1
DCL L FIXED BIN(31,0) STATIC;	LOFRDR	2
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LOFRDR	3
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LOFRDR	4
DCL SETDIR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	LOFRDR	5
FIXED BIN(31,0),FIXED BIN(31,0))	LOFRDR	6
RETURNS(FIXED BIN(31,0));	LOFRDR	7
L=LNKL(CONT(K+1));	LOFRDR	8
CALL SETDIR(0,L,L,L);	LOFRDR	9
RETURN(L);	LOFRDR	10
END LOFRDR;	LOFRDR	11

LPNTR: PROC(K) FIXED BIN(31,0);	LPNTR	0
DCL(K) FIXED BIN(31,0);	LPNTR	1
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LPNTR	2
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LPNTR	3
RETURN(LNKL(CONT(K)));	LPNTR	4

END LPNTR;

LPNTR 5

```

LPURGE: PRDC(LST)      FIXED BIN(31,0);
DCL (LST)              FIXED BIN(31,0);
DCL (X,LP,K,J,L)      FIXED BIN(31,0) STATIC;
  DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
  DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
  DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
    FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
  DCL LCNTR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
  DCL NAMTST ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
  DCL LRDRDV ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
  DCL IRARDR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
  DCL DELETE ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
  DCL LPNTR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
  DCL LSTPRO ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
  DCL ADVSWR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
  DCL ADVLWR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
  DCL CAL FIXED BIN(31,0) STATIC;
  LP=0;
L9: K=LRDRDV(LST);
L3: X=ADVSWR(K,J);
L6: IF J=0 THEN GOTO L2;
  IF NAMTST(X)=0 | LSTPRO(X,K)=0 THEN GOTO L3;
  L=LPNTR(K);
  CALL SETIND(-1,-1,LCNTR(X)-1,X+1);
  X=ADVLWR(K,L);
  CAL=DELETE(L);

```

```

LPURGE 0
LPURGE 1
LPURGE 2
LPURGE 3
LPURGE 4
LPURGE 5
LPURGE 6
LPURGE 7
LPURGE 8
LPURGE 9
LPURGE 10
LPURGE 11
LPURGE 12
LPURGE 13
LPURGE 14
LPURGE 15
LPURGE 16
LPURGE 17
LPURGE 18
LPURGE 19
LPURGE 20
LPURGE 21
LPURGE 22
LPURGE 23
LPURGE 24
LPURGE 25
LPURGE 26
LPURGE 27
LPURGE 28
LPURGE 29

```



```

LP=LP+1;
GOTO L6;
L2: IF LNKR(CONT(K)) $\neq$ 0 THEN DO;
      CAL= IRARDR(K);
      GOTO L9;
      END;
      CAL= IRARDR(K);
      RETURN (LP);
      END LPURGE;

```

```

LPURGE 30
LPURGE 31
LPURGE 32
LPURGE 33
LPURGE 34
LPURGE 35
LPURGE 36
LPURGE 37
LPURGE 38

```

```

LRDRCP: PROC(K)   FIXED BIN(31,0);
      DCL (K)   FIXED BIN(31,0);
      DCL (LR,NEW,NDW)   FIXED BIN(31,0) STATIC;
      DCL CONT   ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
      DCL LNKR   ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
      DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
                        FIXED BIN(31,0),FIXED BIN(31,0))
                        RETURNS(FIXED BIN(31,0));
      DCL STRIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
                        RETURNS(FIXED BIN(31,0));
      DCL NUCELL ENTRY RETURNS(FIXED BIN(31,0));
      DCL NEW FIXED BIN(31,0);
      DCL CAL FIXED BIN(31,0) STATIC;
      LR=NUCELL;
      NEW=LR;
      NDW=K;
L3: CAL= STRIND(CONT(NDW),NEW);
      CAL= STRIND(CONT(NDW+1),NEW+1);
      NDW=LNKR(CONT(NDW));
      IF NDW=0 THEN RETURN(LR);
      NEW=NUCELL;
      CALL SETIND(-1,-1,NEW,NEW);

```

```

LRDRCP 0
LRDRCP 1
LRDRCP 2
LRDRCP 3
LRDRCP 4
LRDRCP 5
LRDRCP 6
LRDRCP 7
LRDRCP 8
LRDRCP 9
LRDRCP 10
LRDRCP 11
LRDRCP 12
LRDRCP 13
LRDRCP 14
LRDRCP 15
LRDRCP 16
LRDRCP 17
LRDRCP 18
LRDRCP 19
LRDRCP 20
LRDRCP 21

```

TOP: NEWR=NEW; GOTO LB;	LRDRCP	22
END LRDRCP;	LRDRCP	23
LRDRDV: PROC(P) FIXED BIN(31,0);	LRDRDV	0
DCL (P) FIXED BIN(31,0);	LRDRDV	1
DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	LRDRDV	2
FIXED BIN(31,0),FIXED BIN(31,0))	LRDRDV	3
RETURNS(FIXED BIN(31,0));	LRDRDV	4
DCL LOCT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LRDRDV	5
DCL NUCELL ENTRY RETURNS(FIXED BIN(31,0));	LRDRDV	6
DCL LRDRDA FIXED BIN(31,0) STATIC;	LRDRDV	7
LRDRDA=NUCELL;	LRDRDV	8
CALL SETIND(3,LOCT(P),0,LRDRDA);	LRDRDV	9
CALL SETIND(0,P,0,LRDRDA+1);	LRDRDV	10
RETURN(LRDRDA);	LRDRDV	11
END LRDRDV;	LRDRDV	12
LSSCPY: PROC(LA) FIXED BIN(31,0) RECURSIVE;	LSSCPY	0
DCL (LA) FIXED BIN(31,0);	LSSCPY	1
DCL (LB,LR,LI,LRA,LRB,IDA,IDB) FIXED BIN(31,0);	LSSCPY	2
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSSCPY	3
DCL LIST ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSSCPY	4
DCL CAL FIXED BIN(31,0) STATIC;	LSSCPY	5
DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSSCPY	6
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSSCPY	7
DCL LSSCPY ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSSCPY	8
DCL NEWBOT ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	LSSCPY	9
RETURNS(FIXED BIN(31,0));	LSSCPY	10
DCL LOCT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSSCPY	11
LB=LIST(LB); LR=LNKR(CONT(LA));	LSSCPY	12

TOP: LI=ID(CONT(LR));	LSSCPY	13
IF LI=0 THEN DO; CAL= NEWBOT(CONT(LR+1),LB);	LSSCPY	14
LR=LNKR(CONT(LR));	LSSCPY	15
GOTO TOP;	LSSCPY	16
IF LI=1 THEN DO; CAL=NEWBOT(LSSCPY(LCCT(CONT(LR+1))),LB);	LSSCPY	17
LR=LNKR(CONT(LR));	LSSCPY	18
GOTO TOP;	LSSCPY	19
END;	LSSCPY	20
IF LI=2 THEN RETURN(LB);	LSSCPY	21
PUT LIST('INVALID ID FIELD ENCOUNTERED IN LSSCPY'); STOP;	LSSCPY	22
END LSSCPY;	LSSCPY	23
	LSSCPY	24
LSTEQL: PROC(LA,LB) FIXED BIN(31,0) RECURSIVE;	LSTEQL	0
DCL (LA,LB) ENTRY(FIXED BIN(31,0));	LSTEQL	1
DCL (IDA,IDB,LRA,LRB) FIXED BIN(31,0);	LSTEQL	2
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSTEQL	3
DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSTEQL	4
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSTEQL	5
DCL LSTEQL ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	LSTEQL	6
RETURNS(FIXED BIN(31,0));	LSTEQL	7
LRA=LNKR(CONT(LA)); LRB= LNKR(CONT(LB));	LSTEQL	8
TOP: IDA=ID(CONT(LRA)); IDB=ID(CONT(LRB));	LSTEQL	9
IF IDA≠IDB THEN RETURN(-1);	LSTEQL	10
IF IDA=0 THEN IF CONT(LRA+1)=CONT(LRB+1)	LSTEQL	11
THEN DO LRA=LNKR(CONT(LRA));	LSTEQL	12
LRB=LNKR(CONT(LRB));	LSTEQL	13
GOTO TOP;	LSTEQL	14
END;	LSTEQL	15
ELSE RETURN(-1);	LSTEQL	16
IF IDA=1 THEN IF LSTEQL(CONT(LRA+1),CONT(LRB+1))=0	LSTEQL	17
THEN DO LRA=LNKR(CONT(LRA));	LSTEQL	18

RETURNS(FIXED LRB=LNKR(CONT(LRB)));	LSTEQL	19
DCL CAL FIXED BIN(31,0) GOTO TOP;	LSTEQL	20
END;	LSTEQL	21
IF LCNTR=0 THEN ELSE RETURN(-1);	LSTEQL	22
IF IDA=2 THEN RETURN(0);	LSTEQL	23
PUT LIST('INVALID ID FIELD ENCOUNTERED IN LSTEQL');	LSTEQL	24
STOP;	LSTEQL	25
END LSTEQL;	LSTEQL	26

LSTPRO: PROC(L,K)	FIXED BIN(31,0);	LSTPRO	0
DCL(L,K)	FIXED BIN(31,0);	LSTPRO	1
DCL (NEXT)	FIXED BIN(31,0) STATIC;	LSTPRO	2
DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSTPRO	3
DCL LNKL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSTPRO	4
DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LSTPRO	5
NEXT=K;		LSTPRO	6
L3: IF LNKL(CONT(NEXT+1))=LNKR(L) THEN RETURN (0);		LSTPRO	7
NEXT=LNKR(CONT(NEXT));		LSTPRO	8
IF NEXT=0 THEN GOTO L3;		LSTPRO	9
RETURN(-1);		LSTPRO	10
END LSTPRO;		LSTPRO	11

LVLRV1: PROC(K)	FIXED BIN(31,0);	LVLRV1	0
DCL (K)	FIXED BIN(31,0);	LVLRV1	1
DCL (LVL RVA,L)	FIXED BIN(31,0) STATIC;	LVLRV1	2
DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LVLRV1	3
DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LVLRV1	4
DCL LCNTR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LVLRV1	5
DCL RCELL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LVLRV1	6
DCL STRIND	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	LVLRV1	7

PROC(K) RETURNS(FIXED BIN(31,0));	LVLRV1	8
DCL CAL FIXED BIN(31,0) STATIC;	LVLRV1	9
LVLRVA=K;	LVLRV1	10
IF LCNTR(K)=0 THEN RETURN(LVLRVA);	LVLRV1	11
L=LNKR(CONT(LVLRVA));	LVLRV1	12
CAL= STRIND(CONT(L),LVLRVA);	LVLRV1	13
CAL= STRIND(CONT(L+1),LVLRVA+1);	LVLRV1	14
CALL RCELL(L);	LVLRV1	15
RETURN(LVLRVA);	LVLRV1	16
END LVLRV1;	LVLRV1	17
LVLVRT: PROC(K) FIXED BIN(31,0);	LVLVRT	0
DCL (K) FIXED BIN(31,0);	LVLVRT	1
DCL (LVLRVA,L) FIXED BIN(31,0) STATIC;	LVLVRT	2
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LVLVRT	3
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LVLVRT	4
DCL STRIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	LVLVRT	5
RETURNS(FIXED BIN(31,0));	LVLVRT	6
DCL LCNTR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LVLVRT	7
DCL RCELL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	LVLVRT	8
DCL CAL FIXED BIN(31,0) STATIC;	LVLVRT	9
LVLRVA=K;	LVLVRT	10
L1: IF LCNTR(K)=0 THEN RETURN(LVLRVA) ;	LVLVRT	11
L=LNKR(CONT(LVLRVA));	LVLVRT	12
CAL= STRIND(CONT(L),LVLRVA);	LVLVRT	13
CAL= STRIND(CONT(L+1),LVLRVA+1);	LVLVRT	14
CALL RCELL(L); GOTO L1;	LVLVRT	15
END LVLVRT;	LVLVRT	16
MADLEFT:	MADLEFT	0

ML: PROC(K) FIXED BIN(31,0);	MADLFT	1
DCL (K) FIXED BIN(31,0);	MADLFT	2
DCL M FIXED BIN(31,0) STATIC;	MADLFT	3
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADLFT	4
DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADLFT	5
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADLFT	6
DCL SETDIR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	MADLFT	7
FIXED BIN(31,0),FIXED BIN(31,0))	MADLFT	8
RETURNS(FIXED BIN(31,0));	MADLFT	9
M=LNKL(CONT(K));	MADLFT	10
IF ID(CONT(M))=2 THEN CALL SETDIR(0,M,M,M) ;	MADLFT	11
RETURN(M);	MADLFT	12
END MADLFT;	MADLFT	13

MADNBT: PROC(P,N)	FIXED BIN(31,0);	MADNBT	0
DCL (P,N)	FIXED BIN(31,0);	MADNBT	1
DCL L	FIXED BIN(31,0) STATIC;	MADNBT	2
DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADNBT	3
DCL ID	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADNBT	4
DCL LNKL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADNBT	5
DCL SETDIR	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	MADNBT	6
FIXED BIN(31,0),FIXED BIN(31,0))	MADNBT	7	
RETURNS(FIXED BIN(31,0));	MADNBT	8	
DCL LOCT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADNBT	9
DCL 1	FIXED BIN(31,0) STATIC;	MADNBT	10
L=LOCT(P);	MADNBT	11	
DO I=1 TO N ; L=LNKL(CONT(L)); END;	MADNBT	12	
IF ID(CONT(L))=2 THEN CALL SETDIR(0,L,L,L);	MADNBT	13	
RETURN (L);	MADNBT	14	
END MADNBT;	MADNBT	15	

MADNTP: PROC(P,N)	FIXED BIN(31,0);	MADNTP	0
-------------------	------------------	--------	---

DCL (P,N)	FIXED BIN(31,0);	MADNTP	1
DCL L	FIXED BIN(31,0) STATIC;	MADNTP	2
DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADNTP	3
DCL ID	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADNTP	4
DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADNTP	5
DCL SETDIR	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	MADNTP	6
	FIXED BIN(31,0),FIXED BIN(31,0))	MADNTP	7
	RETURNS(FIXED BIN(31,0));	MADNTP	8
DCL LOCT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADNTP	9
DCL I	FIXED BIN(31,0) STATIC;	MADNTP	10
	L=LOCT(P);	MADNTP	11
DO I=1 TO N;	L=LNKR(CONT(L)); END;	MADNTP	12
	IF ID(CONT(L))=2 THEN CALL SETDIR(0,L,L,L);	MADNTP	13
	RETURN (L);	MADNTP	14
	END MADNTP;	MADNTP	15

MADRGT:		MADRGT	0	
MR:	PROC(K)	FIXED BIN(31,0);	MADRGT	1
	DCL (K)	FIXED BIN(31,0);	MADRGT	2
	DCL M	FIXED BIN(31,0) STATIC;	MADRGT	3
	DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADRGT	4
	DCL ID	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADRGT	5
	DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	MADRGT	6
	DCL SETDIR	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	MADRGT	7
		FIXED BIN(31,0),FIXED BIN(31,0))	MADRGT	8
		RETURNS(FIXED BIN(31,0));	MADRGT	9
	M=LNKR(CONT(K));	MADRGT	10	
	IF ID(CONT(M))=2 THEN CALL SETDIR(0,M,M,M);	MADRGT	11	
	RETURN (M);	MADRGT	12	
	END MADRGT;	MADRGT	13	

MRKLSS:	PROC(M,LST)	FIXED BIN(31,0);	MRKLSS	0
---------	-------------	------------------	--------	---

MRKLSS	1
MRKLSS	2
MRKLSS	3
MRKLSS	4
MRKLSS	5
MRKLSS	6
MRKLSS	7
MRKLSS	8
MRKLSS	9
MRKLSS	10
MRKLSS	11
MRKLSS	12
MRKLSS	13
MRKLSS	14
MRKLSS	15
MRKLSS	16
MRKLSS	17
MRKLSS	18

MRKLST	0
MRKLST	1
MRKLST	2
MRKLST	3
MRKLST	4
MRKLST	5
MRKLST	6
MRKLST	7
MRKLST	8
MRKLST	9

MTLIST 0


```

/* RESTORES LIST P TO LAVS AND LEAVES
   THE NAME OF THE EMPTY LIST AS ITS
   VALUE */
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL SETDIR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
                 FIXED BIN(31,0),FIXED BIN(31,0))
            RETURNS(FIXED BIN(31,0));
DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
                 FIXED BIN(31,0),FIXED BIN(31,0))
            RETURNS(FIXED BIN(31,0));
DCL LISTMT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL LOCT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL(M,LR,LL) FIXED BIN(31,0) STATIC;
DCL P FIXED BIN(31,0);
DCL AVSL FIXED BIN(31,0) STATIC EXT;
M=LOCT(P);
IF LISTMT(P)=0 THEN RETURN(M);
LR=LNKR(CONT(M)); LL=LNKL(CONT(M)); CALL SETIND(-1,M,M,M);
CALL SETIND(-1,-1,LR,LNKL(AVSL));
CALL SETDIR(-1,LL,-1,AVSL);
L5: CALL SETIND(-1,-1,0,LNKL(AVSL));
RETURN(M);
END MTLIST;

```

```

MTLIST 1
MTLIST 2
MTLIST 3
MTLIST 4
MTLIST 5
MTLIST 6
MTLIST 7
MTLIST 8
MTLIST 9
MTLIST 10
MTLIST 11
MTLIST 12
MTLIST 13
MTLIST 14
MTLIST 15
MTLIST 16
MTLIST 17
MTLIST 18
MTLIST 19
MTLIST 20
MTLIST 21
MTLIST 22
MTLIST 23
MTLIST 24
MTLIST 25

```

```

NAMTST: PROC(K) FIXED BIN(31,0);
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL K FIXED BIN(31,0);

```

```

NAMTST 0
NAMTST 1
NAMTST 2
NAMTST 3
NAMTST 4
NAMTST 5

```

```

DCL LAVSIZ FIXED BIN(31,0) STATIC EXT;
  IF LNKR(K)>LAVSIZ-1 THEN RETURN(-1);
  IF LNKL(K)≠LNKR(K) THEN RETURN(-1);
  IF ID(CONT(K))≠2 THEN RETURN(-1);
  IF CONT(LNKR(CONT(LNKL(CONT(K))))≠CONT(K) THEN RETURN(-1);
RETURN(0);
END NAMTST;

```

```

NAMTST 6
NAMTST 7
NAMTST 8
NAMTST 9
NAMTST 10
NAMTST 11
NAMTST 12

```

NEWBOT:

NB:

```

PROC(P,Q)          FIXED BIN(31,0);
DCL (P,Q)          FIXED BIN(31,0);
DCL LOCT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL NXTLEFT ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
RETURN(NXTLEFT(P,LOCT(Q)));
END NEWBOT;

```

```

NEWBOT 0
NEWBOT 1
NEWBOT 2
NEWBOT 3
NEWBOT 4
NEWBOT 5
NEWBOT 6
NEWBOT 7

```

NEWTOP:

NT:

```

PROC(P,Q)          FIXED BIN(31,0);
DCL (P,Q)          FIXED BIN(31,0);
DCL LOCT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL NXTRGTT ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
RETURN(NXTRGTT(P,LOCT(Q)));
END NEWTOP;

```

```

NEWTOP 0
NEWTOP 1
NEWTOP 2
NEWTOP 3
NEWTOP 4
NEWTOP 5
NEWTOP 6
NEWTOP 7

```

NUCELL:

```

PROC          FIXED BIN(31,0);
/* TAKES A CELL FROM LAVS AND RETURNS MACHINE ADDRESS */

```

```

NUCELL 0
NUCELL 1

```

DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NUCELL	2
DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NUCELL	3
DCL ID	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NUCELL	4
DCL STRIND	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	NUCELL	5
	RETURNS(FIXED BIN(31,0));	NUCELL	6
DCL SETDIR	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	NUCELL	7
	FIXED BIN(31,0),FIXED BIN(31,0))	NUCELL	8
	RETURNS(FIXED BIN(31,0));	NUCELL	9
DCL IRALST	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NUCELL	10
DCL	M FIXED BIN(31,0) STATIC;	NUCELL	11
DCL AVSL	FIXED BIN(31,0) STATIC EXT;	NUCELL	12
DCL CAL	FIXED BIN(31,0) STATIC;	NUCELL	13
M=LNKR(AVSL);		NUCELL	14
IF M=C THEN DO;		NUCELL	15
	PUT EDIT (' LIST OF AVAILABLE SPACE',	NUCELL	16
	' EXHAUSTED - PROGRAM TERMINATED') (A);	NUCELL	17
SIGNAL CONDITION(ALPHA);		NUCELL	18
STOP;		NUCELL	19
END;		NUCELL	20
	ELSE IF ID(CONT(M))=1 THEN CAL= IRALST(CONT(M+1));	NUCELL	21
CALL SETDIR	(-1,-1,LNKR(CONT(M)),AVSL);	NUCELL	22
CAL=	STRIND(C,M); CAL= STRIND(C,M+1);	NUCELL	23
RETURN(M);		NUCELL	24
END NUCELL;		NUCELL	25

NULSTL:	PROC(LNKP,LNKH)	FIXED BIN(31,0);	NULSTL	0
	DCL (LNKP,LNKH)	FIXED BIN(31,0);	NULSTL	1
	DCL (NULSTA,LTOP,LSUC)	FIXED BIN(31,0) STATIC;	NULSTL	2
	DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NULSTL	3
	DCL ID	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NULSTL	4
	DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NULSTL	5
	DCL SETIND	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	NULSTL	6

CALL SETIND(-1,-1,NULSTA, FIXED BIN(31,0), FIXED BIN(31,0))	NULSTL	7
CALL SETIND(-1,NULSTL, RETURNS(FIXED BIN(31,0)));	NULSTL	8
DCL LIST ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NULSTL	9
NULSTA=LIST(9);	NULSTL	10
IF ID (CONT(LNKP))=2 THEN RETURN(NULSTA);	NULSTL	11
LTOP=LNKR(CONT(LNKH));	NULSTL	12
LSUC=LNKR(CONT(LNKP));	NULSTL	13
CALL SETIND(-1,-1,LSUC,LNKH);	NULSTL	14
CALL SETIND(-1,LNKH,-1,LSUC);	NULSTL	15
CALL SETIND(2,LNKP,LTOP,NULSTA);	NULSTL	16
CALL SETIND(-1,-1,NULSTA,LNKP);	NULSTL	17
CALL SETIND(-1,NULSTA,-1,LTOP);	NULSTL	18
RETURN(NULSTA);	NULSTL	19
END NULSTL;	NULSTL	20
NULSTR: PROC(LNKP,LNKH) FIXED BIN(31,0);	NULSTR	0
DCL (LNKP,LNKH) FIXED BIN(31,0);	NULSTR	1
DCL (NULSTA,LBOT,LPRE) FIXED BIN(31,0) STATIC;	NULSTR	2
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NULSTR	3
DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NULSTR	4
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NULSTR	5
DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	NULSTR	6
FIXED BIN(31,0),FIXED BIN(31,0));	NULSTR	7
RETURNS(FIXED BIN(31,0));	NULSTR	8
DCL LIST ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NULSTR	9
NULSTA=LIST(9);	NULSTR	10
IF ID(CONT(LNKP))=2 THEN RETURN(NULSTA);	NULSTR	11
LBOT=LNKL(CONT(LNKH));	NULSTR	12
LPRE=LNKL(CONT(LNKP));	NULSTR	13
CALL SETIND(-1,LPRE,-1,LNKH);	NULSTR	14
CALL SETIND(-1,-1,LNKH,LPRE);	NULSTR	15
CALL SETIND(2,LBOT,LNKP,NULSTA);	NULSTR	16

CALL SETIND(-1,-1,NULSTA,LBOT);	NULSTR	17
CALL SETIND(-1,NULSTA,-1,LNKP);	NULSTR	18
RETURN(NULSTA);	NULSTR	19
END NULSTR;	NULSTR	20

NXTLEFT:		NXTLEFT	0	
NXL:	PRDC(M,A)	FIXED BIN(31,0);	NXTLEFT	1
	DCL(M,A)	FIXED BIN(31,0);	NXTLEFT	2
	DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NXTLEFT	3
	DCL LNKL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NXTLEFT	4
	DCL SETIND	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	NXTLEFT	5
		FIXED BIN(31,0),FIXED BIN(31,0))	NXTLEFT	6
		RETURNS(FIXED BIN(31,0));	NXTLEFT	7
	DCL STRIND	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	NXTLEFT	8
		RETURNS(FIXED BIN(31,0));	NXTLEFT	9
	DCL LCNTR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NXTLEFT	10
	DCL NUCCELL	ENTRY RETURNS(FIXED BIN(31,0));	NXTLEFT	11
	DCL NAMTST	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NXTLEFT	12
	DCL(IL,LL)	FIXED BIN(31,0) STATIC;	NXTLEFT	13
	DCL CAL	FIXED BIN(31,0) STATIC;	NXTLEFT	14
	IL=NUCELL;		NXTLEFT	15
	LL=LNKL(CONT(A));		NXTLEFT	16
	CALL SETIND(-1,-1,IL,LL);		NXTLEFT	17
	CALL SETIND(-1,LL,-1,A);		NXTLEFT	18
	CALL SETIND(0,LL,A,IL);		NXTLEFT	19
	IF NAMTST(M)=0 THEN DO;		NXTLEFT	20
		CALL SETIND(1,-1,-1,IL);	NXTLEFT	21
		CALL SETIND(-1,-1,LCNTR(M)+1,M+1); END;	NXTLEFT	22
	CAL=STRIND(M,IL+1);		NXTLEFT	23
	RETURN(IL);		NXTLEFT	24
	END NXTLEFT;		NXTLEFT	25

NXTRGT:		NXTRGT	0
---------	--	--------	---

NXR:	PROC(N,A)	FIXED BIN(31,0);	NXTRGT	1
	DCL(M,A)	FIXED BIN(31,0);	NXTRGT	2
	DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NXTRGT	3
	DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NXTRGT	4
	DCL SETIND	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	NXTRGT	5
		FIXED BIN(31,0),FIXED BIN(31,0))	NXTRGT	6
		RETURNS(FIXED BIN(31,0));	NXTRGT	7
	DCL STRIND	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	NXTRGT	8
		RETURNS(FIXED BIN(31,0));	NXTRGT	9
	DCL LCNTR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NXTRGT	10
	DCL NUCELL	ENTRY RETURNS(FIXED BIN(31,0));	NXTRGT	11
	DCL NAMTST	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	NXTRGT	12
	DCL(IR,LR)	FIXED BIN(31,0) STATIC;	NXTRGT	13
	DCL CAL	FIXED BIN(31,0) STATIC;	NXTRGT	14
	IR=NUCELL;		NXTRGT	15
	LR=LNKR(CONT(A));		NXTRGT	16
	CALL SETIND(-1,IR,-1,LR);		NXTRGT	17
	CALL SETIND(-1,-1,IR,A);		NXTRGT	18
	CALL SETIND(0,A,LR,IR);		NXTRGT	19
	IF NAMTST(M)=C THEN DD;		NXTRGT	20
	CALL SETIND(1,-1,-1,IR);		NXTRGT	21
	CALL SETIND(-1,-1,LCNTR(M)+1,M+1);		NXTRGT	22
	END;		NXTRGT	23
	CAL= STRIND(M,IR+1);		NXTRGT	24
	RETURN(IR);		NXTRGT	25
	END NXTRGT;		NXTRGT	26
POPBDT:			POPBDT	0
PB:	PROC(P)	FIXED BIN(31,0);	POPBDT	1
	DCL (P)	FIXED BIN(31,0);	POPBDT	2
	DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	POPBDT	3
	DCL LNKL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	POPBDT	4

```

DCL LOCT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL DELETE ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
RETURN(DELETE(LNKL(CONT(LOCT(P)))));
END POPBOT;

```

```

POPBOT 5
POPBOT 6
POPBOT 7
POPBOT 8

```

POPTOP:
PT:

```

PROC(P)          FIXED BIN(31,0);
DCL (P)          FIXED BIN(31,0);
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL LOCT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL DELETE ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
RETURN(DELETE(LNKR(CONT(LOCT(P)))));
END POPTOP;

```

```

POPTOP 0
POPTOP 1
POPTOP 2
POPTOP 3
POPTOP 4
POPTOP 5
POPTOP 6
POPTOP 7
POPTOP 8

```

```

RCCELL:  PROC(CELL)          FIXED BIN(31,0);
          /* RETURNS THE CELL AT ADDRESS 'CELL' TO
          LAVS */
DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL SETDIR ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
          FIXED BIN(31,0),FIXED BIN(31,0))
          RETURNS(FIXED BIN(31,0));
DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),
          FIXED BIN(31,0),FIXED BIN(31,0))
          RETURNS(FIXED BIN(31,0));
DCL CELL FIXED BIN(31,0);
DCL AVSL FIXED BIN(31,0) STATIC EXT;
CALL SETIND(-1,-1,CELL,LNKL(AVSL));
CALL SETDIR (-1,CELL,-1,AVSL);
CALL SETIND(-1,-1,0,CELL);

```

```

RCCELL 0
RCCELL 1
RCCELL 2
RCCELL 3
RCCELL 4
RCCELL 5
RCCELL 6
RCCELL 7
RCCELL 8
RCCELL 9
RCCELL 10
RCCELL 11
RCCELL 12
RCCELL 13
RCCELL 14

```

END RCELL;

RCELL 15

```

REED:  PROC(K)          FIXED BIN(31,0);
       DCL (K)          FIXED BIN(31,0);
       DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
       DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
       RETURN(CONT(LNKL(CONT(K))+1));
       END REED;

```

```

REED 0
REED 1
REED 2
REED 3
REED 4
REED 5

```

```

SEQLL: PROC(Z,N)        FIXED BIN(31,0);
       DCL (Z,N)        FIXED BIN(31,0);
       DCL L FIXED BIN(31,0) STATIC;
       DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
       DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
       DCL LNKL ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
       L=LNKL(Z); Z=CONT(L);
       N=ID(Z)-1;
       RETURN (CONT(L+1));
       END SEQLL;

```

```

SEQLL 0
SEQLL 1
SEQLL 2
SEQLL 3
SEQLL 4
SEQLL 5
SEQLL 6
SEQLL 7
SEQLL 8
SEQLL 9

```

```

SEQLR: PROC(Z,N)        FIXED BIN(31,0);
       DCL (Z,N)        FIXED BIN(31,0);
       DCL L FIXED BIN(31,0) STATIC;
       DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
       DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
       DCL ID ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
       L=LNKR(Z); Z=CONT(L);
       N=ID(Z)-1;

```

```

SEQLR 0
SEQLR 1
SEQLR 2
SEQLR 3
SEQLR 4
SEQLR 5
SEQLR 6
SEQLR 7

```


RETURN (CONT(L+1));	SEQLR	8
END SEQLR;	SEQLR	9

SEQRDR: PROC(LST)	FIXED BIN(31,0);	
DCL LST	FIXED BIN(31,0);	0
DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	1
DCL LOCT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	2
RETURN (CONT(LOCT(LST)));		3
END SEQRDR;		4

SEQSL: PROC(Z,N)	FIXED BIN(31,0);	SEQSL	0
DCL (Z,N)	FIXED BIN(31,0);	SEQSL	1
DCL L	FIXED BIN(31,0) STATIC;	SEQSL	2
DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SEQSL	3
DCL ID	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SEQSL	4
DCL LNKL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SEQSL	5
DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SEQSL	6
IF ID(Z)=-1 THEN GOTO L4;		SEQSL	7
L=LNKL(CONT(CONT(LNKL(CONT(LNKR(Z)))+1));		SEQSL	8
GOTO L3;		SEQSL	9
L4: L=LNKL(Z);		SEQSL	10
L3: IF ID(CONT(L))=1 THEN GOTO L2;		SEQSL	11
Z=CONT(L);		SEQSL	12
N=ID(Z)-1;		SEQSL	13
L2: L=LNKL(CONT(CONT(L+1)));		SEQSL	14
RETURN (CONT(L+1));		SEQSL	15
END SEQSL;		SEQSL	16

SEQSR: PROC(Z,N)	FIXED BIN(31,0);	SEQSR	0
------------------	------------------	-------	---

DCL (Z,N)	ENTRY(FIXED BIN(31,0));	SEQSR	1
DCL L	FIXED BIN(31,0) STATIC;	SEQSR	2
DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SEQSR	3
DCL ID	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SEQSR	4
DCL LNKL	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SEQSR	5
DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SEQSR	6
IF ID(Z)≠1	THEN GOTO L4;	SEQSR	7
L=LNKR	(CONT(LNKL(CONT(LNKR(Z)))+1));	SEQSR	8
GOTO	L3;	SEQSR	9
L4:	L=LNKR(Z);	SEQSR	10
L3:	IF ID(CONT(L))=1 THEN GOTO L2;	SEQSR	11
Z=CONT	(L);	SEQSR	12
N=ID	(Z)-1;	SEQSR	13
RETURN	(CONT(L+1));	SEQSR	14
L2:	L=LNKR(CONT(CONT(L+1)));	SEQSR	15
GOTO	L3;	SEQSR	16
END	SEQSR;	SEQSR	17
SUBSBT:	PROC(DAT,LST)	FIXED BIN(31,0);	SUBSBT 0
DCL	(DAT,LST)	FIXED BIN(31,0);	SUBSBT 1
DCL CONT	ENTRY(FIXED BIN(31,0))	RETURNS(FIXED BIN(31,0));	SUBSBT 2
DCL LNKL	ENTRY(FIXED BIN(31,0))	RETURNS(FIXED BIN(31,0));	SUBSBT 3
DCL SUBST	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	RETURNS(FIXED BIN(31,0));	SUBSBT 4
RETURN	(SUBST(DAT,LNKL(CONT(LST))));	SUBSBT 5	
END	SUBSBT;	SUBSBT 6	
SUBST:	PROC(D,N)	FIXED BIN(31,0);	SUBST 0
DCL	(D,N)	FIXED BIN(31,0);	SUBST 1
DCL CONT	ENTRY(FIXED BIN(31,0))	RETURNS(FIXED BIN(31,0));	SUBST 2

DCL STRIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	SUBST	3
RETURNS(FIXED BIN(31,0));	SUBST	4
DCL CAL FIXED BIN(31,0) STATIC;	SUBST	5
DCL SETIND ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	SUBST	6
FIXED BIN(31,0),FIXED BIN(31,0))	SUBST	7
RETURNS(FIXED BIN(31,0));	SUBST	8
DCL LCNTR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SUBST	9
DCL SUB FIXED BIN(31,0) STATIC;	SUBST	10
DCL NAMTST ENTRY(FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	SUBST	11
SUB=CONT(N+1);	SUBST	12
CAL=STRIND(D,N+1);	SUBST	13
CALL SETIND(0,-1,-1,N);	SUBST	14
IF NAMTST(D)=-1 THEN GOTO LA;	SUBST	15
CALL SETIND(1,-1,-1,N);	SUBST	16
CALL SETIND(-1,-1,LCNTR(D)+1,D+1);	SUBST	17
LA: IF NAMTST(SUB)=0 THEN CAL= IRALST(SUB);	SUBST	18
RETURN(SUB);	SUBST	19
END SUBST;	SUBST	20

SUBSTP: PROC(DAT,LST) FIXED BIN(31,0);	SUBSTP	0
DCL (DAT,LST) FIXED BIN(31,0);	SUBSTP	1
DCL CONT ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SUBSTP	2
DCL LNKR ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SUBSTP	3
DCL SUBST ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	SUBSTP	4
RETURNS(FIXED BIN(31,0));	SUBSTP	5
RETURN(SUBST(DAT,LNKR(CONT(LST))));	SUBSTP	6
END SUBSTP;	SUBSTP	7

TOP:		TOP	0
T: PROC(P) FIXED BIN(31,0);		TOP	1

DCL (P)	FIXED BIN(31,0);	TOP	2
DCL CONT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	TOP	3
DCL LNKR	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	TOP	4
DCL LOCT	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	TOP	5
RETURN(CONT(LNKR(CONT(LOCT(P)))+1));	TOP	6
END TOP;		TOP	7

SECTION 12 - PRIMITIVES OF THE SYMMETRIC LIST PROCESSOR

1

A0	L 4,0(1)	INHALT 12
A1	L 2,0(1)	INHALT 13
A2	L 1,0(1)	INHALT 14
INHALT	CSECT	INHALT 0
	EXTRN RELOC	INHALT 1
	SAVE (1,4)	INHALT 2
	BALR 4,0	INHALT 3
	USING *,4	INHALT 4
	L 2,4(1)	INHALT 5
	L 1,0(1) ADDR OF PARAM TO R1	INHALT 6
	L 1,0(1) CONT OF ADDR TO R1	INHALT 7
	SLL 1,17	INHALT 8
	SRL 1,15	INHALT 9
	L 3,ADCON	INHALT 10
	A 1,0(3)	INHALT 11
	MVC 0(4,2),0(1)	INHALT 12
	RETURN (1,4)	INHALT 13
ADCON	DC A(RELOC)	INHALT 14
	END	INHALT 15
SETIND	CSECT	SETIND 0
	EXTRN RELOC	SETIND 1
	SAVE (2,7)	SETIND 2
	BALR 2,0	SETIND 3
	USING *,2	SETIND 4
	L 7,12(1) ADDR OF CELL TO R7	SETIND 5
	L 7,0(7) ACTUAL ADDR OF CELL	SETIND 6
	SLL 7,17	SETIND 7
	SRL 7,15	SETIND 8
	L 3,ADCON	SETIND 9
	A 7,0(3)	SETIND 10
	L 3,0(7) IND CONT OF CELL TO R3	SETIND 11

SECTION 12 - PRIMITIVES OF THE SYMMETRIC LIST PROCESSOR

2

A0	LAVE	4,8(1)		SETIND	12
	LALR	4,0(4)	LNKR TO R4	SETIND	13
	LSING	6,=X'FFFFFFFF'		SETIND	14
	CLR	4,6		SETIND	15
	BC	8,A1	BRANCH ON ZERO	SETIND	16
	SRDL	4,15	SHIFT NEW LNKR TO R5	SETIND	17
	B	A2		SETIND	18
A1	LR	4,3		SETIND	19
	SRDL	4,15	SHIFT OLD LNKR TO R5	SETIND	20
A2	L	4,4(1)		SETIND	21
	L	4,0(4)	LNKL TO R4	SETIND	22
	SRL	3,15		SETIND	23
	CLR	4,6		SETIND	24
	BC	8,A3		SETIND	25
	SRDL	4,15	NEW LNKL TO R5	SETIND	26
	B	A4		SETIND	27
A3	LR	4,3		SETIND	28
	SRDL	4,15	OLD LNKL TO R5	SETIND	29
A4	L	4,0(1)		SETIND	30
A0	L	4,0(4)	LLNKL TO R4	SETIND	31
	SRL	3,15		SETIND	32
	CLR	4,6		SETIND	33
	BC	8,A5	IF ID = -1 THEN GOTO A5	SETIND	34
	SRDL	4,2	NEW ID TO R5	SETIND	35
	B	A6		SETIND	36
A5	LR	4,3		SETIND	37
A1	SRDL	4,2		SETIND	38
A6	ST	5,0(7)		SETIND	39
A2	RETURN	(2,7)		SETIND	40
ACCON	DC	A(RELOC)		SETIND	41
	END			SETIND	42
	CLR	4,6			
	BC	8,A3			
	SRDL	4,15	NEW LNKL TO R5		
STRDIR	CSECT			STRDIR	0

```

SAVE    (1,4)
BALR    2,0
USING   *,2
L        4,8(1)
L        3,0(1)      ADDR OD DATUM TO R3
L        1,4(1)      ADDR OF CELL TO R1
MVC      0(4,1),0(3)  STORE DATUM IN CELL
MVC      0(4,4),0(3)  RETURN DATUM AS FUNCTION VALUE
RETURN  (1,4)
END

```

```

STRDIR   1
STRDIR   2
STRDIR   3
STRDIR   4
STRDIR   5
STRDIR   6
STRDIR   7
STRDIR   8
STRDIR   9
STRDIR  10

```

```

SETDIR   CSECT
SAVE    (2,7)
BALR    2,0
USING   *,2
L        7,12(1)      ADDR OF CELL TO R7
L        3,0(7)      CONT OF CELL TO R3
A0      L        4,8(1)
L        4,0(4)      LNKR TO R4
L        6,=X'FFFFFFFF'
CLR      4,6
BC       8,A1          BRANCH ON ZERO
SRDL     4,15          SHIFT NEW LNKR TO R5
B        A2
A1      LR        4,3
SRDL     4,15          SHIFT OLD LNKR TO R5
A2      L        4,4(1)
L        4,0(4)      LNKL TO R4
SRL      3,15
CLR      4,6
BC       8,A3
SRDL     4,15          NEW LNKL TO R5

```

```

SETDIR   0
SETDIR   1
SETDIR   2
SETDIR   3
SETDIR   4
SETDIR   5
SETDIR   6
SETDIR   7
SETDIR   8
SETDIR   9
SETDIR  10
SETDIR  11
SETDIR  12
SETDIR  13
SETDIR  14
SETDIR  15
SETDIR  16
SETDIR  17
SETDIR  18
SETDIR  19
SETDIR  20

```

SECTION 12 - PRIMITIVES OF THE SYMMETRIC LIST PROCESSOR

4

ADCON	B	A4		SETDIR	21
A3	LR	4,3		SETDIR	22
	SRDL	4,15	OLD LNKL TO R5	SETDIR	23
A4	L	4,0(1)		SETDIR	24
	L	4,0(4)	LLNKL TO R4	SETDIR	25
10	SRL	3,15		SETDIR	26
	CLR	4,6		SETDIR	27
	BC	8,A5	IF ID = -1 THEN GOTO A5	SETDIR	28
	SRDL	4,2	NEW ID TO R5	SETDIR	29
	B	A6		SETDIR	30
A5	LR	4,3		SETDIR	31
	SRDL	4,2		SETDIR	32
A6	ST	5,0(7)		SETDIR	33
	RETURN	(2,7)		SETDIR	34
	END			SETDIR	35

STR IND	CSECT		STR IND	0
	EXTRN RELOC		STR IND	1
	SAVE (1,5)		STR IND	2
	BALR 2,0		STR IND	3
	USING *,2		STR IND	4
	L 4,8(1)		STR IND	5
	L 3,0(1)	ADDR OF DATUM TO R3	STR IND	6
	L 1,4(1)	ADDR OF CELL TO R1	STR IND	7
	L 1,0(1)	ADDR IN CELL TO R1	STR IND	8
	SLL 1,17		STR IND	9
	SRL 1,15		STR IND	10
	L 5,ADCON		STR IND	11
	A 1,0(5)		STR IND	12
	MVC 0(4,1),0(3)	STORE DATUM IN LCC WHOSE ADDR IS IN CELL	STR IND	13
	MVC 0(4,4),0(3)	RETURN DATUM AS FUNCTION VALUE	STR IND	14
	RETURN (1,5)		STR IND	15

SECTION 12 - PRIMITIVES OF THE SYMMETRIC LIST PROCESSOR

5

ADCON	DC A(RELOC)	STRIND	16
	END	STRIND	17
	SRL 1,17		
	ST 1,0(2)		
	RETURN (1,2)		
ID	CSECT	ID	0
	SAVE (1,2)	ID	1
	L 2,4(1)	ID	2
	L 1,0(1)	ID	3
	L 1,0(1)	ID	4
	SRL 1,30	ID	5
	ST 1,0(2)	ID	6
	RETURN (1,2)	ID	7
	END	ID	8
LNKR	CSECT	LNKR	0
	SAVE (1,2)	LNKR	1
	L 2,4(1)	LNKR	2
	L 1,0(1)	LNKR	3
	L 1,0(1)	LNKR	4
	SLL 1,17	LNKR	5
	SRL 1,17	LNKR	6
	ST 1,0(2)	LNKR	7
	RETURN (1,2)	LNKR	8
	END	LNKR	9
LNKL	CSECT	LNKL	0
	SAVE (1,2)	LNKL	1
	L 2,4(1)	LNKL	2
	L 1,0(1)	LNKL	3

SECTION 12 - PRIMITIVES OF THE SYMMETRIC LIST PROCESSOR

6

```

L    1,0(1)
SLL  1,2
SRL  1,17
ST    1,0(2)
RETURN (1,2)
END

```

```

LNKL  4
LNKL  5
LNKL  6
LNKL  7
LNKL  8
LNKL  9

```

```

MADOV CSECT
      SAVE (1,2)
      L    2,4(1)
      L    1,0(1)
      ST    1,0(2)
      RETURN (1,2)
      END

```

```

MADOV 0
MADOV 1
MADOV 2
MADOV 3
MADOV 4
MADOV 5
MADOV 6

```

```

CONT  CSECT
      EXTRN RELOC
      SAVE (1,4)
      BALR  4,0
      USING *,4
      L    2,4(1)
      L    1,0(1)
      L    1,0(1)
      SLL  1,17
      SRL  1,15
      L    3,ADCON
      A    1,0(3)
      MVC  0(4,2),0(1)
      RETURN (1,4)
ADCON DC    A(RELOC)

```

ADDR OF PARAM TO R1
CONT OF ADDR TO R1

```

CONT  0
CONT  1
CONT  2
CONT  3
CONT  4
CONT  5
CONT  6
CONT  7
CONT  8
CONT  9
CONT 10
CONT 11
CONT 12
CONT 13
CONT 14

```

SECTION 12 - PRIMITIVES OF THE SYMMETRIC LIST PROCESSOR

7

END

CONT 15

```

CANONCL: PROC(STR,MOD)      RECURSIVE;
/* PRINT IS 1 IF PRINTING OF INTERMEDIATE STAGES IS WANTED */
/* BETA IS THE PRINCIPAL STRUCTURE */
DCL I
  MODNTR,SUBST,SUBAST,INLSTR,NXTLFT,
  JENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  ASTURN(FIXED BIN(31,0))
DCL J
  LA,L1,L2,R,R1,R2,STR,A,LA
  JENTRY(FIXED BIN(31,0))
DCL K
  STR,SR,SL
  JENTRY(FIXED BIN(31,0))
DCL
  CANONCL,KINWOUT,LINWOUT,RPRODUT,LPRODUT,JACOBI,ORDPROD,
  STRPA,LED,RED,EXPRD
  JENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
DCL
  SYMB
  JENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL
  IDONT,NTERM,CORCHN
  JENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));
DCL
  PERMCON,INVPRD,IRVINV,DRPIONT,CANCEL,
  SUMPA,EXPANV,TRVLX,PRUDEAP,SUNINT,PRODINT
  JENTRY(FIXED BIN(31,0))
DCL
  GOT,CONT,WT,NTEIST,TLP,LEAF
  JENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0))
DCL

```

```

CANONCL 0
CANONCL 1
CANONCL 2
CANONCL 3
CANONCL 4
CANONCL 5
CANONCL 6
CANONCL 7
CANONCL 8
CANONCL 9
CANONCL 10
CANONCL 11
CANONCL 12
CANONCL 13
CANONCL 14
CANONCL 15
CANONCL 16
CANONCL 17
CANONCL 18
CANONCL 19
CANONCL 20
CANONCL 21
CANONCL 22
CANONCL 23
CANONCL 24
CANONCL 25
CANONCL 26
CANONCL 27
CANONCL 28
CANONCL 29
CANONCL 30

```

SECTION 13 - PROCEDURES OF THE COLLECTING PROCESS

1

```

CANONCL: PROC(STR,MOD)      RECURSIVE;
/* PRNT IS 1 IF PRINTING OF INTERMEDIATE STAGES IS WANTED */
/* BETA IS THE PRINCIPAL STRUCTURE */
DCL (
  MADNTP,SUBST,SUBSBT,INLSTR,NXTLFT
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
DCL (
  CAL,L,LR,R,RR,MOD,STR,A,LA
)  FIXED BIN(31,0);
DCL (
  SYM,SR,SL
)  CHAR(1);
DCL(
  CANONCL,RINVOUT,LINVOUT,RPRDOUT,LPRDOUT,JACOBI,ORDPROD,
  STRPN,LED,RED,EXPRD
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0));
DCL (
  SYMB
)ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL (
  IDNT,NTerm,COMCHK
)ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));
DCL (
  PERMCOM,INVPRD,INVINV,DRPIDNT,CANCEL,
  SUMEX,EXPINV,INVEX,PRODEXP,SUMINT,PRODINT
)ENTRY(FIXED BIN(31,0));
DCL (
  BDT,CONT,WT,MILIST,TOP,LNKR
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (

```

```

CANONCL 31
CANONCL 32
CANONCL 33
CANONCL 0
CANONCL 1
CANONCL 2
CANONCL 3
CANONCL 4
CANONCL 5
CANONCL 6
CANONCL 7
CANONCL 8
CANONCL 9
CANONCL 10
CANONCL 11
CANONCL 12
CANONCL 13
CANONCL 14
CANONCL 15
CANONCL 16
CANONCL 17
CANONCL 18
CANONCL 19
CANONCL 20
CANONCL 21
CANONCL 22
CANONCL 23
CANONCL 24
CANONCL 25
CANONCL 26
CANONCL 27
CANONCL 28
CANONCL 29
CANONCL 30

```


EQ,LT	CANONCL 31
ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	CANONCL 32
RETURNS(BIT(1));	CANONCL 33
DCL (CANONCL 34
LG	CANONCL 35
ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))	CANONCL 36
RETURNS(FIXED BIN(31,0));	CANONCL 37
DCL SCAN ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	CANONCL 38
FIXED BIN(31,0),FIXED BIN(31,0));	CANONCL 39
DCL PRNT FIXED BIN(31,0) EXT;	CANONCL 40
DCL BETA FIXED BIN(31,0) EXT;	CANONCL 41
DCL DI FIXED DEC;	CANONCL 42
DCL C3 CHAR(3);	CANONCL 43
DCL CT CHAR(2),DT FIXED DEC;	CANONCL 44
DCL JJJ STATIC INITIAL(0);	CANONCL 45
IF PRNT=1 THEN CALL STRPN(BETA,0);	CANONCL 46
IF PRNT=1 THEN JJJ=1;	CANONCL 47
JJJ=JJJ+1; IF JJJ>50 THEN	CANONCL 48
DO; JJJ=0; CALL STRPN(BETA,0); END;	CANONCL 49
IF -NTERM(STR) THEN RETURN;	CANONCL 50
L=CONT(MADNTP(STR,2)+1); R=BOT(STR); SYM=SYMB(STR);	CANONCL 51
	CANONCL 52
	CANONCL 53
IF SYM=''' THEN DO;	CANONCL 54
IF -NTERM(R) THEN DO;	CANONCL 55
IF IDNT(R) THEN DO; CALL CANCEL(STR); GOTO RET; END;END;	CANONCL 56
SR=SYMB(R);	CANONCL 57
IF SR='.' THEN DO;	CANONCL 58
CALL INVPRD(STR);	CANONCL 59
CALL CANONCL(CONT(MADNTP(STR,2)+1),MOD);	CANONCL 60
CALL CANONCL(BOT(STR),MOD); GOTO RET; END;	CANONCL 61
IF SR=''' THEN DO; CALL INVINV(STR); GOTO RET; END;	CANONCL 62
IF SR=':' THEN DO; CALL INVEX(STR); GOTO RET; END;	CANONCL 63
RETURN;	CANONCL 64

CALL REDU; END; /*' ' '*/	CANONCL 65
IF SYMB(L)=' ' THEN DO;	CANONCL 66
IF NTERM(L) THEN	CANONCL 67
IF SYMB(L)=' ' & BOT(L)=R THEN	CANONCL 68
DO; CALL CANCEL(STR); GOTO RET; END;	CANONCL 69
IF NTERM(R) THEN	CANONCL 70
IF SYMB(R)=' ' & BOT(R)=L THEN	CANONCL 71
DO; CALL CANCEL(STR); GOTO RET; END;	CANONCL 72
IF -NTERM(L) & IDNT(L) THEN	CANONCL 73
DO; CALL DRPIDNT(STR); GOTO RET; END;	CANONCL 74
IF -NTERM(R) & IDNT(R) THEN	CANONCL 75
DO; CALL DRPIDNT(STR); GOTO RET; END;	CANONCL 76
CALL ORDPD(STR,MOD); RETURN;	CANONCL 77
END; /*' ' '*/	CANONCL 78
	CANONCL 79
	CANONCL 80
	CANONCL 81
	CANONCL 82
IF SYMB(L)=' ' THEN DO;	CANONCL 83
IF IDNT(L) IDNT(R) THEN DO;	CANONCL 84
CAL=CANCEL(STR); GOTO RET; END;	CANONCL 85
IF SYMB(L)=':' & SYMB(R)=':' THEN	CANONCL 86
IF EQ(L,R) THEN DO; CAL=CANCEL(STR); GOTO RET; END;	CANONCL 87
IF SYMB(L)=';' & SYMB(R)=';' THEN	CANONCL 88
IF EQ(L,R) THEN DO; CAL=CANCEL(STR); GOTO RET; END;	CANONCL 89
IF SYMB(L)=' ' THEN DO;	CANONCL 90
CAL=LINVOUT(STR,MOD); GOTO RES; END;	CANONCL 91
IF SYMB(R)=' ' THEN DO;	CANONCL 92
CAL=RINVOUT(STR,MOD); GOTO RES; END;	CANONCL 93
IF SYMB(L)='.' THEN DO;	CANONCL 94
CAL=LPRDOUT(STR,MOD); GOTO RET; END;	CANONCL 95
IF SYMB(R)='.' THEN DO;	CANONCL 96
CAL=RPRDOUT(STR,MOD); GOTO RET; END;	CANONCL 97
IF SYMB(R)=':' THEN DO;	CANONCL 98

```

      CALL RED(STR,MOD); GOTO RET; END;
    IF SYMB(L)=':' THEN DO;
      CALL LED(STR,MOD); GOTO RET; END;
    IF COMCHK(L) & COMCHK(R) & LT(L,R) THEN DO;
      CAL=PERMCOM(STR); GOTO RET; END;
    LR=CONT(MADNTP(L,2)+1); SR=SYMB(L); RR=BOT(L);
    IF SR='.' THEN RETURN;
    IF COMCHK(L) & COMCHK(R) & COMCHK(LR)
      & LT(R,RR) & LT(RR,LR) THEN DO;
      CALL JACOBI(STR,MOD); GOTO RES; END;
    RETURN;
  END; /* */

```

```

IF SYM=':' THEN DO;
  SR=SYMB(L);
  IF SR='.' THEN DO; CALL EXPINV(STR); GOTO RES; END;
  IF SR=':' THEN DO; CALL PRODEXP(STR); GOTO RES; END;
  IF SR='.' THEN DO;
    RR=TOP(R); UNSPEC(SL)=SUBSTR(UNSPEC(RR),9,8);
    IF SUBSTR(UNSPEC(RR),9,24)=UNSPEC('DOU') THEN DO;
      IF SUBSTR(UNSPEC(BOT(BOT(R))),9,8)≠UNSPEC('I')
        THEN GOTO K2;
      UNSPEC(CT)=SUBSTR(UNSPEC(BOT(BOT(R))),17,16);
      DI=CT; L,LR=LG('.',L,0); GOTO K1;
    END;
    IF SL='I' THEN DO;
      UNSPEC(CT)=SUBSTR(UNSPEC(RR),17,16);
      DI=CT; LR=L; GOTO K1;
    END;
  END;

```

```

K2: CALL EXPRD(STR,MOD); GOTO RET;
K1: DO I=2 TO DI; LR=LG('.',LR,L); END;
    CAL=IRALST(INLSTR(LR,MTLIST(STR)));
    GOTO RES; END;

```

CANONCL 99
 CANONCL 100
 CANONCL 101
 CANONCL 102
 CANONCL 103
 CANONCL 104
 CANONCL 105
 CANONCL 106
 CANONCL 107
 CANONCL 108
 CANONCL 109
 CANONCL 110
 CANONCL 111
 CANONCL 112
 CANONCL 113
 CANONCL 114
 CANONCL 115
 CANONCL 116
 CANONCL 117
 CANONCL 118
 CANONCL 119
 CANONCL 120
 CANONCL 121
 CANONCL 122
 CANONCL 123
 CANONCL 124
 CANONCL 125
 CANONCL 126
 CANONCL 127
 CANONCL 128
 CANONCL 129
 CANONCL 130
 CANONCL 131
 CANONCL 132

```

C3=SUBSTR(UNSPEC(TOP(R)),9,24);
IF C3='I 0' | C3='I00' THEN DO; CAL=CANCEL(STR); END;
RETURN; END;

```

```

IF SYM='+' | SYM='-' THEN DO;
CALL SUMINT(STR); RETURN; END;

```

```

IF SYM='U' THEN DO;
IF SYMB(L)='+' | SYMB(L)='-' THEN DO;
CALL SUMINT(L); GOTO RET; END;
IF -NTERM(L) THEN RETURN;
IF SYMB(L)~='U' THEN RETURN;
A=BOT(L); CAL=MTLIST(STR);
LR=LNKR(A); LA=LNKR(CONT(LR));
DO I=1 TO 4;
CAL=NXTLEFT(CONT(MADNTP(A,I)+1),STR);
LA=LNKR(CONT(LA)); IF LA=LR THEN GOTO RET;
END;
END;

```

```

IF SYM='*' THEN DO; CALL PRODINT(STR); RETURN; END;
IF SYM='|' THEN RETURN;

```

```

SIGNAL CONDITION(ALPHA);

```

```

RET: IF PRNT=1 THEN CALL STRPN(BETA,0); RETURN;
RES: IF PRNT=1 THEN CALL STRPN(BETA,0);
CALL SCAN(STR,STR,0,MOD);
RETURN;

```

CANONCL133
 CANONCL134
 CANONCL135
 CANONCL136
 CANONCL137
 CANONCL138
 CANONCL139
 CANONCL140
 CANONCL141
 CANONCL142
 CANONCL143
 CANONCL144
 CANONCL145
 CANONCL146
 CANONCL147
 CANONCL148
 CANONCL149
 CANONCL150
 CANONCL151
 CANONCL152
 CANONCL153
 CANONCL154
 CANONCL155
 CANONCL156
 CANONCL157
 CANONCL158
 CANONCL159
 CANONCL160
 CANONCL161
 CANONCL162
 CANONCL163
 CANONCL164
 CANONCL165
 CANONCL166


```
END CANONCL; FIXED BIN(31,0) RETURNS(CHAR(1));
```

CANONCL167

```
DCL (
  GRPL, COMCHK
  ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));
```

```
CMUTPRD: PROC(STR);
```

```
/* AB -> BA(A,B) */
```

```
DCL (
  CAL
  ) ENTRY(FIXED BIN(31,0));
```

```
DCL (
  A,B,L,STR
  ) ENTRY(FIXED BIN(31,0));
```

```
DCL (
  BOT,CONT,DELETE
  ) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
```

```
DCL (
  MADNTP
  ) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
```

```
DCL (
  LG
  ) ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
```

```
A=CONT(MADNTP(STR,2)+1);
```

```
B=BOT(STR);
```

```
STR=LG(' ',LG(' ',B,A),LG(' ',A,B));
```

```
END CMUTPRD;
```

CMUTPRD 0

CMUTPRD 1

CMUTPRD 2

CMUTPRD 3

CMUTPRD 4

CMUTPRD 5

CMUTPRD 6

CMUTPRD 7

CMUTPRD 8

CMUTPRD 9

CMUTPRD 10

CMUTPRD 11

CMUTPRD 12

CMUTPRD 13

CMUTPRD 14

CMUTPRD 15

CMUTPRD 16

CMUTPRD 17

CMUTPRD 18

CMUTPRD 19

CMUTPRD 20

CMUTPRD 21

CMUTPRD 22

```
COMCHK: PROC(COM) BIT(1) RECURSIVE;
```

```
/* RETURNS '1'B IFF COM IS A COMMUTATOR */
```

```
DCL (
  SYMB
```

COMCHK 0

COMCHK 1

COMCHK 2

COMCHK 3

```

    )ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL (
    GRPL,COMCHK
    )ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));
DCL (
    COM
    ) FIXED BIN(31,0);
DCL (
    BOT,CONT
    )ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
    MADNTP
    )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
IF GRPL(COM) THEN RETURN('1'B);
IF SYMB(COM)='&'&SYMB(COM)='&' THEN RETURN('0'B);
IF SYMB(COM)='&' THEN RETURN(COMCHK(BOT(COM)));
RETURN(COMCHK(BOT(COM))&COMCHK(CONT(MADNTP(COM,2)+1)));
END COMCHK;

```

```

COMCHK 4
COMCHK 5
COMCHK 6
COMCHK 7
COMCHK 8
COMCHK 9
COMCHK 10
COMCHK 11
COMCHK 12
COMCHK 13
COMCHK 14
COMCHK 15
COMCHK 16
COMCHK 17
COMCHK 18
COMCHK 19
COMCHK 20
COMCHK 21
COMCHK 22

```

```

COMCNT: PROC(STR)      CHAR(1)  RECURSIVE;
/* PROC ENTERS COMMA COUNTS IN THE STRUCTURE */
DCL (
    A,B,STR,WTB,CAL
    ) FIXED BIN(31,0);
DCL (
    COMCNT
    )ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL (
    SYMB
    )ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL (

```

```

COMCNT 0
COMCNT 1
COMCNT 2
COMCNT 3
COMCNT 4
COMCNT 5
COMCNT 6
COMCNT 7
COMCNT 8
COMCNT 9
COMCNT 10
COMCNT 11

```

MRKA,MRKB	COMCNT	12
) CHAR(1);	COMCNT	13
DCL (COMCNT	14
SUBSTP,MADNTP	COMCNT	15
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	COMCNT	16
RETURNS(FIXED BIN(31,0));	COMCNT	17
DCL (COMCNT	18
WT,CONT,BOT,TOP	COMCNT	19
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	COMCNT	20
DCL MRKC CHAR(4);	COMCNT	21
DCL MRK(4) PACKED CHAR(1) DEFINED MRKC ;	COMCNT	22
DCL (COMCNT	23
NTERM	COMCNT	24
)ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));	COMCNT	25
DCL INK ENTRY(CHAR(1))RETURNS(CHAR(1));	COMCNT	26
ON CONVERSION SIGNAL CONDITION (ALPHA);	COMCNT	27
IF ~NTERM(STR) THEN RETURN('0');	COMCNT	28
A=CONT(MADNTP(STR,2)+1);	COMCNT	29
B=BOT(STR);	COMCNT	30
UNSPEC(MRKC)=UNSPEC(TOP(STR));	COMCNT	31
MRKA=COMCNT(A); MRKB=COMCNT(B);	COMCNT	32
IF SYMB(STR)~=',' THEN RETURN('0');	COMCNT	33
IF MRKA >MRKB THEN MRK(3)=MRKA; ELSE MRK(3)=INK(MRKB);	COMCNT	34
CAL=SUBSTP(UNSPEC(MRKC),STR); RETURN(MRK(3));	COMCNT	35
END COMCNT;	COMCNT	36
COPYPRD: PROC(STR) FIXED BIN(31,0) RECURSIVE;	COPYPRD	0
	COPYPRD	1
DCL STR FIXED BIN(31,0);	COPYPRD	2
	COPYPRD	3
DCL (COPYPRD	4
	COPYPRD	5

```

      LG
      )ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
      RETURNS(FIXED BIN(31,0));

DCL (
CONT,BOT,COPYPRD
      )ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));

DCL (
      MADNTP
      )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
      RETURNS(FIXED BIN(31,0));

DCL (
      SYMB
      )ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));

DCL (
      NTERM
      )ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));

IF -NTERM(STR) THEN RETURN(STR);

```

```

COPYPRD 6
COPYPRD 7
COPYPRD 8
COPYPRD 9
COPYPRD 10
COPYPRD 11
COPYPRD 12
COPYPRD 13
COPYPRD 14
COPYPRD 15
COPYPRD 16
COPYPRD 17
COPYPRD 18
COPYPRD 19
COPYPRD 20
COPYPRD 21
COPYPRD 22
COPYPRD 23
COPYPRD 24
COPYPRD 25
COPYPRD 26
COPYPRD 27
COPYPRD 28
COPYPRD 29
COPYPRD 30
COPYPRD 31
COPYPRD 32
COPYPRD 33
COPYPRD 34
COPYPRD 35
COPYPRD 36
COPYPRD 37
COPYPRD 38
COPYPRD 39

```



```

IF SYMB(STR)~='.' THEN RETURN(STR);
RETURN(LG('.',COPYPRD(CONT(MADNTP(STR,2)+1))
      ,COPYPRD(BOT(STR)))));
END COPYPRD;

```

```

COPYPRD 40
COPYPRD 41
COPYPRD 42
COPYPRD 43
COPYPRD 44
COPYPRD 45
COPYPRD 46

```

```

DRPIDNT: PROC(STR);
DCL(
  CAL,STR,J,LA,LR,I
  ) FIXED BIN(31,0);
DCL(
  TOP,BOT,CONT,MTLIST,LNKR
  ) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL(
  MADNTP,INLSTR,NXTLFT
  ) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
IF SUBSTR(UNSPEC(TOP(BOT(STR))),9,8)=UNSPEC('U')
THEN J=CONT(MADNTP(STR,2)+1);
ELSE J=BOT(STR);
CAL=MTLIST(STR); LR=LNKR(J); LA=LNKR(CONT(LR));
DO I=1 TO 4;
  CAL=NXTLFT(CONT(MADNTP(J,I)+1),STR);
  LA=LNKR(CONT(LA)); IF LA=LR THEN RETURN;
END;
SIGNAL CONDITION(ALPHA);
END DRPIDNT;

```

```

DRPIDNT 0
DRPIDNT 1
DRPIDNT 2
DRPIDNT 3
DRPIDNT 4
DRPIDNT 5
DRPIDNT 6
DRPIDNT 7
DRPIDNT 8
DRPIDNT 9
DRPIDNT 10
DRPIDNT 11
DRPIDNT 12
DRPIDNT 13
DRPIDNT 14
DRPIDNT 15
DRPIDNT 16
DRPIDNT 17
DRPIDNT 18
DRPIDNT 19
DRPIDNT 20

```

```

/* RETURNS '1'B IFF A=B */
DCL (ANT,BNT) BIT(1);
DCL (
    MADNTP
    )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL (
    L,A,B
    )    FIXED BIN(31,0);
DCL (
    NTERM
    )ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));
DCL (
    SYMB
    )ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL (
    CONT,BDT
    )ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
    EQ
    )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(BIT(1));
DCL (SA,SB) CHAR(1);
ANT=NTERM(A); BNT=NTERM(B);
IF -ANT & -BNT THEN RETURN(SYMB(A)=SYMB(B));
SA=SYMB(A); SB=SYMB(B);
IF SA=':' | SA=''' THEN RETURN(EQ(CONT(MADNTP(A,2)+1),B));
IF SB=':' | SB=''' THEN RETURN(EQ(A,CONT(MADNTP(B,2)+1)));
IF SA=- SB THEN RETURN('0'B);
RETURN(EQ(CONT(MADNTP(A,2)+1),CONT(MADNTP(B,2)+1))
    &EQ(BDT(A),BDT(B)));
END EQ;

```

DCL (EXP INV	1
STR,A,N,CAL	EXP INV	2
) FIXED BIN(31,0);	EXP INV	3
DCL (EXP INV	4
BOT, TOP, CONT, MTLIST	EXP INV	5
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	EXP INV	6
DCL (EXP INV	7
LG	EXP INV	8
) ENTRY(CHAR(1), FIXED BIN(31,0), FIXED BIN(31,0))	EXP INV	9
RETURNS(FIXED BIN(31,0));	EXP INV	10
DCL (EXP INV	11
MADNTP, INLSTR	EXP INV	12
) ENTRY(FIXED BIN(31,0), FIXED BIN(31,0))	EXP INV	13
RETURNS(FIXED BIN(31,0));	EXP INV	14
DCL IRALST ENTRY(FIXED BIN(31,0));	EXP INV	15
A=BOT(CONT(MADNTP(STR,2)+1));	EXP INV	16
N=BOT(STR);	EXP INV	17
IF SUBSTR(UNSPEC(TOP(N)),9,24)=UNSPEC('00U')	EXP INV	18
THEN A=LG(':',A,BOT(N));	EXP INV	19
ELSE A=LG(':',A,LG('U',N,0));	EXP INV	20
CAL=IRALST(INLSTR(A,MTLIST(STR)));	EXP INV	21
END EXP INV;	EXP INV	22

EXPRD:	PROC(STR,MOD);	EXPRD	0
	/* THIS PROCEDURE GENERATES ACANONIC EXPANSION OF	EXPRD	1
	(A.B):N MOD WEIGHT 7 OR LESS	EXPRD	2
	*/	EXPRD	3
DCL (STR, EX, L, LA, A, B, N, BA, BAA, BAB, BAAA, BAAB, BABB, BAAAA,	EXPRD	4
	BAAAB, BAABB, BABBB, BAABA, BABBA, BAAAAA, BAAAAAB, BAAABB,	EXPRD	5
	BAABBB, BABBBB, BAAABA, BAABBA, BABBBBA, BABBBAA,	EXPRD	6
	MOD, BETS, CAL, WA, WB,	EXPRD	7
	12,13,14,15,16,17,18,19,110,	EXPRD	8

I12,I13,I18,I21,I20,I24,I27,I30,I36,I52,I54	EXPRD	9
) FIXED BIN(31,0);	EXPRD	10
DCL (EXPRD	11
LG	EXPRD	12
)ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))	EXPRD	13
RETURNS(FIXED BIN(31,0));	EXPRD	14
DCL (EXPRD	15
MADNTP,INLSTR,NXTLFT	EXPRD	16
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	EXPRD	17
RETURNS(FIXED BIN(31,0));	EXPRD	18
DCL (EXPRD	19
WT,CONT,BOT,MTLIST,LIST	EXPRD	20
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	EXPRD	21
DCL PRNTI BIT(1) EXT;	EXPRD	22
DCL BETA FIXED BIN(31,0) EXT;	EXPRD	23
DCL (EXPRD	24
CANONCL	EXPRD	25
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0));	EXPRD	26
DCL(EXPRD	27
STRPN	EXPRD	28
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0));	EXPRD	29
IF PRNTI THEN PUT SKIP(2) LIST('ENTER EXPRD');	EXPRD	30
L=CONT(MADNTP(STR,2)+1); N=BOT(STR);	EXPRD	31
A=CONT(MADNTP(L,2)+1); B=BOT(L);	EXPRD	32
I2=LIST(9); CAL =NXTLFT(UNSPEC('0I02'),I2);	EXPRD	33
I3=LIST(9); CAL =NXTLFT(UNSPEC('0I03'),I3);	EXPRD	34
I4=LIST(9); CAL =NXTLFT(UNSPEC('0I04'),I4);	EXPRD	35
I5=LIST(9); CAL =NXTLFT(UNSPEC('0I05'),I5);	EXPRD	36
I6=LIST(9); CAL =NXTLFT(UNSPEC('0I06'),I6);	EXPRD	37
I7=LIST(9); CAL =NXTLFT(UNSPEC('0I07'),I7);	EXPRD	38
I8=LIST(9); CAL =NXTLFT(UNSPEC('0I08'),I8);	EXPRD	39
I9=LIST(9); CAL =NXTLFT(UNSPEC('0I09'),I9);	EXPRD	40
I10=LIST(9); CAL =NXTLFT(UNSPEC('0I10'),I10);	EXPRD	41
I12=LIST(9); CAL =NXTLFT(UNSPEC('0I12'),I12);	EXPRD	42

I13=LIST(9); CAL =NXTLFT(UNSPEC('0I13'),I13);	EXPRD	43
I18=LIST(9); CAL =NXTLFT(UNSPEC('0I18'),I18);	EXPRD	44
I20=LIST(9); CAL =NXTLFT(UNSPEC('0I20'),I20);	EXPRD	45
I21=LIST(9); CAL =NXTLFT(UNSPEC('0I21'),I21);	EXPRD	46
I24=LIST(9); CAL =NXTLFT(UNSPEC('0I24'),I24);	EXPRD	47
I27=LIST(9); CAL =NXTLFT(UNSPEC('0I27'),I27);	EXPRD	48
I30=LIST(9); CAL =NXTLFT(UNSPEC('0I30'),I30);	EXPRD	49
I36=LIST(9); CAL =NXTLFT(UNSPEC('0I36'),I36);	EXPRD	50
I52=LIST(9); CAL =NXTLFT(UNSPEC('0I52'),I52);	EXPRD	51
I54=LIST(9); CAL =NXTLFT(UNSPEC('0I54'),I54);	EXPRD	52
IF PRNTI THEN CALL STRPN(STR,0);	EXPRD	53
BETS=BETA;	EXPRD	54
LA=LG(' ',A,N);	EXPRD	55
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD	56
BETA=LA; CALL CANONCL(LA,MOD);	EXPRD	57
L=LG(' ',B,N);	EXPRD	58
IF PRNTI THEN CALL STRPN(L,0);	EXPRD	59
BETA=L; CALL CANONCL(L,MOD);	EXPRD	60
LA=LG(' ',LA,L);	EXPRD	61
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD	62
BETA=LA; CALL CANONCL(LA,MOD);	EXPRD	63
WA=WT(A); WB=WT(B);	EXPRD	64
	EXPRD	65
IF WA+WB< MOD THEN DO;	EXPRD	66
BA=LG(' ',B,A);	EXPRD	67
IF PRNTI THEN CALL STRPN(BA,0);	EXPRD	68
BETA=BA; CALL CANONCL(BA,MOD);	EXPRD	69
EX=LG(' ',N,I2);	EXPRD	70
L=LG(' ',BA,EX);	EXPRD	71
IF PRNTI THEN CALL STRPN(L,0);	EXPRD	72
BETA=L; CALL CANONCL(L,MOD);	EXPRD	73
LA=LG(' ',LA,L);	EXPRD	74
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD	75
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD	76

```

IF 2*WA+WB < MOD THEN DO;
  BAA=LG(' ', ' ', BA, A);
  IF PRNTI THEN CALL STRPN(BAA, 0);
  BETA=BAA; CALL CANONCL(BAA, MOD);
  EX=LG(' | ', N, I3);
  L=LG(' : ', BAA, EX);
  IF PRNTI THEN CALL STRPN(L, 0);
  BETA=L; CALL CANONCL(L, MOD);
  LA=LG(' % ', LA, L);
  IF PRNTI THEN CALL STRPN(LA, 0);
  BETA=LA; CALL CANONCL(LA, MOD); END;

```

```

IF WA+2*WB < MOD THEN DO;
  BAB=LG(' ', ' ', BA, B);
  IF PRNTI THEN CALL STRPN(BAB, 0);
  BETA=BAB; CALL CANONCL(BAB, MOD);
  EX=LG(' + ', LG(' | ', N, I2),
    LG(' * ', I2, LG(' | ', N, I3)));
  L=LG(' : ', BAB, EX);
  IF PRNTI THEN CALL STRPN(L, 0);
  BETA=L; CALL CANONCL(L, MOD);
  LA=LG(' % ', LA, L);
  IF PRNTI THEN CALL STRPN(LA, 0);
  BETA=LA; CALL CANONCL(LA, MOD); END;

```

```

IF 3*WA+WB < MOD THEN DO;
  BAAA=LG(' ', ' ', BAA, A);
  IF PRNTI THEN CALL STRPN(BAAA, 0);
  BETA=BAAA; CALL CANONCL(BAAA, MOD);
  EX=LG(' | ', N, I4);
  L=LG(' : ', BAAA, EX);
  IF PRNTI THEN CALL STRPN(L, 0);
  BETA=L; CALL CANONCL(L, MOD);

```

```

EXPRD 77
EXPRD 78
EXPRD 79
EXPRD 80
EXPRD 81
EXPRD 82
EXPRD 83
EXPRD 84
EXPRD 85
EXPRD 86
EXPRD 87
EXPRD 88
EXPRD 89
EXPRD 90
EXPRD 91
EXPRD 92
EXPRD 93
EXPRD 94
EXPRD 95
EXPRD 96
EXPRD 97
EXPRD 98
EXPRD 99
EXPRD 100
EXPRD 101
EXPRD 102
EXPRD 103
EXPRD 104
EXPRD 105
EXPRD 106
EXPRD 107
EXPRD 108
EXPRD 109
EXPRD 110

```

LA=LG(' ',LA,L);	EXPRD 111
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 112
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 113
IF 2*WA+2*WB < MOD THEN DO;	EXPRD 114
BAAB=LG(' ',BAA,B);	EXPRD 115
IF PRNTI THEN CALL STRPN(BAAB ,0);	EXPRD 116
BETA=BAAB ; CALL CANONCL(BAAB ,MOD);	EXPRD 117
EX=LG(' '+',LG(' '*',I2,LG(' ' ',N,I3)),	EXPRD 118
LG(' '*',I3,LG(' ' ',N,I4)));	EXPRD 119
L=LG(' ':',BAAB,EX);	EXPRD 120
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 121
BETA=L; CALL CANONCL(L,MOD);	EXPRD 122
LA=LG(' ',LA,L);	EXPRD 123
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 124
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 125
IF WA+3*WB < MOD THEN DO;	EXPRD 126
BABB=LG(' ',BAB,B);	EXPRD 127
IF PRNTI THEN CALL STRPN(BABB ,0);	EXPRD 128
BETA=BABB ; CALL CANONCL(BABB ,MOD);	EXPRD 129
EX=LG(' '+',LG(' '*',I2,LG(' ' ',N,I3)),	EXPRD 130
LG(' '*',I3,LG(' ' ',N,I4)));	EXPRD 131
L=LG(' ':',BABB,EX);	EXPRD 132
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 133
BETA=L; CALL CANONCL(L,MOD);	EXPRD 134
LA=LG(' ',LA,L);	EXPRD 135
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 136
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 137
IF 4*WA+WB < MOD THEN DO;	EXPRD 138
BAAAA=LG(' ',BAAA,A);	EXPRD 139
IF PRNTI THEN CALL STRPN(BAAAA ,0);	EXPRD 140
BETA=BAAAA ; CALL CANONCL(BAAAA ,MOD);	EXPRD 141
	EXPRD 142
	EXPRD 143
	EXPRD 144

EX=LG(' ',N,15);	EXPRD 145
L=LG(':',BAAAA,EX);	EXPRD 146
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 147
BETA=L; CALL CANONCL(L,MOD);	EXPRD 148
LA=LG('.',LA,L);	EXPRD 149
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 150
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 151
IF 3*WA+2*WB < MOD THEN DO;	EXPRD 152
BAAAAB=LG('',BAAA,B);	EXPRD 153
IF PRNTI THEN CALL STRPN(BAAAAB,0);	EXPRD 154
BETA=BAAAAB; CALL CANONCL(BAAAAB,MOD);	EXPRD 155
EX=LG('+',LG('* ',I3,LG(' ',N,I4)),	EXPRD 156
LG('* ',I4,LG(' ',N,I5)))	EXPRD 157
L=LG(':',BAAAAB,EX);	EXPRD 158
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 159
BETA=L; CALL CANONCL(L,MOD);	EXPRD 160
LA=LG('.',LA,L);	EXPRD 161
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 162
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 163
IF 2*WA+3*WB < MOD THEN DO;	EXPRD 164
BAABBB=LG('',BAAB,B);	EXPRD 165
IF PRNTI THEN CALL STRPN(BAABBB,0);	EXPRD 166
BETA=BAABBB; CALL CANONCL(BAABBB,MOD);	EXPRD 167
EX=LG('+',LG(' ',N,I3),	EXPRD 168
LG('+',LG('* ',I6,LG(' ',N,I4)),	EXPRD 169
LG('* ',I6,LG(' ',N,I5)))	EXPRD 170
L=LG(':',BAABBB,EX);	EXPRD 171
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 172
BETA=L; CALL CANONCL(L,MOD);	EXPRD 173
LA=LG('.',LA,L);	EXPRD 174
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 175
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 176
	EXPRD 177
	EXPRD 178

IF WA+4*WB < MOD THEN DO;	EXPRD 179
BABBB=LG(' ', ' ', BABB, B);	EXPRD 180
IF PRNTI THEN CALL STRPN(BABBB, 0);	EXPRD 181
BETA=BABBB; CALL CANONCL(BABBB, MOD);	EXPRD 182
EX=LG('+' , LG('*' , I3, LG(' ' , N, I4)),	EXPRD 183
LG('*' , I4, LG(' ' , N, I5)));	EXPRD 184
L=LG(':' , BABBB, EX);	EXPRD 185
IF PRNTI THEN CALL STRPN(L, 0);	EXPRD 186
BETA=L; CALL CANONCL(L, MOD);	EXPRD 187
LA=LG('.' , LA, L);	EXPRD 188
IF PRNTI THEN CALL STRPN(LA, 0);	EXPRD 189
BETA=LA; CALL CANONCL(LA, MOD); END;	EXPRD 190
	EXPRD 191
	EXPRD 192
IF 3*WA+2*WB < MOD THEN DO;	EXPRD 193
BAABA=LG(' ', ' ', BAA, BA);	EXPRD 194
IF PRNTI THEN CALL STRPN(BAABA, 0);	EXPRD 195
BETA=BAABA; CALL CANONCL(BAABA, MOD);	EXPRD 196
EX=LG('+' , LG(' ' , N, I3),	EXPRD 197
LG('+' , LG('*' , I7, LG(' ' , N, I4)),	EXPRD 198
LG('*' , I6, LG(' ' , N, I5)));	EXPRD 199
L=LG(':' , BAABA, EX);	EXPRD 200
IF PRNTI THEN CALL STRPN(L, 0);	EXPRD 201
BETA=L; CALL CANONCL(L, MOD);	EXPRD 202
LA=LG('.' , LA, L);	EXPRD 203
IF PRNTI THEN CALL STRPN(LA, 0);	EXPRD 204
BETA=LA; CALL CANONCL(LA, MOD); END;	EXPRD 205
	EXPRD 206
IF 2*WA+3*WB < MOD THEN DO;	EXPRD 207
BABBA=LG(' ', ' ', BAB, BA);	EXPRD 208
IF PRNTI THEN CALL STRPN(BABBA, 0);	EXPRD 209
BETA=BABBA; CALL CANONCL(BABBA, MOD);	EXPRD 210
EX=LG('+' , LG('*' , I6, LG(' ' , N, I3)), LG('+' , LG('*' ,	EXPRD 211
I18, LG(' ' , N, I4)), LG('*' , I12, LG(' ' , N, I5)));	EXPRD 212

L=LG(':',BABBA,EX);	EXPRD 213
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 214
BETA=L; CALL CANONCL(L,MOD);	EXPRD 215
LA=LG('.'. ,LA,L);	EXPRD 216
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 217
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 218
	EXPRD 219
IF 5*WA+WB < MOD THEN DO;	EXPRD 220
BAAAAA=LG(' ',BAAAA,A);	EXPRD 221
IF PRNTI THEN CALL STRPN(BAAAAA,0);	EXPRD 222
BETA=BAAAAA; CALL CANONCL(BAAAAA,MOD);	EXPRD 223
EX=LG(' ' ',N,16);	EXPRD 224
L=LG(' ':',BAAAAA,EX);	EXPRD 225
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 226
BETA=L; CALL CANONCL(L,MOD);	EXPRD 227
LA=LG('.'. ,LA,L);	EXPRD 228
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 229
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 230
	EXPRD 231
IF 4*WA+2*WB < MOD THEN DO;	EXPRD 232
BAAAAAB=LG(' ',BAAAA,B);	EXPRD 233
IF PRNTI THEN CALL STRPN(BAAAAAB,0);	EXPRD 234
BETA=BAAAAAB; CALL CANONCL(BAAAAAB,MOD);	EXPRD 235
EX=LG(' '+',LG(' '*',14,LG(' ' ',N,15)),	EXPRD 236
LG(' '*',15,LG(' ' ',N,16)))	EXPRD 237
L=LG(' ':',BAAAAAB,EX);	EXPRD 238
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 239
BETA=L; CALL CANONCL(L,MOD);	EXPRD 240
LA=LG('.'. ,LA,L);	EXPRD 241
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 242
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 243
	EXPRD 244
IF 3*WA+3*WB < MOD THEN DO;	EXPRD 245
BAAABB=LG(' ',BAAAB,B);	EXPRD 246

IF PRNTI THEN CALL STRPN(BAAABB,0);	EXPRD 247
BETA=BAAABB; CALL CANONCL(BAAABB,MOD);	EXPRD 248
EX=LG('+',LG('*',I3,LG(' ',N,I4)),	EXPRD 249
LG('+',LG('*',I12,LG(' ',N,I5)),	EXPRD 250
LG('*',I10,LG(' ',N,I6))));	EXPRD 251
L=LG(':',BAAABB,EX);	EXPRD 252
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 253
BETA=L; CALL CANONCL(L,MOD);	EXPRD 254
LA=LG('.',LA,L);	EXPRD 255
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 256
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 257
IF 2*WA+4*WB < MOD THEN DO;	EXPRD 258
BAABBB=LG(' ',BAABB,B);	EXPRD 259
IF PRNTI THEN CALL STRPN(BAABBB,0);	EXPRD 260
BETA=BAABBB; CALL CANONCL(BAABBB,MOD);	EXPRD 261
EX=LG('+',LG('*',I3,LG(' ',N,I4)),	EXPRD 262
LG('+',LG('*',I12,LG(' ',N,I5)),	EXPRD 263
LG('*',I10,LG(' ',N,I6))));	EXPRD 264
L=LG(':',BAABBB,EX);	EXPRD 265
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 266
BETA=L; CALL CANONCL(L,MOD);	EXPRD 267
LA=LG('.',LA,L);	EXPRD 268
IF PRNTI THEN CALL STRPN(LA,0);	EXPRD 269
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 270
IF WA+5*WB < MOD THEN DO;	EXPRD 271
BABBBB=LG(' ',BABBB,B);	EXPRD 272
IF PRNTI THEN CALL STRPN(BABBBB,0);	EXPRD 273
BETA=BABBBB; CALL CANONCL(BABBBB,MOD);	EXPRD 274
EX=LG('+',LG('*',I4,LG(' ',N,I5)),	EXPRD 275
LG('*',I5,LG(' ',N,I6))));	EXPRD 276
L=LG(':',BABBBB,EX);	EXPRD 277
IF PRNTI THEN CALL STRPN(L,0);	EXPRD 278
	EXPRD 279
	EXPRD 280

BETA=L; CALL CANONCL(L,MOD);	EXPRD 281
LA=LG(' ',LA,L);	EXPRD 282
IF PRNT1 THEN CALL STRPN(LA,0);	EXPRD 283
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 284
	EXPRD 285
IF 4*WA+2*WB < MOD THEN DO;	EXPRD 286
BAAABA=LG(' ',BAAA,BA);	EXPRD 287
IF PRNT1 THEN CALL STRPN(BAAABA,0);	EXPRD 288
BETA=BAAABA; CALL CANONCL(BAAABA,MOD);	EXPRD 289
EX=LG(' ',LG(' '*',I3,LG(' ' ',N,I4)),	EXPRD 290
LG(' '+',LG(' '*',I13,LG(' ' ',N,I5)),	EXPRD 291
LG(' '*',I10,LG(' ' ',N,I6)))));	EXPRD 292
L=LG(' ':',BAAABA,EX);	EXPRD 293
IF PRNT1 THEN CALL STRPN(L,0);	EXPRD 294
BETA=L; CALL CANONCL(L,MOD);	EXPRD 295
LA=LG(' ',LA,L);	EXPRD 296
IF PRNT1 THEN CALL STRPN(LA,0);	EXPRD 297
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 298
	EXPRD 299
IF 3*WA+3*WB < MOD THEN DO;	EXPRD 300
BAABBA=LG(' ',BAAB,BA);	EXPRD 301
IF PRNT1 THEN CALL STRPN(BAABBA,0);	EXPRD 302
BETA=BAABBA; CALL CANONCL(BAABBA,MOD);	EXPRD 303
EX=LG(' ',LG(' '*',I2,LG(' ' ',N,I3)),	EXPRD 304
LG(' '+',LG(' '*',I24,LG(' ' ',N,I4)),	EXPRD 305
LG(' '+',LG(' '*',I52,LG(' ' ',N,I5)),	EXPRD 306
LG(' '*',I30,LG(' ' ',N,I6)))));	EXPRD 307
L=LG(' ':',BAABBA,EX);	EXPRD 308
IF PRNT1 THEN CALL STRPN(L,0);	EXPRD 309
BETA=L; CALL CANONCL(L,MOD);	EXPRD 310
LA=LG(' ',LA,L);	EXPRD 311
IF PRNT1 THEN CALL STRPN(LA,0);	EXPRD 312
BETA=LA; CALL CANONCL(LA,MOD); END;	EXPRD 313
	EXPRD 314

IF 2*WA+4*WB < MOD THEN DO;	EXPRD 315
BABBBBA=LG(' ', ' ', BABB, BA);	EXPRD 316
IF PRNTI THEN CALL STRPN(BABBBBA, 0);	EXPRD 317
BETA=BABBBBA; CALL CANONCL(BABBBBA, MOD);	EXPRD 318
EX=LG('+' , LG('*' , 13, LG(' ' , N, 13)),	EXPRD 319
LG('+' , LG('*' , 127, LG(' ' , N, 14)),	EXPRD 320
LG('+' , LG('*' , 154, LG(' ' , N, 15)),	EXPRD 321
LG('*' , 130, LG(' ' , N, 16)))));	EXPRD 322
L=LG(':' , BABBBBA, EX);	EXPRD 323
IF PRNTI THEN CALL STRPN(L, 0);	EXPRD 324
BETA=L; CALL CANONCL(L, MOD);	EXPRD 325
LA=LG('.' , LA, L);	EXPRD 326
IF PRNTI THEN CALL STRPN(LA, 0);	EXPRD 327
BETA=LA; CALL CANONCL(LA, MOD); END;	EXPRD 328
	EXPRD 329
IF 3*WA+3*WB < MOD THEN DO;	EXPRD 330
BABBAA=LG(' ', ' ', BAB, BAA);	EXPRD 331
IF PRNTI THEN CALL STRPN(BABBAA, 0);	EXPRD 332
BETA=BABBAA; CALL CANONCL(BABBAA, MOD);	EXPRD 333
EX=LG('+' , LG('*' , 14, LG(' ' , N, 13)),	EXPRD 334
LG('+' , LG('*' , 121, LG(' ' , N, 14)),	EXPRD 335
LG('+' , LG('*' , 136, LG(' ' , N, 15)),	EXPRD 336
LG('*' , 120, LG(' ' , N, 16)))));	EXPRD 337
L=LG(':' , BABBAA, EX);	EXPRD 338
IF PRNTI THEN CALL STRPN(L, 0);	EXPRD 339
BETA=L; CALL CANONCL(L, MOD);	EXPRD 340
LA=LG('.' , LA, L);	EXPRD 341
IF PRNTI THEN CALL STRPN(LA, 0);	EXPRD 342
BETA=LA; CALL CANONCL(LA, MOD); END;	EXPRD 343
	EXPRD 344
CAL=IRALST(INLSTR(LA, MTLIST(STR)));	EXPRD 345
BETA=BETS;	EXPRD 346
IF PRNTI THEN PUT SKIP(2) LIST('EXIT' EXPRD);	EXPRD 347
END EXPRD;	EXPRD 348

GRPL:	PROC(STR) BIT(1);	GRPL	0
	/* RETURNS '1'B IFF STR IS A GROUP ELEMENT */	GRPL	1
	DCL (GRPL	2
	CH	GRPL	3
) CHAR(1);	GRPL	4
	DCL (GRPL	5
	STR	GRPL	6
) FIXED BIN(31,0);	GRPL	7
	DCL (GRPL	8
	TOP	GRPL	9
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	GRPL	10
	UNSPEC(CH)=SUBSTR(UNSPEC(TOP(STR)),9,8);	GRPL	11
	RETURN(CH='G' CH='U');	GRPL	12
	END GRPL;	GRPL	13
IDNT:	PROC(STR) BIT(1);	IDNT	0
	/* RETURNS '1'B IFF STR IS THE IDENTITY */	IDNT	1
	DCL (IDNT	2
	STR	IDNT	3
) FIXED BIN(31,0);	IDNT	4
	DCL (IDNT	5
	TOP	IDNT	6
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	IDNT	7
	RETURN(SUBSTR(UNSPEC(TOP(STR)),9,8)=UNSPEC('U'));	IDNT	8
	END IDNT;	IDNT	9
INVEX:	PROC(STR) FIXED BIN(31,0);	INVEX	0

DCL (INVEX	1
STR,A,B,N	INVEX	2
) FIXED BIN(31,0);	INVEX	3
DCL (INVEX	4
BOT,CONT,TOP,MTLIST	INVEX	5
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INVEX	6
DCL (INVEX	7
LG	INVEX	8
)ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))	INVEX	9
RETURNS(FIXED BIN(31,0));	INVEX	10
DCL (INVEX	11
MADNTP,INLSTR	INVEX	12
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	INVEX	13
RETURNS(FIXED BIN(31,0));	INVEX	14
/* (A:N) -> A:-N */	INVEX	15
B=BOT(STR);	INVEX	16
A=CONT(MADNTP(B,2)+1);	INVEX	17
N=BOT(B);	INVEX	18
IF SUBSTR(UNSPEC(TOP(N)),9,24)=UNSPEC('000')	INVEX	19
THEN A=LG(':',A,BOT(N));	INVEX	20
ELSE A=LG(':',A,LG('U',N,0));	INVEX	21
CALL IKALST(INLSTR(A,MTLIST(STR)));	INVEX	22
END INVEX;	INVEX	23
INVINV: PROC(STR);	INVINV	0
/* A -> A */	INVINV	1
DCL (INVINV	2
STR,A,CAL,LA,LR,I	INVINV	3
) FIXED BIN(31,0);	INVINV	4
DCL (INVINV	5
BOT,MTLIST,LNKR,CONT	INVINV	6
)ENTRY(FIXED BIN(31,0))RETURNS(FIXED BIN(31,0));	INVINV	7

DCL (INVINV	8
INLSTR,NXTLFT,MADNTP	INVINV	9
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	INVINV	10
RETURNS(FIXED BIN(31,0));	INVINV	11
A=BOT(BOT(STR));	INVINV	12
CAL=MTLIST(STR); LR=LNKR(A); LA=LNKR(CONT(LR));	INVINV	13
DO I=1 TO 4;	INVINV	14
CAL=NXTLFT(CONT(MADNTP(A,I)+1),STR);	INVINV	15
LA=LNKR(CONT(LA)); IF LA=LR THEN RETURN;	INVINV	16
END;	INVINV	17
SIGNAL CONDITION(ALPHA);	INVINV	18
END INVINV;	INVINV	19

INVPRD: PROC(STR);	INVPRD	0
/* (AB) ⁰ -> B ⁰ A ⁰ */	INVPRD	1
DCL (INVPRD	2
L	INVPRD	3
) FIXED BIN(31,0);	INVPRD	4
DCL (INVPRD	5
A,B,STR,CAL,LA,LB	INVPRD	6
) FIXED BIN(31,0);	INVPRD	7
DCL (INVPRD	8
LG	INVPRD	9
)ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))	INVPRD	10
RETURNS(FIXED BIN(31,0));	INVPRD	11
DCL (INVPRD	12
BOT,CONT,MTLIST	INVPRD	13
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	INVPRD	14
DCL (INVPRD	15
MADNTP,INLSTR	INVPRD	16
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	INVPRD	17
RETURNS(FIXED BIN(31,0));	INVPRD	18

DCL (INVPRD 19
SYMB	INVPRD 20
)ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));	INVPRD 21
A=CONT(MADNTP(BOT(STR),2)+1);	INVPRD 22
B=BOT(BOT(STR));	INVPRD 23
IF SYMB(A)=' ' THEN LA=CONT(MADNTP(A,2)+1);	INVPRD 24
ELSE LA=LG(' ',A,0);	INVPRD 25
IF SYMB(B)=' ' THEN LB=CONT(MADNTP(B,2)+1);	INVPRD 26
ELSE LB=LG(' ',B,0);	INVPRD 27
A=LG(' ',LB,LA);	INVPRD 28
CALL IRALST(INLSTR(A,MTLIST(STR)));	INVPRD 29
END INVPRD;	INVPRD 30

JACOBI: PROC(STR,MOD) FIXED BIN(31,0);
 /* JACOBI APPLIES THE JACOBI IDENTITY TO THE STRUCTURE STR */

DCL (JACOBI 0
A,B,C,L,V1,V2,V3,CAL,WA,WB,WC,BETS	JACOBI 1
) FIXED BIN(31,0);	JACOBI 2
DCL (AC,CB,BA,CA,BAC,CAB) FIXED BIN(31,0);	JACOBI 3
DCL (JACOBI 4
STR,MOD	JACOBI 5
) FIXED BIN(31,0);	JACOBI 6
DCL (JACOBI 7
CONT,BOT,WT,LIST,MLIST	JACOBI 8
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	JACOBI 9
DCL (JACOBI 10
MADNTP,NXILFT,INLSTR	JACOBI 11
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	JACOBI 12
RETURNS(FIXED BIN(31,0));	JACOBI 13
DCL (JACOBI 14
LG	JACOBI 15
)ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))	JACOBI 16
	JACOBI 17
	JACOBI 18

```

      RETURNS(FIXED BIN(31,0));
DCL (STRPN,CANONCL) ENTRY (FIXED BIN(31,0),FIXED BIN(31,0));
DCL PRNT1 BIT(1) EXT,BETA FIXED BIN(31,0)EXT;
IF PRNT1 THEN PUT SKIP(2) LIST('ENTER JACOBI');
IF PRNT1 THEN CALL STRPN(STR,0);
BETS=BETA;
A=BOT(STR);
L=CONT(MADNTP(STR,2)+1);
B=BOT(L);
C=CONT(MADNTP(L,2)+1);
AC=LG(' ',A,C); BETA=AC; CALL CANONCL(AC,MOD);
CB=LG(' ',C,B); BETA=CB; CALL CANONCL(CB,MOD);
BA=LG(' ',B,A); BETA=BA; CALL CANONCL(BA,MOD);
CA=LG(' ',C,A); BETA=CA; CALL CANONCL(CA,MOD);
BAC=LG(' ',BA,C); BETA=BAC; CALL CANONCL(BAC,MOD);
CAB=LG(' ',CA,B); BETA=CAB; CALL CANONCL(CAB,MOD);
WA=WT(A); WB=WT(B); WC=WT(C);
IF WA+WB+2*WC >= MOD THEN DO;
  L=LIST(9); CAL=NXTLFT(UNSPEC('0000'),L); END;
IF WA+WB+2*WC < MOD THEN L=LG(' ',AC,CB);
IF WA+2*WB+2*WC < MOD THEN L=LG(' ',L,LG(' ',CB,BAC));
L=LG(' ',L,LG(' ',BAC,0));
IF WA+2*WB+WC < MOD THEN L=LG(' ',L,LG(' ',CB,BA));
IF 2*WA+2*WB+WC < MOD THEN L=LG(' ',L,LG(' ',BA,LG(' ',AC,B)));
IF 2*WA+WB+2*WC < MOD THEN
  L=LG(' ',L,LG(' ',AC,LG(' ',CA,B),0));
L=LG(' ',L,CAB);
IF 2*WA+WB+WC < MOD THEN L=LG(' ',L,LG(' ',BA,AC));
IF 2*WA+WB+2*WC < MOD THEN L=LG(' ',L,LG(' ',AC,LG(' ',CB,A)));
IF PRNT1 THEN CALL STRPN(L,0);
CALL IRALST(INLSTR(L,MLIST(STR)));
IF PRNT1 THEN PUT SKIP(2) LIST('EXIT JACOBI');
BETA=BETS;
END JACOBI;

```

JACOBI 19
 JACOBI 20
 JACOBI 21
 JACOBI 22
 JACOBI 23
 JACOBI 24
 JACOBI 25
 JACOBI 26
 JACOBI 27
 JACOBI 28
 JACOBI 29
 JACOBI 30
 JACOBI 31
 JACOBI 32
 JACOBI 33
 JACOBI 34
 JACOBI 35
 JACOBI 36
 JACOBI 37
 JACOBI 38
 JACOBI 39
 JACOBI 40
 JACOBI 41
 JACOBI 42
 JACOBI 43
 JACOBI 44
 JACOBI 45
 JACOBI 46
 JACOBI 47
 JACOBI 48
 JACOBI 49
 JACOBI 50
 JACOBI 51
 JACOBI 52

```

LEO:  PROC(STR,MOD);
      /* LEFT EXPONENT OUT */
      /* N IS THE GREATEST INTEGER SUCH THAT
         N.WT(B)+WT(A) < MOD */
      DCL K FIXED BIN(31,0);
      DCL(
        STR,MOD,A,B,X,BX,N,L,LA,LB,J,CAL,WTB,BETS,
        BAB,BABB,BA,LC,L1,L2,L3,L4
      )FIXED BIN(31,0);
      DCL(
        WT,CONT,BOT,LIST,MILIST,TCP
      )ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
      DCL(
        NXLFT,MADNTP,INLSTR
      )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
        RETURNS(FIXED BIN(31,0));
      DCL(
        LG
      )ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
        RETURNS(FIXED BIN(31,0));
      DCL INT BIT(1);
      DCL KC CHAR(10), (KT,NX) CHAR(2);
      DCL PRNTI BIT(1) EXT;
      DCL(
        STRPN
      )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0));
      DCL(
        CANDNCL
      )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0));
      DCL NCR ENTRY(CHAR(2),CHAR(2)) RETURNS(FIXED BIN(31,0));
      DCL BETA FIXED BIN(31,0) EXT;

```

```

LEO 0
LEO 1
LEO 2
LEO 3
LEO 4
LEO 5
LEO 6
LEO 7
LEO 8
LEO 9
LEO 10
LEO 11
LEO 12
LEO 13
LEO 14
LEO 15
LEO 16
LEO 17
LEO 18
LEO 19
LEO 20
LEO 21
LEO 22
LEO 23
LEO 24
LEO 25
LEO 26
LEO 27
LEO 28
LEO 29
LEO 30

```

DCL SCNP ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	LEO	31
LA=LG(0,0,0,0); FIXED BIN(31,0),FIXED BIN(31,0));	LEO	32
IF PRNTI THEN PUT SKIP(2) LIST('ENTER LEO');	LEO	33
IF PRNTI THEN CALL STRPN(STR,0);	LEO	34
BETS=BETA; BETA=L; INT='0'B;	LEO	35
BX=CONT(MADNTP(STR,2)+1); A=BOT(STR);	LEO	36
B=CONT(MADNTP(BX,2)+1); X=BOT(BX);	LEO	37
LA=LIST(9); CAL=NXTLEFT(UNSPEC('0000'),LA);	LEO	38
LB=TOP(X);	LEO	39
IF SUBSTR(UNSPEC(LB),9,8)=UNSPEC('I') THEN DO;	LEO	40
UNSPEC(NX)=SUBSTR(UNSPEC(LB),17,16); INT='1'B;	LEO	41
IF SUBSTR(NX,1,1)='0' THEN SUBSTR(NX,1,1)=' ';	LEO	42
END;	LEO	43
IF INT THEN IF NX<'1' THEN GOTO S2;	LEO	44
WTA=WT(A); WTB=WT(B); IF WTA+WTB>= MOD THEN GOTO S2;	LEO	45
IF MOD>7 THEN SIGNAL CONDITION(ALPHA);	LEO	46
CAL=IRALST(LA);	LEO	47
L=LG(0,0,B,A);	LEO	48
IF PRNTI THEN CALL STRPN(L,0);	LEO	49
BETA=L; CALL CANONCL(L,MOD);	LEO	50
BA,LA=L;	LEO	51
IF INT THEN IF NX='1' THEN GOTO S2;	LEO	52
LA=LG(0,0,L,X);	LEO	53
IF PRNTI THEN CALL STRPN(LA,0);	LEO	54
BETA=LA; CALL CANONCL(LA,MOD);	LEO	55
N=(MOD-1-WTA)/WTB; KC=(9)'0' '1';	LEO	56
BABB=0;	LEO	57
DO J=2 TO N;	LEO	58
KC=KC+'1'; KT=SUBSTR(KC,8,2);	LEO	59
IF INT THEN IF NX<KT THEN GOTO S1;	LEO	60
L=LG(0,0,L,B);	LEO	61
IF PRNTI THEN CALL STRPN(L,0);	LEO	62
BETA=L; CALL CANONCL(L,MOD);	LEO	63
IF J=2 THEN BAB=L; IF J=3 THEN BABB=L;	LEO	64

IF INT THEN IF NX=KT THEN DO;	LEO	65
LA=LG(' ',LA,L);	LEO	66
IF PRNTI THEN CALL STRPN(LA,0);	LEO	67
GOTO S1; END;	LEO	68
IF INT THEN LB=NCR(NX,KT);	LEO	69
ELSE DO; LB=LIST(9); CAL=NXTLFT(UNSPEC('01' KT),LB);	LEO	70
LB=LG(' ' X,LB);	LEO	71
END;	LEO	72
LB=LG(' ',L,LB);	LEO	73
IF PRNTI THEN CALL STRPN(LB,0);	LEO	74
BETA=LB; CALL CANONCL(LB,MOD);	LEO	75
LA=LG(' ',LA,LB);	LEO	76
END;	LEO	77
S1: IF 2*WTA+3*WTB < MOD THEN DO;	LEO	78
L=LG(' ',BAB,BA);	LEO	79
BETA=L; CALL CANONCL(L,MOD);	LEO	80
L1=LIST(9); CAL=NXTLFT(UNSPEC('0101'),L1);	LEO	81
L3=LIST(9); CAL=NXTLFT(UNSPEC('0103'),L3);	LEO	82
IF INT THEN DO;	LEO	83
KT=SUBSTR(CHAR(NX+'01'),8,2);	LEO	84
LC=LG(' ',NCR(NX,'03'),NCR(KT,'03'));	LEO	85
BETA=LC; CALL CANONCL(LC,MOD);	LEO	86
END;	LEO	87
ELSE DO;	LEO	88
LC=LG(' ',LG(' ' X,L3),LG(' ' LG(' '+X,L1),L3));	LEO	89
END;	LEO	90
L=LG(' ',L,LC);	LEO	91
BETA=L; CALL CANONCL(L,MOD);	LEO	92
LA=LG(' ',LA,L);	LEO	93
END;	LEO	94
IF 2*WTA+4*WTB < MOD THEN DO;	LEO	95
IF BABB=0 THEN GOTO S2;	LEO	96
L=LG(' ',BABB,BA);	LEO	97
	LEO	98

```

DCL BETA=L; CALL CANONCL(L,MOD);
L2=LIST(9); CAL=NXTLFT(UNSPEC('0102'),L2);
L4=LIST(9); CAL=NXTLFT(UNSPEC('0104'),L4);
DCL IF INT THEN DO;
    LC=LG('*',L2,NCR(KT,'04'));
    BETA=LC; CALL CANONCL(LC,MOD);
    LC=LG('+',NCR(NX,'04'),LC);
DCL BETA=LC; CALL CANONCL(LC,MOD);
    END;
ELSE DO;
    LC=LG('+',LG(']',X,L4),
    LG('*',L2,LG(']',LG('+',X,L1),L4)));
    END;
L=LG(':',L,LC);
BETA=L; CALL CANONCL(L,MOD);
LA=LG('.',LA,L);
END;

```

```

S2: BETA=LA;
CALL SCNP(LA,LA,0,MOD);
IF PRNTI THEN CALL STRPN(LA,0);
CAL=IRALST(INLSTR(LA,MILIST(STR)));
BETA=BETS;
IF PRNTI THEN PUT SKIP(2) LIST('EXIT LEO');
END LEO;

```

```

LG: PRDC(SYMB,LEFT,RIGHT) FIXED BIN(31,0);
/* LG GENERATES A LIST STRUCTURE OF STANDARD FORMAT */
DCL (
    LEFT,RIGHT
) FIXED BIN(31,0);
DCL SYMB CHAR(1);

```

```

LEO      99
LEO      100
LEO      101
LEO      102
LEO      103
LEO      104
LEO      105
LEO      106
LEO      107
LEO      108
LEO      109
LEO      110
LEO      111
LEO      112
LEO      113
LEO      114
LEO      115
LEO      116
LEO      117
LEO      118
LEO      119
LEO      120
LEO      121
LEO      122
LEO      123

```

```

LG      0
LG      1
LG      2
LG      3
LG      4
LG      5

```

```

DCL (
    A,CAL
) ENTRY(FIXED BIN(31,0));
DCL (
    NXTLEFT
) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL (
    LSSCPY,LIST
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
A=LIST(9);
IF SYMB=' ' || (SYMB='U' & RIGHT=0)
    THEN CAL=NXTLEFT(UNSPEC('000' || SYMB),NXTLEFT(LEFT,A));
    ELSE CAL=NXTLEFT(UNSPEC('000' || SYMB),
        NXTLEFT(LEFT,NXTLEFT(RIGHT,A)));
RETURN(A);
END LG;

```

```

LG 6
LG 7
LG 8
LG 9
LG 10
LG 11
LG 12
LG 13
LG 14
LG 15
LG 16
LG 17
LG 18
LG 19
LG 20
LG 21
LG 22

```

```

LINVDOUT: PROC(STR,MOD);
/* A,B ->(A,B)^(B,A,A) */
DCL (
    A,B,L,CAL,STR,MOD
) ENTRY(FIXED BIN(31,0));
DCL (
    LG
) ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL (
    CONT,BOT,WT,MTLIST
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
    MADNTP,INLSTR

```

```

LINVDOUT 0
LINVDOUT 1
LINVDOUT 2
LINVDOUT 3
LINVDOUT 4
LINVDOUT 5
LINVDOUT 6
LINVDOUT 7
LINVDOUT 8
LINVDOUT 9
LINVDOUT 10
LINVDOUT 11
LINVDOUT 12
LINVDOUT 13

```

```

L=0) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
A=BOT(CONT(MADNTP(STR,2)+1));
B=BOT(STR);
L=LG('','','',LG('','',A,B),0);
IF WT(B)+2*WT(A)<MOD THEN
  L=LG('','',L,LG('','',LG('','',B,A),LG('','',A,0)));
CALL IRALST(INLSTR(L,MTLIST(STR)));
END LINVOUT;

```

```

LINVOUT 14
LINVOUT 15
LINVOUT 16
LINVOUT 17
LINVOUT 18
LINVOUT 19
LINVOUT 20
LINVOUT 21
LINVOUT 22

```

```

LPRDDUT: PROC(STR,MOD);
/* AB,C -> (A,C)(A,C,B)(B,C) */
DCL (
  C,L,B,A,CAL,STR,MOD,WA,WB,WC,BETS,LA
) LG FIXED BIN(31,0);
DCL (
  BOT,CONT,WT,LIST,MTLIST
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
  MADNTP,NXTLEFT,INLSTR
) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
DCL (
  LG
) ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
DCL (STRPN,CANONCL) ENTRY (FIXED BIN(31,0),FIXED BIN(31,0));
DCL PRNTI BIT(1) EXT,BETA FIXED BIN(31,0)EXT;
IF PRNTI THEN PUT SKIP(2) LIST('ENTER LPRDDUT');
IF PRNTI THEN CALL STRPN(STR,0);
BETS=BETA;
C=BOT(STR);

```

```

LPRDDUT 0
LPRDDUT 1
LPRDDUT 2
LPRDDUT 3
LPRDDUT 4
LPRDDUT 5
LPRDDUT 6
LPRDDUT 7
LPRDDUT 8
LPRDDUT 9
LPRDDUT 10
LPRDDUT 11
LPRDDUT 12
LPRDDUT 13
LPRDDUT 14
LPRDDUT 15
LPRDDUT 16
LPRDDUT 17
LPRDDUT 18
LPRDDUT 19
LPRDDUT 20
LPRDDUT 21

```



```

L=CONT(MADNTP(STR,2)+1);
B=BUT(L);
A=CONT(MADNTP(L,2)+1);
WA=WT(A); WB=WT(B); WC=WT(C);
IF WA+WC >= MOD THEN DO;
  L=LIST(9); CAL=NXTLFT(UNSPEC('0000'),L); END;
IF WA+WC < MOD THEN DO;
  L=LG(' ', ' ', A, C);
  IF PRNTI THEN CALL STRPN(L,0);
  BETA=L; CALL CANONCL(L,MOD); END;
IF WA+WB+WC < MOD THEN DO;
  LA=LG(' ', ' ', L, B);
  IF PRNTI THEN CALL STRPN(LA,0);
  BETA=LA; CALL CANONCL(LA,MOD);
  L=LG(' ', ' ', L, LA); END;
IF WB+WC < MOD THEN DO;
  LA=LG(' ', ' ', B, C);
  IF PRNTI THEN CALL STRPN(LA,0);
  BETA=LA; CALL CANONCL(LA,MOD);
  L=LG(' ', ' ', L, LA);
  IF PRNTI THEN CALL STRPN(L,0);
  BETA=L; CALL CANONCL(L,MOD); END;
CALL IRALST(INLSTR(L,MTLIST(STR)));
IF PRNTI THEN PUT SKIP(2) LIST('EXIT LPRDOUT');
BETA=BETS;
END LPRDOUT;

```

```

LPRDOUT 22
LPRDOUT 23
LPRDOUT 24
LPRDOUT 25
LPRDOUT 26
LPRDOUT 27
LPRDOUT 28
LPRDOUT 29
LPRDOUT 30
LPRDOUT 31
LPRDOUT 32
LPRDOUT 33
LPRDOUT 34
LPRDOUT 35
LPRDOUT 36
LPRDOUT 37
LPRDOUT 38
LPRDOUT 39
LPRDOUT 40
LPRDOUT 41
LPRDOUT 42
LPRDOUT 43
LPRDOUT 44
LPRDOUT 45
LPRDOUT 46
LPRDOUT 47

```

```

LSTSTR: PROC(LST) FIXED BIN(31,0);
/* THIS PROCEDURE CONVERTS THE LINEAR LIST 'LST'
   TO A STANDARD BINARY STRUCTURE. */
DCL (
  LST,PH,P1,LA,P2,CAL

```

```

LSTSTR 0
LSTSTR 1
LSTSTR 2
LSTSTR 3
LSTSTR 4

```

```

/* L) RET FIXED BIN(31,0); */
DCL (
  LG, TOP, CONT, BOT
  )ENTRY(CHAR(1), FIXED BIN(31,0), FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
DCL (
  LNKR, LIST, CONT, LCNTR
  )ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
  NXTLEFT, MADNBT
  )ENTRY(FIXED BIN(31,0), FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
DCL SETIND ENTRY(FIXED BIN(31,0), FIXED BIN(31,0),
  FIXED BIN(31,0), FIXED BIN(31,0));
PH, P1=LNKR(LST);
P1=LNKR(CONT(P1));
IF P1=PH THEN DO;
  LA=LIST(9); CAL=NXTLEFT(UNSPEC('OU00'), LA);
  RETURN(LA); END;
P2=LNKR(CONT(P1));
IF P2=PH THEN DO;
  LA=CONT(MADNBT(CONT(P1+1), 2)+1);
  CALL SETIND(-1, -1, LCNTR(LA)+1, LA+1);
  RETURN(LA); END;
LA=LG('.', CONT(MADNBT(CONT(P1+1), 2)+1),
  CONT(MADNBT(CONT(P2+1), 2)+1));
L1:
P2=LNKR(CONT(P2));
IF P2=PH THEN RETURN(LA);
LA=LG('.', LA, CONT(MADNBT(CONT(P2+1), 2)+1));
GOTO L1;
END LSTSTR;

```

```

LSTSTR 5
LSTSTR 6
LSTSTR 7
LSTSTR 8
LSTSTR 9
LSTSTR 10
LSTSTR 11
LSTSTR 12
LSTSTR 13
LSTSTR 14
LSTSTR 15
LSTSTR 16
LSTSTR 17
LSTSTR 18
LSTSTR 19
LSTSTR 20
LSTSTR 21
LSTSTR 22
LSTSTR 23
LSTSTR 24
LSTSTR 25
LSTSTR 26
LSTSTR 27
LSTSTR 28
LSTSTR 29
LSTSTR 30
LSTSTR 31
LSTSTR 32
LSTSTR 33
LSTSTR 34
LSTSTR 35

```

```

/* LT RETURNS '1'B IFF ALB */
DCL (
    WT, TOP, CONT, BOT
) ENTRY (FIXED BIN(31,0)) RETURNS (FIXED BIN(31,0));
DCL (
    WTA, WTB, A1, B1, A2, B2, A, B, LDA, LDB, TRA, TRB
) ENTRY (FIXED BIN(31,0));
DCL (
    MADNTP
) ENTRY (FIXED BIN(31,0), FIXED BIN(31,0))
    RETURNS (FIXED BIN(31,0));
DCL (
    LT, EQ
) ENTRY (FIXED BIN(31,0), FIXED BIN(31,0))
    RETURNS (BIT(1));
DCL SYMB ENTRY (FIXED BIN(31,0)) RETURNS (CHAR(1));
DCL (SA, SB) CHAR(1);
WTA=WT(A); WTB=WT(B); SA=SYMB(A); SB=SYMB(B);
IF WTA>WTB THEN RETURN('0'B);
IF WTA<WTB THEN RETURN('1'B);
IF SA=':' | SA=''' THEN RETURN(LT(CONT(MADNTP(A,2)+1), B));
IF SB=':' | SB=''' THEN RETURN(LT(A, CONT(MADNTP(B,2)+1)));
IF WTA=1 THEN RETURN(SA<SB);
A1=CONT(MADNTP(A,2)+1); A2=BOT(A);
B1=CONT(MADNTP(B,2)+1); B2=BOT(B);
IF LT(A1, A2)
    THEN DO; LDA=A2; TRA=A1; END;
    ELSE DO; LDA=A1; TRA=A2; END;
IF LT(B1, B2)
    THEN DO; LDB=B2; TRB=B1; END;
    ELSE DO; LDB=B1; TRB=B2; END;
RETURN(LT(LDA, LDB) |
    (EQ(LDA, LDB) & LT(TRA, TRB)) |
    (EQ(LDA, LDB) & EQ(TRA, TRB) & LT(B1, A1)));

```

LT	1
LT	2
LT	3
LT	4
LT	5
LT	6
LT	7
LT	8
LT	9
LT	10
LT	11
LT	12
LT	13
LT	14
LT	15
LT	16
LT	17
LT	18
LT	19
LT	20
LT	21
LT	22
LT	23
LT	24
LT	25
LT	26
LT	27
LT	28
LT	29
LT	30
LT	31
LT	32
LT	33
LT	34

```

END LT;  NUMBER OF COMBINATIONS OF N THINGS R AT A TIME */
DCL (L,I,CALL) FIXED BIN(31,0);
DCL (N,R) CHAR(27);
DCL (NN,W,PRO) FIXED DECIMAL;

```

```

NCOM:  PROC(NGEN,W) FIXED BIN(31,0) RECURSIVE;
/* THIS PROCEDURE CALCULATES AN UPPER BOUND TO THE
   NUMBER OF BASIC COMMUTATORS FOR EACH WEIGHT UP TO
   A GIVEN WEIGHT FOR A GIVEN NUMBER OF GENERATORS
   NGEN.
   THE VALLES ARE ENTERED IN AN ARRAY NC */
DCL NC(10) FIXED BIN(31,0) STATIC EXT;
DCL (
    NGEN,W,NN,I
) FIXED BIN(31,0);
DCL (
    NCOM
) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
   RETURNS(FIXED BIN(31,0));
IF W>10 THEN SIGNAL CONDITION (ALPHA);
IF W>10 THEN DO; PUT LIST("DUMMY STATEMENT");
               SIGNAL CONDITION (ALPHA); END;
IF W=1 THEN DO; NC(1)=NGEN; RETURN(NGEN); END;
NN=0;
I=(W+1)/2; J=W-1;
DO I=I TO J;
    NN=NN+NCOM(NGEN,I)*NCOM(NGEN,W-I);
END;
NC(W)=NN; RETURN(NN);
END NCOM;

```

LT	35
NCR	3
NCR	4
NCR	5
NCOM	0
NCOM	1
NCOM	2
NCOM	3
NCOM	4
NCOM	5
NCOM	6
NCOM	7
NCOM	8
NCOM	9
NCOM	10
NCOM	11
NCOM	12
NCOM	13
NCOM	14
NCOM	15
NCOM	16
NCOM	17
NCOM	18
NCOM	19
NCOM	20
NCOM	21
NCOM	22
NCOM	23
NCOM	24

```

NCR:  PROC(N,R) FIXED BIN(31,0);
/* NCR RETURNS A STRUCTURE CONTAINING THE INTEGER EQUQL

```

NCR	0
NCR	1


```

/* TO THE NUMBER OF COMBINATIONS OF N THINGS R AT A TIME */
DCL (L,I,CAL) FIXED BIN(31,0);
DCL (N,R) CHAR(2);
DCL (DN,DR,PRD) FIXED DECIMAL;
DCL (
    LIST
    )ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
    NXTLFT
    )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
IF N=0 THEN DO;
    L=LIST (9); CAL=NXTLFT(UNSPEC('0100'),L);
    RETURN(L); END;
IF R=0 THEN DO;
    L=LIST(9); CAL=NXTLFT(UNSPEC('0100'),L);
    RETURN(L); END;
DN=N; DR=R; PRD=1;
DO I=1 TO DR; PRD=PRD*(DN-I+1); END;
DO I=1 TO DR; PRD=PRD/(DR-I+1); END;
IF PRD>99 THEN SIGNAL CONDITION(ALPHA);
L=LIST(9);
CAL=NXTLFT(UNSPEC('01' || SUBSTR(CHAR(PRD),7,2)),L);
RETURN(L);
END NCR;

```

NCR	2
NCR	3
NCR	4
NCR	5
NCR	6
NCR	7
NCR	8
NCR	9
NCR	10
NCR	11
NCR	12
NCR	13
NCR	14
NCR	15
NCR	16
NCR	17
NCR	18
NCR	19
NCR	20
NCR	21
NCR	22
NCR	23
NCR	24
NCR	25
NCR	26

```

NTERM: PROC(STR)          BIT(1);
/* RETURNS '1' IF STRUCTURE IS NON TERMINAL */
DCL (
    STR
    ) FIXED BIN(31,0);
DCL (

```

NTERM	0
NTERM	1
NTERM	2
NTERM	3
NTERM	4
NTERM	5

```
LNKL,LNKR,CONT
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
RETURN(LNKL(CONT(STR))=LNKR(CONT(STR)));
END NTERM;
```

```
NTERM 6
NTERM 7
NTERM 8
NTERM 9
```

```
NU: PROC(STR) FLOAT BIN(53) RECURSIVE;
DCL (
    WT,CONT,BOT
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
    MADNTP
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL (
    SYMB
)ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL NN FLOAT BIN(53);
DCL (
    L,R,WL,WR,WS,I,STR
)    FIXED BIN(31,0);
DCL NU ENTRY(FIXED BIN(31,0)) RETURNS(FLOAT BIN(53));
DCL NC(10) FIXED BIN(31,0) STATIC EXT;
DCL NA FIXED BIN(15,0);
DCL SYM CHAR(1);
DCL NTERM ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));
IF -NTERM(STR) THEN DO;
    UNSPEC(NA)=(10)'0'B||SUBSTR(UNSPEC(SYMB(STR)),3,6); NN=NA;
    RETURN(NN);    END;
L=CONT(MADNTP(STR,2)+1); R=BOT(STR);
SYM=SYMB(STR);
IF SYM=':' | SYM='''' THEN RETURN(NU(L));
WS=WT(STR); WR=WT(R); WL=WS-1;
```

```
NU 0
NU 1
NU 2
NU 3
NU 4
NU 5
NU 6
NU 7
NU 8
NU 9
NU 10
NU 11
NU 12
NU 13
NU 14
NU 15
NU 16
NU 17
NU 18
NU 19
NU 20
NU 21
NU 22
NU 23
NU 24
NU 25
NU 26
```

SECTION 13 - PROCEDURES OF THE COLLECTING PROCESS

40

```

I=(WS+1)/2; NN=C;
DO I=1 TO WL;
    NN=NN+NC(I)*NC(WS-I)+NC(WR)*NU(L)+NU(R);
END;
RETURN(NN);
END NU;

```

```

ORDPROD 27
ORDPROD 28
ORDPROD 29
ORDPROD 30
ORDPROD 31
ORDPROD 32
ORDPROD 33

```

ORDPROD: PROC(STR,MOD);

```

DCL (
    P1,P2,P3,P4,PH,W1,W2,W3,K1,K2,K3,KA,KB,WA,WB,ORLA,ORLB,
    WGT,ORDL,ORDR,ORL1,ORL2,ORL3,ORR1,ORR2,ORR3,ORRA,ORRB,
    L,MOD,LST,CAL,STR,LA,HOLDLST,P5,LR
) FIXED BIN(31,0);
DCL (
    NEW
) BIT(1);
DCL (
    LNKR,LNKL,CONT,DELETE,SUMEX,TCP,BOT,LIST,LSTSTR,MTLIST
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
    SUBST,MADNTP,MADNBT,NXTLFT,NXTRGT,INLSTR
) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL (
    CANONCL,PROSPEC,STRPN
) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0));
DCL (
    LG
) ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL PRNT FIXED BIN(31,0) EXT, PRNTI BIT(1) EXT;
DCL BETA FIXED BIN(31,0) EXT,BETS FIXED BIN(31,0);

```

```

ORDPROD 0
ORDPROD 1
ORDPROD 2
ORDPROD 3
ORDPROD 4
ORDPROD 5
ORDPROD 6
ORDPROD 7
ORDPROD 8
ORDPROD 9
ORDPROD 10
ORDPROD 11
ORDPROD 12
ORDPROD 13
ORDPROD 14
ORDPROD 15
ORDPROD 16
ORDPROD 17
ORDPROD 18
ORDPROD 19
ORDPROD 20
ORDPROD 21
ORDPROD 22
ORDPROD 23
ORDPROD 24

```

DCL IDNT ENTRY(FIXED BIN(31,0))RETURNS(BIT(1));	ORDPROD 25
DCL F BIT(1);	ORDPROD 26
BETS=BETA;	ORDPROD 27
IF PRNT1 THEN PUT SKIP(2) LIST('ENTER ORDPD');	ORDPROD 28
IF PRNT1 THEN CALL STRPN(STR,0);	ORDPROD 29
LST=LIST(9); CAL=PROSPEC(STR,LST);	ORDPROD 30
PH,P1,P2,P3,P4=LNKR(LST);	ORDPROD 31
WGT,ORDL,ORDR='0'B {(31)'1'B; NEW='1'B;	ORDPROD 32
F='0'B;	ORDPROD 33
	ORDPROD 34
	ORDPROD 35
L1: /* SCAN LIST FOR LEAST COMMUTATOR */	ORDPROD 36
/* P1 POINTS TO THE RIGHTMOST COLLECTED PART */	ORDPROD 37
/* P2 POINTS TO THE LEFTMOST &LEAST FACTOR	ORDPROD 38
CURRENTLY KNOWN DURING SCANNING */	ORDPROD 39
/* P3 IS USED FOR SCANNING */	ORDPROD 40
/* PH MARKS THE HEADER */	ORDPROD 41
	ORDPROD 42
P3=LNKR(CONT(P3)); K3=CONT(P3+1);	ORDPROD 43
IF P3=PH THEN GOTO L2;	ORDPROD 44
W3=BUT(K3); ORL3=TOP(K3); ORR3=CONT(MADNTP(K3,2)+1);	ORDPROD 45
IF W3 > WGT THEN GOTO L1;	ORDPROD 46
IF W3 < WGT THEN GOTO L7;	ORDPROD 47
IF ORL3 > ORDL THEN GOTO L1;	ORDPROD 48
IF ORL3 < ORDL THEN GOTO L7;	ORDPROD 49
IF ORR3 >= ORDR THEN GOTO L1;	ORDPROD 50
	ORDPROD 51
L7: P2=P3; NEW='0'B; K2=K3;	ORDPROD 52
WGT=W3; ORDL=ORL3; ORDR=ORR3; GOTO L1;	ORDPROD 53
	ORDPROD 54
L2: /* THIS SECTION MOVES THE LEAST COMMUTATOR FOUND	ORDPROD 55
IN THE LAST SECTION */	ORDPROD 56
IF NEW='1'B THEN GOTO L3;	ORDPROD 57
IF P2=LNKR(CONT(PH)) THEN DO; P1=P2; GOTO L3; END;	ORDPROD 58

IF P2=LNKR(CONT(P1)) THEN DO; P1=P2; GOTO L3; END;	ORDPROD 59
IF P3=LNKR(CONT(P1)) THEN DO;	ORDPROD 60
HOLDLST=LIST(9); CAL=NXTLFT(CONT(P1+1),HOLDLST);	ORDPROD 61
GOTO L4; END;	ORDPROD 62
IF P1=PH THEN DO;	ORDPROD 63
CAL=NXTRGT(K2,P1); P1=LNKR(CONT(P1)); P2=LNKR(CONT(P2));	ORDPROD 64
HOLDLST=LIST(9); CAL=NXTLFT(CONT(P1+1),HOLDLST);	ORDPROD 65
CAL=DELETE(LNKL(CONT(P2))); P3=P1; GOTO L4; END;	ORDPROD 66
K1=CONT(P1+1); W1=BOT(K1); ORL1=TOP(K1);	ORDPROD 67
ORR1=CONT(MADNTP(K1,2)+1);	ORDPROD 68
HOLDLST=LIST(9); CAL=NXTLFT(K2,HOLDLST);	ORDPROD 69
CAL=NXTRGT(K2,P1); P1=LNKR(CONT(P1));	ORDPROD 70
P3=P1; P2=LNKR(CONT(P2)); CAL=DELETE(LNKL(CONT(P2)));	ORDPROD 71
	ORDPROD 72
	ORDPROD 73
/* SCAN REGION BETWEEN P1 & P2 TO INTRODUCE EXTRA	ORDPROD 74
TERMS IF NECESSARY */	ORDPROD 75
	ORDPROD 76
L4: P3=LNKR(CONT(P3)); K3=CONT(P3+1);	ORDPROD 77
IF P3=PH THEN DO; CAL=IRALST(HOLDLST); GOTO L3; END;	ORDPROD 78
	ORDPROD 79
IF P3=P2 THEN DO;	ORDPROD 80
IF PRNT1 & F THEN DO;	ORDPROD 81
F='0'B; L=LSTSTR(LST);	ORDPROD 82
CALL STRPN(L,0); CAL=IRALST(L);	ORDPROD 83
END;	ORDPROD 84
P3=LNKL(CONT(P3));	ORDPROD 85
NEW='1'B; CAL=IRALST(HOLDLST);	ORDPROD 86
L8: P3=LNKR(CONT(P3)); K3=CONT(P3+1);	ORDPROD 87
IF P3=PH THEN GOTO L2;	ORDPROD 88
W3=BOT(K3); ORL3=TOP(K3); ORR3=CONT(MADNTP(K3,2)+1);	ORDPROD 89
IF W3>WGT THEN GOTO L8;	ORDPROD 90
IF ORL3>ORDL THEN GOTO L8;	ORDPROD 91
IF ORR3>ORDR THEN GOTO L8;	ORDPROD 92

P2=P3; NEW='0'B; K2=K3;	ORDPROD 93
IF P3=LNKR(CONT(P3)); K3=CONT(P3+1);	ORDPROD 94
CA GOTO L2; END;	ORDPROD 95
GOTO L3;	ORDPROD 96
J2: W3=BOT(K3); URL3=TOP(K3);	ORDPROD 97
URR3=CONT(MADNTP(K3,2)+1);	ORDPROD 98
K1=BOT(HOLDLST); W1=BOT(K1);	ORDPROD 99
IF W3+W1 < MOD THEN DO;	ORDPROD100
L=LG(' ', ' ', CONT(MADNBT(K3,2)+1), CONT(MADNBT(K1,2)+1));	ORDPROD101
BETA=L; CALL CANONCL(L,MOD);	ORDPROD102
P5=LNKR(CONT(P3));	ORDPROD103
LA=LIST(9); CAL=PROSPEC(L,LA); CAL=INLSTR(LA,P3);	ORDPROD104
P3=LNKL(CONT(P5));	ORDPROD105
F='1'B;	ORDPROD106
IF P3=PH THEN GOTO L4; END;	ORDPROD107
GOTO L4;	ORDPROD108
L3: IF P4=PH THEN P4=LNKR(CONT(P4));	ORDPROD109
P3=P4; KA=CONT(P3+1); WA=BOT(KA);	ORDPROD110
ORLA=TOP(KA); ORRA=CONT(MADNTP(KA,2)+1);	ORDPROD111
L5: P3=LNKR(CONT(P3)); IF P3=PH THEN GOTO L6;	ORDPROD112
KB=CONT(P3+1); WB=BOT(KB);	ORDPROD113
ORLB=TOP(KB); ORRB=CONT(MADNTP(KB,2)+1);	ORDPROD114
IF WA=WB THEN GOTO J2;	ORDPROD115
IF ORLA=ORLB THEN GOTO J2;	ORDPROD116
IF ORRA=ORRB THEN GOTO J2;	ORDPROD117
F='1'B;	ORDPROD118
L=LG(' ', ' ', CONT(MADNBT(KA,2)+1), CONT(MADNBT(KB,2)+1));	ORDPROD119
CAL=SUMEX(L);	ORDPROD120
BETA=BOT(L); CALL CANONCL(BOT(L),MOD);	ORDPROD121
BETA=L; CALL CANONCL(L,MOD);	ORDPROD122
LA=LIST(9); CAL=PROSPEC(L,LA); CAL=INLSTR(LA,P3);	ORDPROD123
IF P1=LNKL(CONT(P3)) THEN P1=P3;	ORDPROD124
CAL=DELETE(LNKL(CONT(P3)));	ORDPROD125
	ORDPROD126

SECTION 13 - PROCEDURES OF THE COLLECTING PROCESS

44

```

P3=LNKR(CONT(P3));
IF P1=LNKL(CONT(P3)) THEN P1=P3;
CAL=DELETE(LNKL(CONT(P3))); KA=CONT(P3+1);
GOTO L5;

```

```

J2: /* SUMMATION NOT POSSIBLE */
WA=WB; ORLA=ORLB; ORRA=ORRB; KA=KB;
GOTO L5;

```

```

L6: ORDL,ORDR,WGT='0'B||(31)'1'B; NEW='1'B;
IF PRNT1 & F THEN DO;
    F='0'B; L=LSTSTR(LST);
    CALL STRPN(L,0); CAL=IRALST(L);
    END;
P2,P3,P4=P1;
IF LNKR(CONT(P1))= PH THEN GOTO L1;

```

```

PREC: L=LSTSTR(LST); CAL=IRALST(LST);
/* THE P CAL =IRALST(INLSTR(L,MTLIST(STR)));
THE PROD IF PRNT1 THEN CALL STRPN(STR,0);
PRECEDEN BETA=BETS;
THE STR IF PRNT1 THEN PUT SKIP(2) LIST('EXIT ORDPROD');
IT RETUR END ORDPROD;

```

```

PERMCOM: PROC(STR);
/* A,B -> (B,A) */
DCL (
    L,A,B,STR,CAL
) FIXED BIN(31,0);
DCL (
    CONT,BOT,MTLIST
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));

```

```

ORDPROD127
ORDPROD128
ORDPROD129
ORDPROD130
ORDPROD131
ORDPROD132
ORDPROD133
ORDPROD134
ORDPROD135
ORDPROD136
ORDPROD137
ORDPROD138
ORDPROD139
ORDPROD140
ORDPROD141
ORDPROD142
ORDPROD143
ORDPROD144
ORDPROD145
ORDPROD146
ORDPROD147
ORDPROD148
ORDPROD149

```

```

PERMCOM 0
PERMCOM 1
PERMCOM 2
PERMCOM 3
PERMCOM 4
PERMCOM 5
PERMCOM 6
PERMCOM 7

```

DCL (PERMCOM 8
MADNTP, INLSTR	PERMCOM 9
)ENTRY(FIXED BIN(31,0), FIXED BIN(31,0))	PERMCOM 10
RETURNS(FIXED BIN(31,0));	PERMCOM 11
DCL (PERMCOM 12
LG	PERMCOM 13
)ENTRY(CHAR(1), FIXED BIN(31,0), FIXED BIN(31,0))	PERMCOM 14
RETURNS(FIXED BIN(31,0));	PERMCOM 15
DCL IRALST ENTRY(FIXED BIN(31,0));	PERMCOM 16
A=CONT(MADNTP(STR,2)+1);	PERMCOM 17
B=BOT(STR);	PERMCOM 18
A= (LG('','','', LG('','',B,A),0));	PERMCOM 19
CALL IRALST(INLSTR(A, MTLIST(STR)));	PERMCOM 20
END PERMCOM;	PERMCOM 21

PREC: PROC(STR) FIXED BIN(31,0);	PREC 0
/* THE PARAM STR IS A STRUCTURE.	PREC 1
THE PROC RETURNS A VALUE WHICH INDICATES THE ORDER OF	PREC 2
PRECEDENCE OF AN OPERATOR. THE LARGER THE VALUE	PREC 3
THE STRONGER THE OPERATOR BOND.	PREC 4
IT RETURNS THE VALUE 20 IF THE STR IS TERMINAL */	PREC 5
DCL (STR,1) FIXED BIN(31,0);	PREC 6
DCL S CHAR(1);	PREC 7
DCL (PREC 8
SYMB	PREC 9
)ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));	PREC 10
DCL (PREC 11
NTERM	PREC 12
)ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));	PREC 13
DCL CH(10) CHAR(1) STATIC	PREC 14
INITIAL('U', '-','+', '*', ' ', '/', '.', ':');	PREC 15
DCL PR(10) FIXED BIN(31,0) STATIC	PREC 16

SECTION 13 - PROCEDURES OF THE COLLECTING PROCESS

46

INITIAL(10,1,1,2,3,4,5,6,10,7);	PREC	17
IF -NTERM(STR) THEN RETURN(20);	PREC	18
S=SYMB(STR);	PREC	19
DO I=1 TO 10; IF CH(I)=S THEN RETURN(PR(I)); END;	PREC	20
PUT LIST('INVALID OPERATOR'); RETURN(20);	PREC	21
END PREC;	PREC	22

PROD:	PROC(A,B) FIXED BIN(31,0);	PROD	0
	/* ADDS TWO INTEGERS IN LIST FORMAT AND RETURNS THE SUM	PROD	1
	IN LIST FORMAT */	PROD	2
	DCL (AM,BM) BIT(1);	PROD	3
	DCL (DL,DR) FIXED DEC(5,0);	PROD	4
	DCL (A,B,L,R,CAL) FIXED BIN(31,0);	PROD	5
	DCL CT CHAR(2);	PROD	6
	DCL (PROD	7
	LIST,BUT	PROD	8
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	PROD	9
	DCL (PROD	10
	SYMB	PROD	11
)ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));	PROD	12
	DCL (PROD	13
	NTERM	PROD	14
)ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));	PROD	15
	DCL (PROD	16
	NXTLEFT	PROD	17
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	PROD	18
	RETURNS(FIXED BIN(31,0));	PROD	19
	DCL (PROD	20
	LG	PROD	21
)ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))	PROD	22
	RETURNS(FIXED BIN(31,0));	PROD	23
	AM='0'B; BM='0'B;	PROD	24

IF NTERM(A) THEN IF SYMB(A)='U' THEN	PROD	25
DO; AM='1'B; L=BCT(BOT(A)); END;	PROD	26
ELSE SIGNAL CONDITION(ALPHA);	PROD	27
ELSE L=BOT(A);	PROD	28
IF NTERM(B) THEN IF SYMB(B)='U' THEN	PROD	29
DO; BM='1'B; R=BCT(BOT(B)); END;	PROD	30
ELSE SIGNAL CONDITION(ALPHA);	PROD	31
ELSE R=BOT(B);	PROD	32
UNSPEC(CT)=SUBSTR(UNSPEC(L),17,16);	PROD	33
DL=CT; IF AM THEN DL=-DL;	PROD	34
UNSPEC(CT)=SUBSTR(UNSPEC(R),17,16);	PROD	35
DR=CT; IF BM THEN DR=-DR;	PROD	36
DL=DL*DR;	PROD	37
IF DL>99 THEN SIGNAL CONDITION(ALPHA);	PROD	38
IF DL=0 THEN DO;	PROD	39
R=LIST(9); CAL=NXTLEFT(UNSPEC('0100'),R);	PROD	40
RETURN(R); END;	PROD	41
IF DL<0 THEN DO;	PROD	42
DL=-DL;	PROD	43
R=LIST(9); CAL=NXTLEFT(UNSPEC('01') SUBSTR(CHAR(DL),7,2),R);	PROD	44
R=LG('U',R,0);	PROD	45
RETURN(R); END;	PROD	46
IF DL>0 THEN DO;	PROD	47
R=LIST(9); CAL=NXTLEFT(UNSPEC('01') SUBSTR(CHAR(DL),7,2),R);	PROD	48
RETURN(R);	PROD	49
SIGNAL CONDITION (ALPHA);	PROD	50
END PROD;	PROD	51

PRODEXP: PROC(STR) FIXED BIN(31,0);	PRODEXP	0
/* (A:M):N -> A:(MN) */	PRODEXP	1
/* UNARY MINUS IF PRESENT IS RAISED TO THE TOP OF	PRODEXP	2
THE STRUCTURE */	PRODEXP	3

SECTION 13 - PROCEDURES OF THE COLLECTING PROCESS

48

```

/* STR IS DCL (
    STR,A,B,M,N,L,CAL
    ) ENTRY(FIXED BIN(31,0));
DCL (
    TOP,BOT,CONT,MTLIST
    ) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
    LG
    ) ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL (
    MADNTP,INLSTR
    ) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL (
    LU,RU
    ) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL IRALST ENTRY(FIXED BIN(31,0));
L=CONT(MADNTP(STR,2)+1);
A=CONT(MADNTP(L,2)+1);
M=BOT(L);
N=BOT(STR);
LU='0'B; RU='0'B;
IF SUBSTR(UNSPEC(TOP(M)),9,24)=UNSPEC('COU') THEN LU='1'B;
IF SUBSTR(UNSPEC(TOP(N)),9,24)=UNSPEC('COU') THEN RU='1'B;
IF ~LU & ~RU THEN B=LG('*',M,N);
IF LU & RU THEN B=LG('*',BOT(M),BOT(N));
IF LU & ~RU THEN B=LG('U',LG('*',BOT(M),N),0);
IF ~LU & RU THEN B=LG('U',LG('*',M,BOT(N)),0);
CAL= IRALST(INLSTR(LG(':',A,B),MTLIST(STR)));
END PRODEXP;

```

```

PRODEXP 4
PRODEXP 5
PRODEXP 6
PRODEXP 7
PRODEXP 8
PRODEXP 9
PRODEXP 10
PRODEXP 11
PRODEXP 12
PRODEXP 13
PRODEXP 14
PRODEXP 15
PRODEXP 16
PRODEXP 17
PRODEXP 18
PRODEXP 19
PRODEXP 20
PRODEXP 21
PRODEXP 22
PRODEXP 23
PRODEXP 24
PRODEXP 25
PRODEXP 26
PRODEXP 27
PRODEXP 28
PRODEXP 29
PRODEXP 30
PRODEXP 31
PRODEXP 32
PRODEXP 33
PRODEXP 34

```

/* STR IS A STRUCTURE WITH OPERATOR '*'. IT IS ASSUMED THAT THE
LEFT AND RIGHT PARTS ARE IN CANONIC FORM AND IT
CONVERTS THE STRUCTURE TO CANONIC FORM. */

```

DCL(
  L,R,BR,CAL,STR
)FIXED BIN(31,0);
DCL(
  TL,TR
)CHAR(4);
DCL(
  SUBST,MADNTP,INLSTR,PROD,SUBSTP,NXTLFT
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
DCL(
  LG
)ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
DCL(
  SYMB
)ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL(
  CONT,MTLIST,TOP,BOT,LIST
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
  NTERM
)ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));
DCL(WL,WR) CHAR(1);
DCL (LT,RT) CHAR(3);
DCL (WRA,WRB) CHAR(1);
L=CONT(MADNTP(STR,2)+1); R=BOT(STR);
UNSPEC(TL)=UNSPEC(TOP(L)); UNSPEC(TR)=UNSPEC(TOP(R));
WL=SUBSTR(TL,2,1); WR=SUBSTR(TR,2,1);
LT=SUBSTR(TL,2,3); RT=SUBSTR(TR,2,3);
IF WL='I' & WR='I' THEN

```

PRODINT 1
 PRODINT 2
 PRODINT 3
 PRODINT 4
 PRODINT 5
 PRODINT 6
 PRODINT 7
 PRODINT 8
 PRODINT 9
 PRODINT 10
 PRODINT 11
 PRODINT 12
 PRODINT 13
 PRODINT 14
 PRODINT 15
 PRODINT 16
 PRODINT 17
 PRODINT 18
 PRODINT 19
 PRODINT 20
 PRODINT 21
 PRODINT 22
 PRODINT 23
 PRODINT 24
 PRODINT 25
 PRODINT 26
 PRODINT 27
 PRODINT 28
 PRODINT 29
 PRODINT 30
 PRODINT 31
 PRODINT 32
 PRODINT 33
 PRODINT 34


```

DO; L=PROD(L,R); GOTO RET; END;
IF LT='000' & WR='1' THEN DO;
UNSPEC(WRA)=SUBSTR(UNSPEC(TOP(BOT(L))),9,8);
IF WRA='1' THEN DO; L=PROD(L,R); GOTO RET; END;
END;
IF WL='1' & RT='000' THEN DO;
UNSPEC(WRB)=SUBSTR(UNSPEC(TOP(BOT(R))),9,8);
IF WRB='1' THEN DO; L=PROD(L,R); GOTO RET; END;
END;
IF LT='000' & RT='000' THEN DO;
UNSPEC(WRA)=SUBSTR(UNSPEC(TOP(BOT(L))),9,8);
UNSPEC(WRB)=SUBSTR(UNSPEC(TOP(BOT(R))),9,8);
IF WRA='1' & WRB='1' THEN DO;
L=PROD(L,R); GOTO RET; END;
END;
IF -NTERM(L) & -NTERM(R) THEN RETURN;
IF LT='100' | LT='1 0' | RT='100' | RT='1 0' THEN DO;
L=LIST(9); CAL=NXTLEFT(UNSPEC('0100'),L); GOTO RET; END;
IF SUBSTR(TR,2,1)='1' THEN DO; L=LG('*',R,L); GOTO RET; END;
IF SUBSTR(TR,2,3)='000' THEN DO;
/C.SYSIN DD *
END PRODINT;

```

```

PRODINT 35
PRODINT 36
PRODINT 37
PRODINT 38
PRODINT 39
PRODINT 40
PRODINT 41
PRODINT 42
PRODINT 43
PRODINT 44
PRODINT 45
PRODINT 46
PRODINT 47
PRODINT 48
PRODINT 49
PRODINT 50
PRODINT 51
PRODINT 52
PRODINT 53
PRODINT 54
PRODINT 55
PRODINT 56

```

```

PROSPEC: PROC(STR,LST) RECURSIVE;
/* THIS PROCEDURE TRANSFORMS THEN BINARY STRUCTURE STR
INTO A LINEAR LIST OF LISTS. THE LATTER IS ATTACHED
TO THE EMPTY LIST LST. */
DCL (
CAL,NCL,NLR,STR,LST,LA,L,R
) FIXED BIN(31,0);
DCL NUA FLOAT BIN(53);
DCL (

```

```

PROSPEC 0
PROSPEC 1
PROSPEC 2
PROSPEC 3
PROSPEC 4
PROSPEC 5
PROSPEC 6
PROSPEC 7
PROSPEC 8

```

SECTION 13 - PROCEDURES OF THE COLLECTING PROCESS

51

```

DCL SYMB
  )ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL (
  NXTLFT,MADNTP,PROSPEC
  )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
DCL (
  CONT,BOT,LIST,WT
  )ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL NU ENTRY(FIXED BIN(31,0)) RETURNS(FLCAT BIN(53));
DCL IDNT ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));
IF IDNT(STR) THEN RETURN;
IF SYMB(STR) <= '.' THEN DO;
  LA=LIST(9); NUA=NU(STR);
  NUL=NUA/ 1E31B; NUR=NUA-NUL* 1E31B;
  CAL=NXTLFT(NUL,LA); CAL=NXTLFT(NUR,LA);
  CAL=NXTLFT(STR,LA); CAL=NXTLFT(WT(STR),LA);
  CAL=NXTLFT(LA,LST); RETURN;
END;
L=CONT(MADNTP(STR,2)+1); R=BOT(STR);
CALL PROSPEC(L,LST); CALL PROSPEC(R,LST);
END PROSPEC;

```

```

PROSPEC 9
PROSPEC 10
PROSPEC 11
PROSPEC 12
PROSPEC 13
PROSPEC 14
PROSPEC 15
PROSPEC 16
PROSPEC 17
PROSPEC 18
PROSPEC 19
PROSPEC 20
PROSPEC 21
PROSPEC 22
PROSPEC 23
PROSPEC 24
PROSPEC 25
PROSPEC 26
PROSPEC 27
PROSPEC 28
PROSPEC 29
PROSPEC 30

```

```

REQ: PROC(STR,MOD);
/* RIGHT EXPONENT OUT */
/* N IS THE GREATEST INTEGER SUCH THAT
   WT(A)+N.WT(B) < MOD */
DCL K FIXED BIN(31,0);
DCL(
  STR,MOD,A,B,X,BX,N,L,LA,LB,J,CAL,MTA,WTB,BETS,
  BAB,BABB,BA,LC,L1,L2,L3,L4
  )FIXED BIN(31,0);

```

```

REQ 0
REQ 1
REQ 2
REQ 3
REQ 4
REQ 5
REQ 6
REQ 7
REQ 8

```

DCL(END;	REQ	9
IF WT,CONT,BOT,LIST,MTLIST,TOP		REQ	10
ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));		REQ	11
DCL(REQ	12
NXTLEFT,MADNTP,INLSTR		REQ	13
ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))		REQ	14
RETURNS(FIXED BIN(31,0));		REQ	15
DCL(REQ	16
LG		REQ	17
ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))		REQ	18
RETURNS(FIXED BIN(31,0));		REQ	19
DCL INT BIT(1);		REQ	20
DCL KC CHAR(10), (KT,NX) CHAR(2);		REQ	21
DCL PRNTI BIT(1) EXT;		REQ	22
DCL(REQ	23
STRPN		REQ	24
ENTRY(FIXED BIN(31,0),FIXED BIN(31,0));		REQ	25
DCL(REQ	26
CANONCL		REQ	27
ENTRY(FIXED BIN(31,0),FIXED BIN(31,0));		REQ	28
DCL BETA FIXED BIN(31,0) EXT;		REQ	29
DCL SCNP ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),		REQ	30
FIXED BIN(31,0),FIXED BIN(31,0));		REQ	31
DCL NCR ENTRY(CHAR(2),CHAR(2))RETURNS(FIXED BIN(31,0));		REQ	32
IF PRNTI THEN PUT SKIP(2) LIST('ENTER REQ');		REQ	33
IF PRNTI THEN CALL STRPN(STR,0);		REQ	34
BETS=BETA; BETA=L; INT='0'B;		REQ	35
A=CONT(MADNTP(STR,2)+1); BX=BOT(STR);		REQ	36
B=CONT(MADNTP(BX,2)+1); X=BOT(BX);		REQ	37
LA=LIST(9); CAL=NXTLEFT(UNSPEC('0000'),LA);		REQ	38
LB=TOP(X);		REQ	39
IF SUBSTR(UNSPEC(LB),9,8)=UNSPEC('1') THEN DO;		REQ	40
UNSPEC(NX)=SUBSTR(UNSPEC(LB),17,16); INT='1'B;		REQ	41
IF SUBSTR(NX,1,1)='0' THEN SUBSTR(NX,1,1)=' ';		REQ	42

	END;	REO	43
IF INT THEN IF NX<'1' THEN GOTO S2;		REO	44
WTA=WT(A); WTB=WT(B); IF WTA+WTB>= MOD THEN GOTO S2;		REO	45
IF MOD>7 THEN SIGNAL CONDITION(ALPHA);		REO	46
CAL=IRALST(LA);		REO	47
L=LG(' ', ' ', B, A);		REO	48
IF PRNTI THEN CALL STRPN(L, 0);		REO	49
BETA=L; CALL CANONCL(L, MOD);		REO	50
BA, LA=L;		REO	51
LC=LG('U', x, 0);		REO	52
BETA=LC; CALL CANONCL(LC, MOD);		REO	53
LA=LG(' ', ' ', L, LC);		REO	54
IF PRNTI THEN CALL STRPN(LA, 0);		REO	55
BETA=LA; CALL CANONCL(LA, MOD);		REO	56
N=(MOD-1-WTA)/WTB; KC=(9)'0' '1';		REO	57
BABB=0;		REO	58
DO J=2 TO N;		REO	59
KC=KC+'1'; KT=SUBSTR(KC, 8, 2);		REO	60
IF INT THEN IF NX<KT THEN GOTO S1;		REO	61
L=LG(' ', ' ', L, B);		REO	62
IF PRNTI THEN CALL STRPN(L, 0);		REO	63
BETA=L; CALL CANONCL(L, MOD);		REO	64
IF J=2 THEN BAB=L; IF J=3 THEN BABB=L;		REO	65
IF INT THEN IF NX=KT THEN DO;		REO	66
LB=LIST(9); CAL=NXTLEFT(UNSPEC('0101'), LB);		REO	67
LB=LG(' ', ' ', L, LG('U', LB, 0));		REO	68
IF PRNTI THEN CALL STRPN(LB, 0);		REO	69
BETA=LB; CALL CANONCL(LB, MOD);		REO	70
LA=LG(' ', ' ', LA, LB);		REO	71
IF PRNTI THEN CALL STRPN(LA, 0);		REO	72
GOTO S1; END;		REO	73
IF INT THEN LB=NCR(NX, KT);		REO	74
ELSE DO; LB=LIST(9); CAL=NXTLEFT(UNSPEC('01' KT), LB);		REO	75
LB=LG(' ', ' ', x, LB);		REO	76


```

        END;
        LB=LG(':',L,LG('U',LB,0));
        IF PRNT1 THEN CALL STRPN(LB,0);
        BETA=LB; CALL CANONCL(LB,MOD);
        LA=LG('.',LA,LB);
END;

```

```

S1:  IF 2*WTA+3*WTB < MOD THEN DO;
        L=LG(' ',BAB,BA);
        BETA=L; CALL CANONCL(L,MOD);
        L1=LIST(9); CAL=NXTLFT(UNSPEC('0101'),L1);
        L3=LIST(9); CAL=NXTLFT(UNSPEC('0103'),L3);
        IF INT THEN DO;
            KT=SUBSTR(CHAR(NX+'01'),8,2);
            LC=NCR(KT,'03');
            BETA=LC; CALL CANONCL(LC,MOD);
            END;
        ELSE DO;
            LC=LG(' ',LG(' ',X,L1),L3);
            END;
        L=LG(':',L,LC);
        BETA=L; CALL CANONCL(L,MOD);
        LA=LG('.',LA,L);

```

```

        END;
    IF 2*WTA+4*WTB < MOD THEN DO;
        IF BABB=0 THEN GOTO S2;
        L=LG(' ',BABB,BA);
        BETA=L; CALL CANONCL(L,MOD);
        L2=LIST(9); CAL=NXTLFT(UNSPEC('0102'),L2);
        L4=LIST(9); CAL=NXTLFT(UNSPEC('0104'),L4);
        IF INT THEN DO;
            LC=NCR(KT,'04');
            BETA=LC; CALL CANONCL(LC,MOD);
            END;

```

```

REO      77
REO      78
REO      79
REO      80
REO      81
REO      82
REO      83
REO      84
REO      85
REO      86
REO      87
REO      88
REO      89
REO      90
REO      91
REO      92
REO      93
REO      94
REO      95
REO      96
REO      97
REO      98
REO      99
REO     100
REO     101
REO     102
REO     103
REO     104
REO     105
REO     106
REO     107
REO     108
REO     109
REO     110

```

```

A=LG(0:0, L, LC);
LC=LG(0:0, LG(0:0, X, L1), L4);
END;
L=LG(0:0, L, LC);
BETA=L; CALL CANONCL(L, MOD);
LA=LG(0:0, LA, L);
END;

```

```

S2: BETA=LA;
CALL SCNP(LA, LA, 0, MOD);
IF PRNT1 THEN CALL STRPN(LA, 0);
CAL=IRALST(INLSTR(LA, MTLIST(STR)));
BETA=BETS;
IF PRNT1 THEN PUT SKIP(2) LIST('EXIT REC');
END REC;

```

```

RINVOUT: PROC(STR, MOD);
/* (A,B^0) -> (A,B)*(B,A,B^0) */
DCL (
    A,B,L,CAL,STR,MOD
) FIXED BIN(31,0);
DCL (
    LG
) ENTRY(CHAR(1), FIXED BIN(31,0), FIXED BIN(31,0))
RETURNS(FIXED BIN(31,0));
DCL (
    BOT,CON,WT,MTLIST
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
    MADNTP,INLSTR
) ENTRY(FIXED BIN(31,0), FIXED BIN(31,0))
RETURNS(FIXED BIN(31,0));

```

```

REC OUT 111
REC OUT 112
REC OUT 113
REC OUT 114
REC OUT 115
REC OUT 116
REC OUT 117
REC OUT 118
REC OUT 119
REC OUT 120
REC OUT 121
REC OUT 122
REC OUT 123
REC OUT 124
REC OUT 125

```

```

RINVOUT 0
RINVOUT 1
RINVOUT 2
RINVOUT 3
RINVOUT 4
RINVOUT 5
RINVOUT 6
RINVOUT 7
RINVOUT 8
RINVOUT 9
RINVOUT 10
RINVOUT 11
RINVOUT 12
RINVOUT 13
RINVOUT 14
RINVOUT 15

```

```

A=CONT(MADNTP(STR,2)+1);
B=BOT(BOT(STR));
L=LG(' ',LG(' ',A,B),0);
IF WT(A)+2*WT(B)<MOD THEN
    L=LG(' ',L,LG(' ',LG(' ',B,A),LG(' ',B,0)));
CALL IRALST(INLSTR(L,MTLIST(STR)));
END RINVOUT;

```

```

RINVOUT 16
RINVOUT 17
RINVOUT 18
RINVOUT 19
RINVOUT 20
RINVOUT 21
RINVOUT 22

```

```

RPRDOUT: PROC(STR,MOD);
/* A,BC -> (A,C)(A,B)(A,B,C) */
DCL (
    A,B,C,STR,L,CAL,MOD,WA,WB,WC,BETS,LA
) FIXED BIN(31,0);
DCL (
    LG
) ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
RETURNS(FIXED BIN(31,0));
DCL (
    BOT,CONT,WT,LIST,MLIST
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
    MADNTP,NXTLEFT,INLSTR
) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
RETURNS(FIXED BIN(31,0));
DCL (STRPN,CANONCL) ENTRY (FIXED BIN(31,0),FIXED BIN(31,0));
DCL PRNT1 BIT(1) EXT,BETA FIXED BIN(31,0)EXT;
IF PRNT1 THEN PUT SKIP(2) LIST('ENTER RPRDOUT');
IF PRNT1 THEN CALL STRPN(STR,0);
BETS=BETA;
A=CONT(MADNTP(STR,2)+1);
B=CONT(MADNTP(BOT(STR),2)+1);
C=BOT(BOT(STR));

```

```

RPRDOUT 0
RPRDOUT 1
RPRDOUT 2
RPRDOUT 3
RPRDOUT 4
RPRDOUT 5
RPRDOUT 6
RPRDOUT 7
RPRDOUT 8
RPRDOUT 9
RPRDOUT 10
RPRDOUT 11
RPRDOUT 12
RPRDOUT 13
RPRDOUT 14
RPRDOUT 15
RPRDOUT 16
RPRDOUT 17
RPRDOUT 18
RPRDOUT 19
RPRDOUT 20
RPRDOUT 21
RPRDOUT 22
RPRDOUT 23

```

```

WA=WT(A); WB=WT(B); WC=WT(C);
IF WA+WB >= MOD THEN DO;
    L=LIST(9); CAL=NXTLEFT(UNSPEC('OUOO'),L); END;
IF WA+WB < MOD THEN DO;
    L=LG(' ', ' ', A, B);
    IF PRNTI THEN CALL STRPN(L,0);
    BETA=L; CALL CANONCL(L,MOD); END;
IF WA+WB+WC < MOD THEN DO;
    LA=LG(' ', ' ', L, C);
    IF PRNTI THEN CALL STRPN(LA,0);
    BETA=LA; CALL CANONCL(LA,MOD);
    L=LG(' ', ' ', L, LA); END;
IF WA+WC < MOD THEN DO;
    LA=LG(' ', ' ', A, C);
    IF PRNTI THEN CALL STRPN(LA,0);
    BETA=LA; CALL CANONCL(LA,MOD);
    L=LG(' ', ' ', LA, L);
    IF PRNTI THEN CALL STRPN(L,0);
    BETA=L; CALL CANONCL(L,MOD); END;
CALL IRALST(INLSTR(L,MTLIST(STR)));
IF PRNTI THEN PUT SKIP(2) LIST('EXIT RPRDOUT');
BETA=BETS;
END RPRDOUT;

```

```

SCAN: PROC(STR,STROLD,POS,MOD) RECURSIVE;
/* POS=1 -> STR LEFT OF STROLD */
/* POS=2 -> STR RIGHT OF STROLD */
/* POS=0 -> STR IS MAIN STRUCTURE */
/* STR IS STRUCTURE TO BE SCANNED */
/* STROLD IS STRUCTURE FROM WHICH STR IS DERIVED */
DCL (
    MOD, STR, STROLD, POS, L, R, CAL, MODL, MODR, N

```

```

RPRDOUT 24
RPRDOUT 25
RPRDOUT 26
RPRDOUT 27
RPRDOUT 28
RPRDOUT 29
RPRDOUT 30
RPRDOUT 31
RPRDOUT 32
RPRDOUT 33
RPRDOUT 34
RPRDOUT 35
RPRDOUT 36
RPRDOUT 37
RPRDOUT 38
RPRDOUT 39
RPRDOUT 40
RPRDOUT 41
RPRDOUT 42
RPRDOUT 43
RPRDOUT 44
RPRDOUT 45
RPRDOUT 46

```

```

SCAN 0
SCAN 1
SCAN 2
SCAN 3
SCAN 4
SCAN 5
SCAN 6
SCAN 7

```


)	FIXED BIN(31,0);	SCAN	8
DCL (SCAN	9
SYM		SCAN	10
STRGN: PROC)	CHAR(1);	SCAN	11
DCL (SCAN	12
NTERM		SCAN	13
)	ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));	SCAN	14
DCL (SCAN	15
SYMB		SCAN	16
)	ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));	SCAN	17
DCL (SCAN	18
BOT,CONT,WT		SCAN	19
)	ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SCAN	20
DCL (SCAN	21
CANONCL,STRPN		SCAN	22
)	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0));	SCAN	23
DCL (SCAN	24
MADNTP,SUBST,SUBSBT		SCAN	25
)	ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	SCAN	26
)	RETURNS(FIXED BIN(31,0));	SCAN	27
DCL	SCAN ENTRY(FIXED BIN(31,0),FIXED BIN(31,0),	SCAN	28
)	FIXED BIN(31,0),FIXED BIN(31,0));	SCAN	29
IF	NTERM(STR) THEN RETURN;	SCAN	30
MODL=MOD;	MODR=MOD;	SCAN	31
L=CONT(MADNTP(STR,2)+1);	R=BCT(STR); SYM=SYMB(STR);	SCAN	32
IF SYM=' ' SYM='0'	THEN GOTO L1;	SCAN	33
IF SYM=','	THEN	SCAN	34
DO;	MODL=MOD-WT(R); MODR=MOD-WT(L); END;	SCAN	35
IF NTERM(L)	THEN CALL SCAN(L,STR,1,MODL);	SCAN	36
L1:	IF NTERM(R) THEN CALL SCAN(R,STR,2,MODR);	SCAN	37
IF SYMB(STR)='.'	& SYMB(STRCLD)='.'	SCAN	38
& STR=STRCLD	THEN RETURN;	SCAN	39
L2:	CALL CANONCL(STR,MOD);	SCAN	40
RETURN;		SCAN	41
END	SCAN;	SCAN	41

```

IF SUBSTR(UNSPEC(STRNG),1,25,61) = UNSPEC(1) &
SUBSTR(UNSPEC(STRNG),1,25,61) = UNSPEC(1) THEN
DCL CAL=PORT(STRNG); CAL=PORT(STRNG);
STRGN: PROC(STRNG) FIXED BIN(31,0) RECURSIVE;
/* AT THE FIRST LEVEL OF CALL THIS PROCEDURE TAKES A
STRING LIST STRNG, DIVIDES IT IN TWO PARTS SEPARATED BY
AN APPROPRIATE OPERATOR. IT THEN GENERATES A STRUCTURE
BASED ON THESE TWO PARTS AND THE DIVIDING OPERATOR. */
DCL (
F,S,SA,SB,SC,SD,BC,LR,PNT,STRNG,FA,CAL,P,N,PC,PNC
) FIXED BIN(31,0);
DCL (
NULSTR,SEQLR
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
RETURNS(FIXED BIN(31,0));
DCL (
CONT,
SEQRDR,STRGN,LNKL,LNKR,POPTOP,POPBOT,DELETE,TOP,BOT
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
CH,CHC,C,CC
) CHAR(1);
DCL (
LG
)ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
RETURNS(FIXED BIN(31,0));
DCL PREC ENTRY(CHAR(1)) RETURNS(FIXED BIN(31,0));
DCL (
NTERM
)ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));

IF -NTERM(STRNG) THEN RETURN(TOP(STRNG));

```

```

STRGN 31
STRGN 32
STRGN 33
STRGN 0
STRGN 1
STRGN 2
STRGN 3
STRGN 4
STRGN 5
STRGN 6
STRGN 7
STRGN 8
STRGN 9
STRGN 10
STRGN 11
STRGN 12
STRGN 13
STRGN 14
STRGN 15
STRGN 16
STRGN 17
STRGN 18
STRGN 19
STRGN 20
STRGN 21
STRGN 22
STRGN 23
STRGN 24
STRGN 25
STRGN 26
STRGN 27
STRGN 28
STRGN 29
STRGN 30

```

```

IF SUBSTR(UNSPEC(BOT(STRNG)),25,8) = UNSPEC('') &
END SUBSTR(UNSPEC(TOP(STRNG)),25,8) = UNSPEC('') THEN
DO; CAL=POPTOP(STRNG); CAL=PCPBOT(STRNG);
RETURN(STRGN(STRNG));
END;

```

```

STRNG: BC=0; PC=10; CC='0';

```

```

S=SEQRDR(STRNG);

```

```

SCN: PNT=LNKR(S); SA=SEQLR(S,FA);

```

```

IF FA>0 THEN DO;

```

```

LR=NULSTR(PNC,STRNG); CAL=POPTOP(LR);

```

```

/* ERASE READER */

```

```

RETURN(LG(CHC,STRGN(STRNG),STRGN(LR)));

```

```

END;

```

```

IF FA=0 THEN GOTO SCN;

```

```

UNSPEC(CH)=SUBSTR(UNSPEC(SA),25,8);

```

```

IF CH='(' THEN DO; BC=BC+1; GOTO SCN; END;

```

```

IF CH=')' THEN DO; BC=BC-1; GOTO SCN; END;

```

```

IF BC<0 THEN SIGNAL CONDITION(ALPHA);

```

```

IF BC>0 THEN GOTO SCN;

```

```

P=PREC(CH); UNSPEC(C)=SUBSTR(UNSPEC(SA),17,8);

```

```

IF CH=',' & P=PC THEN

```

```

IF C>=CC THEN

```

```

DO; CC=C; PNC=PNT; CHC=CH; GOTO SCN; END;

```

```

ELSE GOTO SCN;

```

```

IF P<= PC THEN

```

```

DO; PC=P; PNC=PNT; CHC=CH; CC=C; GOTO SCN; END;

```

```

GOTO SCN;

```

```

PREC: PRDC(S) FIXED BIN(31,0);

```

```

DCL I FIXED BIN(31,0);

```

```

DCL S CHAR(1);

```

```

DCL CH(9) CHAR(1) STATIC

```

```

INITIAL('U','-',',','+','|','/',' ',' ',' ',' ',' ',' ');

```

```

DO I=1 TO 9; IF CH(I)=S THEN RETURN(I); END;

```

```
STRGN 31
```

```
STRGN 32
```

```
STRGN 33
```

```
STRGN 34
```

```
STRGN 35
```

```
STRGN 36
```

```
STRGN 37
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```
STRGN 38
```

```
STRGN 39
```

```
STRGN 40
```

```
STRGN 41
```

```
STRGN 42
```

```
STRGN 43
```

```
STRGN 44
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```
STRGN 45
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```
STRGN 46
```

```
STRGN 47
```

```
STRGN 48
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```
STRGN 49
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STRGN 50
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STRGN 51
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```
STRGN 52
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STRGN 53
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STRGN 54
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STRGN 55
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STRGN 56
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STRGN 57
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STRGN 58
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STRGN 59
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```
STRGN 60
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```
STRGN 61
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```
STRGN 62
```

```
STRGN 63
```

```
STRGN 64
```

```

PUT LIST('INVALID OPERATOR'); RETURN(20);
END PREC;
END STRGN; CH=1; GOTO OPNTR; END;

STRNGET: PROC(LMAIN);
/* THIS ROUTINE TAKES COMMUTATOR EXPRESSION FROMCARD
AND GENERATES THE CORRESPONDING LIST STRUCTURE.
BLANKS MAY BE USED FOR SPACING ARBITRARILY.
THE EXPRESSION MUST BE TERMINATED BY '#'.
THE LETTERS A -> L ARE ASSUMED TO BE GROUP ELEMENTS.
THE LETTERS M -> Z ARE ASSUMED TO BE INTEGER LITERALS */
DCL (
    GPL(12), INL(14), C, CH, CHA, CA
) CHAR(1);
DCL (
    GPA(12), INA(14), L, LMAIN, I, J, CAL
) FIXED BIN(31,0);
DCL (
    NXLFT, SUBSBT
) ENTRY(FIXED BIN(31,0), FIXED BIN(31,0))
RETURNS(FIXED BIN(31,0));
DCL (
    LIST, STRGN, NAMIST, BOT
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL CHT CHAR(10), CHF CHAR(4);
GPL=' '; GPA=0; INL=' '; INA=0; K=0;
CAL=LIST(LMAIN); J=0; CH=' ';
NUCH: IF J=1 THEN DO; J=0; GOTO SORT; END;
NXT: GET EDIT(C)(A(1)); IF C=' ' THEN GOTO NXT;
IF C='#' THEN RETURN;
IF CH>= 'A' & CH<= 'Z' THEN
    IF C='(' | (C>= 'A' & C<= 'Z') THEN

```

```

STRGN 65
STRGN 66
STRGN 67
STRNGET 31
STRNGET 32
STRNGET 33
STRNGET 0
STRNGET 1
STRNGET 2
STRNGET 3
STRNGET 4
STRNGET 5
STRNGET 6
STRNGET 7
STRNGET 8
STRNGET 9
STRNGET 10
STRNGET 11
STRNGET 12
STRNGET 13
STRNGET 14
STRNGET 15
STRNGET 16
STRNGET 17
STRNGET 18
STRNGET 19
STRNGET 20
STRNGET 21
STRNGET 22
STRNGET 23
STRNGET 24
STRNGET 25
STRNGET 26
STRNGET 27

```


DO; J=1; CH='.'; GOTO OPRTR; END;	STRNGET 28
IF CH=')' & ((C>='A' & C<='Z') C='(') THEN	STRNGET 29
DO; J=1; CH='.'; GOTO OPRTR; END;	STRNGET 30
SORT: CH=C;	STRNGET 31
IF C>'Z' THEN GOTO INT;	STRNGET 32
K=0;	STRNGET 33
IF C<'A' THEN GOTO OPRTR;	STRNGET 34
IF C<'M' THEN GOTO GPCL;	STRNGET 35
IF C<'O' THEN GOTO INTLIT;	STRNGET 36
SIGNAL CONDITION(ALPHA);	STRNGET 37
OPRTR: I=BOT(LMAIN);	STRNGET 38
IF NAMTST(I)=0 THEN GOTO LA;	STRNGET 39
UNSPEC(CHA)=SUBSTR(UNSPEC(1),25,8);	STRNGET 40
IF CH='.' & CHA='.' THEN DO;	STRNGET 41
UNSPEC(CHF)=UNSPEC(1);	STRNGET 42
CHT='00000000' SUBSTR(CHF,3,1); CHT=CHT+'1';	STRNGET 43
I=UNSPEC('00' SUBSTR(CHT,9,1) CHA);	STRNGET 44
CAL=SUBSBT(1,LMAIN); GOTO NUCH;	STRNGET 45
END;	STRNGET 46
LA: CAL=NXTLFT(UNSPEC('000' CH),LMAIN); GOTO NUCH;	STRNGET 47
GPCL: DO I=1 TO 12;	STRNGET 48
IF GPL(I)=CH THEN	STRNGET 49
DO; CAL=NXTLFT(GPA(I),LMAIN); GOTO NUCH; END;	STRNGET 50
IF GPL(I)=' ' THEN DO;	STRNGET 51
CAL=NXTLFT(UNSPEC('OGO' CH),LIST(L));	STRNGET 52
GPL(I)=CH; CAL=NXTLFT(L,LMAIN); GPA(I)=L; GOTO NUCH; END;	STRNGET 53
END;	STRNGET 54
SIGNAL CONDITION(ALPHA);	STRNGET 55
INTLIT: DO I=1 TO 14;	STRNGET 56
IF INL(I)=CH THEN	STRNGET 57
DO; CAL=NXTLFT(INA(I),LMAIN); GOTO NUCH; END;	STRNGET 58
IF INL(I)=' ' THEN DO;	STRNGET 59
CAL=NXTLFT(UNSPEC('OLO' CH),LIST(L));	STRNGET 60
INL(I)=CH; INA(I)=L;	STRNGET 61

	CAL=NXTLEFT(L,LMAIN); GOTO NUCH; END;	STRNGET 62
	END;	STRNGET 63
	SIGNAL CONDITION(ALPHA);	STRNGET 64
INT:	IF K=1 THEN GOTO INB;	STRNGET 65
	K=1; CAL=NXTLEFT(UNSPEC('010' CA),LMAIN); GOTO NUCH;	STRNGET 66
INB:	K=0; CAL=NXTLEFT(UNSPEC('01' C CA),LMAIN); GOTO NUCH;	STRNGET 67
	END STRNGET;	STRNGET 68
STRPN:	PROC(STR,N) RECURSIVE;	STRPN 0
	/* THIS PROC PRINTS A STRUCTURE, STR, SUPPLIED IN	STRPN 1
	NAME FORMAT. LEFT NORMING OF COMMUTATORS IS ASSUMED. */	STRPN 2
	DCL (STRPN 3
	A,B,CAL,I,NCOM,STR,N,PA,PB,PS	STRPN 4
) FIXED BIN(31,0);	STRPN 5
	DCL (STRPN 6
	SYMB	STRPN 7
)ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));	STRPN 8
	DCL (STRPN 9
	SYM	STRPN 10
) CHAR(1);	STRPN 11
	DCL (STRPN 12
	COMCNT	STRPN 13
)ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));	STRPN 14
	DCL (STRPN 15
	PREC	STRPN 16
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	STRPN 17
	DCL (STRPN 18
	NTERM	STRPN 19
)ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));	STRPN 20
	DCL (STRPN 21
	BOT,CONT,TOP	STRPN 22
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	STRPN 23

DCL (STRPN	24
MADNTP,STRPN	STRPN	25
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	STRPN	26
RETURNS(FIXED BIN(31,0));	STRPN	27
DCL IDNT ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));	STRPN	28
IF N>500 THEN	STRPN	29
DO; PUT SKIP(N-500); CAL=COMCNT(STR); N=1; END;	STRPN	30
IF N=500 THEN DO; CAL=COMCNT(STR); N=1; END;	STRPN	31
IF N>200 THEN DO; N=N-200; CAL=COMCNT(STR); END;	STRPN	32
IF N>87 THEN DO; PUT SKIP(2); N=1; END;	STRPN	33
IF N=0 THEN DO; PUT SKIP(3); CAL=COMCNT(STR); N=1; END;	STRPN	34
SYM=SYMB(STR);	STRPN	35
IF -NTERM(STR) THEN DO;	STRPN	36
IF IDNT(STR) THEN DO; PUT EDIT('1')(A); N=N+1; RETURN; END;	STRPN	37
IF SUBSTR(UNSPEC(TOP(STR)),9,8)=UNSPEC('1') THEN DO;	STRPN	38
UNSPEC(SYM)=SUBSTR(UNSPEC(TOP(STR)),17,8);	STRPN	39
IF SYM-='0' & SYM-=' ' THEN DO; PUT EDIT(SYM)(A); N=N+1; END;	STRPN	40
SYM=SYMB(STR); PUT EDIT(SYM)(A); N=N+1; RETURN; END;	STRPN	41
PUT EDIT(SYM)(A); N=N+1; RETURN;	STRPN	42
END;	STRPN	43
A=CONT(MADNTP(STR,2)+1); B=BOT(STR);	STRPN	44
PS=PREC(STR); IF SYM=':' THEN PS=PS+1;	STRPN	45
PA=PREC(A); PB=PREC(B);	STRPN	46
IF SYM=''' SYM='U' THEN GOTO L1; /*UNARY OPERATION */	STRPN	47
/* STR IS NON TERMINAL & BINARY */	STRPN	48
IF PS>PA THEN DO; PUT EDIT('(')(A); N=N+1; END;	STRPN	49
CAL=STRPN(A,N);	STRPN	50
IF PS>PA THEN DO; PUT EDIT(')')(A); N=N+1; END;	STRPN	51
IF SYM='.' THEN GOTO L2;	STRPN	52
IF SYM-=',' THEN DO;	STRPN	53
PUT EDIT(SYM)(A); N=N+1; GOTO L2; END;	STRPN	54
NCOM=(27)'0'B SUBSTR(UNSPEC(TOP(STR)),21,4);	STRPN	55
	STRPN	56
	STRPN	57

DO I=1 TO NCOM; PUT EDIT('')(A); N=N+1; END;	STRPN	58
L2: IF PS>PB THEN DO; PUT EDIT('')(A); N=N+1; END;	STRPN	59
CAL=STRPN(B,N);	STRPN	60
IF PS>PB THEN DO; PUT EDIT('')(A); N=N+1; END;	STRPN	61
RETURN;	STRPN	62
L1: IF SYM=''' THEN DO;	STRPN	63
IF PS>PA THEN DO; PUT EDIT('')(A); N=N+1; END;	STRPN	64
CAL=STRPN(A,N);	STRPN	65
IF PS>PA THEN DO; PUT EDIT('')(A); N=N+1; END;	STRPN	66
PUT EDIT('''')(A); N=N+1; RETURN; END;	STRPN	67
IF SYM='U' THEN DO;	STRPN	68
PUT EDIT('U')(A); N=N+1;	STRPN	69
IF PS>PA THEN DO; PUT EDIT('')(A); N=N+1; END;	STRPN	70
CAL=STRPN(A,N);	STRPN	71
IF PS>PA THEN DO; PUT EDIT('')(A); N=N+1; END;	STRPN	72
RETURN; END;	STRPN	73
END STRPN;	STRPN	74

SUM: PROC(A,B) FIXED BIN(31,0);	SUM	0
/* ADDS TWO INTEGERS IN LIST FORMAT AND RETURNS THE SUM	SUM	1
IN LIST FORMAT */	SUM	2
DCL (AM,BM) BIT(1);	SUM	3
DCL (DL,DR) FIXED DEC(5,0);	SUM	4
DCL (A,B,L,R,CAL) FIXED BIN(31,0);	SUM	5
DCL CT CHAR(2);	SUM	6
DCL (SUM	7
LIST,BOT	SUM	8
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SUM	9
DCL (SUM	10
SYMB	SUM	11
)ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));	SUM	12
DCL (SUM	13

NTERM	SUM	14
ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));	SUM	15
DCL (SUM	16
NXTLFT	SUM	17
ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))	SUM	18
RETURNS(FIXED BIN(31,0));	SUM	19
DCL (SUM	20
LG	SUM	21
ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))	SUM	22
RETURNS(FIXED BIN(31,0));	SUM	23
AM='0'B; BM='0'B;	SUM	24
IF NTERM(A) THEN IF SYMB(A)='U' THEN	SUM	25
DO; AM='1'B; L=BOT(BOT(A)); END;	SUM	26
ELSE SIGNAL CONDITION(ALPHA);	SUM	27
ELSE L=BOT(A);	SUM	28
IF NTERM(B) THEN IF SYMB(B)='U' THEN	SUM	29
DO; BM='1'B; R=BOT(BOT(B)); END;	SUM	30
ELSE SIGNAL CONDITION(ALPHA);	SUM	31
ELSE R=BOT(B);	SUM	32
UNSPEC(CT)=SUBSTR(UNSPEC(L),17,16);	SUM	33
DL=CT; IF AM THEN DL=-DL;	SUM	34
UNSPEC(CT)=SUBSTR(UNSPEC(R),17,16);	SUM	35
DR=CT; IF BM THEN DR=-DR;	SUM	36
DL=DL+DR;	SUM	37
IF DL>99 THEN SIGNAL CONDITION(ALPHA);	SUM	38
IF DL=0 THEN DO;	SUM	39
R=LIST(9); CAL=NXTLFT(UNSPEC('0100'),R);	SUM	40
RETURN(R); END;	SUM	41
IF DL<0 THEN DO;	SUM	42
DL=-DL;	SUM	43
R=LIST(9); CAL=NXTLFT(UNSPEC('01') SUBSTR(CHAR(DL),7,2),R);	SUM	44
R=LG('U',R,0);	SUM	45
RETURN(R); END;	SUM	46
IF DL>0 THEN DO;	SUM	47

```

R=LIST(9); CAL=NXTLEFT(UNSPEC('01' || SUBSTR(CHAR(DL),7,2)),R);
RETURN(R);
SIGNAL CONDITION (ALPHA);
END SUM;

```

```

SUMEX 48
SUMEX 49
SUMEX 50
SUMEX 51

```

```

SUMEX: PROC(STR) FIXED BIN(31,0);
/* (A:M)(A:N) -> A:(M+N) */
/* UNARY MINUS IS REPLACED BY BINARYMINUS OR RAISED
   TO THE TOP OF THE RESULTING INTEGRAL STRUCTURE */
DCL (
    STR,A,B,M,N,L,R,CAL
) FIXED BIN(31,0);
DCL (
    BOT,TOP,CONT,MLIST,LIST
) ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL (
    LG
) ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL (
    MADNTP,INLSTR,NXTLEFT
) ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
    RETURNS(FIXED BIN(31,0));
DCL (
    LU,RU
) BIT(1);
DCL IRALST ENTRY(FIXED BIN(31,0));
DCL(
    SYMB
) ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL NTERM ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));
DCL (GL,GR) BIT(1);

```

```

SUMEX 0
SUMEX 1
SUMEX 2
SUMEX 3
SUMEX 4
SUMEX 5
SUMEX 6
SUMEX 7
SUMEX 8
SUMEX 9
SUMEX 10
SUMEX 11
SUMEX 12
SUMEX 13
SUMEX 14
SUMEX 15
SUMEX 16
SUMEX 17
SUMEX 18
SUMEX 19
SUMEX 20
SUMEX 21
SUMEX 22
SUMEX 23
SUMEX 24
SUMEX 25
SUMEX 26

```

GL,GR='0'B;	SUMEX	27
L=CONT(MADNTP(STR,2)+1);	SUMEX	28
IF SYMB(L)=':' THEN A=CONT(MADNTP(L,2)+1);	SUMEX	29
ELSE A=L;	SUMEX	30
IF SYMB(A)=':' THEN DO;A=BOT(A); GL='1'B; END;	SUMEX	31
IF SYMB(L)~=':'	SUMEX	32
THEN DO; M=LIST(9); CAL=NXTLFT(UNSPEC('0101'),M); END;	SUMEX	33
ELSE M=BOT(L);	SUMEX	34
R=BOT(STR);	SUMEX	35
IF SYMB(R)~=':'	SUMEX	36
THEN DO; N=LIST(9); CAL=NXTLFT(UNSPEC('0101'),N);	SUMEX	37
IF SYMB(R)=':' THEN GR='1'B;	SUMEX	38
END;	SUMEX	39
ELSE DO; N=BOT(R);	SUMEX	40
IF SYMB(CONT(MADNTP(R,2)+1))=':' THEN GR='1'B;	SUMEX	41
END;	SUMEX	42
LU='0'B; RU='0'B;	SUMEX	43
IF SUBSTR(UNSPEC(TOP(M)),9,24)=UNSPEC('000') THEN LU='1'B;	SUMEX	44
IF SUBSTR(UNSPEC(TOP(N)),9,24)=UNSPEC('000') THEN RU='1'B;	SUMEX	45
IF GL & ~LU THEN DO; M=LG('U',M,0); LU='1'B; GOTO K1; END;	SUMEX	46
IF GL & LU THEN DO; M=BOT(M); LU='0'B; END;	SUMEX	47
K1: IF GR & ~RU THEN DO; N=LG('U',N,0); RU='1'B; GOTO K2; END;	SUMEX	48
IF GR & RU THEN DO; N=BOT(N); RU='0'B; END;	SUMEX	49
K2:	SUMEX	50
IF LU & RU THEN B=LG('U',LG('+',BOT(M),BOT(N)),0);	SUMEX	51
IF LU & ~RU THEN B=LG('-',N,BOT(M));	SUMEX	52
IF ~LU & RU THEN B=LG('-',M,BOT(N));	SUMEX	53
IF ~LU & ~RU THEN B=LG('+',M,N);	SUMEX	54
CAL=IRALST(INLSTR(LG(':',A,B),MTLIST(STR)));	SUMEX	55
END SUMEX;	SUMEX	56

```

/* STR IS A STRUCTURE WITH OPERATOR '+' OR '-'. IT IS ASSUMED THAT
THE LEFT AND RIGHT PARTS ARE IN CANONIC FORM AND IT
CONVERTS THE STRUCTURE TO CANONIC FORM. */

```

```

DCL(
  L,R,BR,CAL,STR
  IF WL)FIXED BIN(31,0);
DCL(
  IF TL,TR
  )CHAR(4);
DCL(
  SUBST,MADNTP,INLSTR,SUM,SUBSTP
  )ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
DCL(
  LG
  )ENTRY(CHAR(1),FIXED BIN(31,0),FIXED BIN(31,0))
  RETURNS(FIXED BIN(31,0));
DCL(
  SYMB
  )ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));
DCL(
  CONT,MTLIST,TOP,BOT
  )ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));
DCL(
  NTERM
  )ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));
DCL(WL,WR) CHAR(1);
DCL(LT,RT) CHAR(3);
DCL(WRA,WRB) CHAR(1);
L=CONT(MADNTP(STR,2)+1); R=BOT(STR);
IF SYMB(STR)='-' THEN DO;
  IF SYMB(R)='U' THEN DO;
    R=BOT(R); CAL=SUBSTP(UNSPEC('000+'),STR);
  END;

```

```

SUMINT 1
SUMINT 2
SUMINT 3
SUMINT 4
SUMINT 5
SUMINT 6
SUMINT 7
SUMINT 8
SUMINT 9
SUMINT 10
SUMINT 11
SUMINT 12
SUMINT 13
SUMINT 14
SUMINT 15
SUMINT 16
SUMINT 17
SUMINT 18
SUMINT 19
SUMINT 20
SUMINT 21
SUMINT 22
SUMINT 23
SUMINT 24
SUMINT 25
SUMINT 26
SUMINT 27
SUMINT 28
SUMINT 29
SUMINT 30
SUMINT 31
SUMINT 32
SUMINT 33
SUMINT 34

```



```

      ELSE R=LG('U',R,0);
      END;
UNSPEC(TL)=UNSPEC(TOP(L)); UNSPEC(TR)=UNSPEC(TOP(R));
WL=SUBSTR(TL,2,1); WR=SUBSTR(TR,2,1);
LT=SUBSTR(TL,2,3); RT=SUBSTR(TR,2,3);
IF WL='I' & WR='I' THEN
  DO; L=SUM(L,R); GOTO RET; END;
IF LT='OCU' & WR='I' THEN DO;
  UNSPEC(WRA)=SUBSTR(UNSPEC(TOP(BOT(L))),9,8);
  IF WRA='I' THEN DO; L=SUM(L,R); GOTO RET; END;
  END;
IF WL='I' & RT='COU' THEN DO;
  UNSPEC(WRB)=SUBSTR(UNSPEC(TOP(BOT(R))),9,8);
  IF WRB='I' THEN DO; L=SUM(L,R); GOTO RET; END;
  END;
IF LT='DOU' & RT='COU' THEN DO;
  UNSPEC(WRA)=SUBSTR(UNSPEC(TOP(BOT(L))),9,8);
  UNSPEC(WRB)=SUBSTR(UNSPEC(TOP(BOT(R))),9,8);
  IF WRA='I' & WRB='I' THEN DO;
    L=SUM(L,R); GOTO RET; END;
  END;
IF -INTERM(L) & -INTERM(R) THEN RETURN;
IF LT='I00' | LT='I 0' THEN DO; L=R; GOTO RET; END;
IF RT='I00' | RT='I 0' THEN GOTO RET;
IF SUBSTR(TR,2,1)='I' THEN DO; L=LG('+',R,L); GOTO RET; END;
IF SUBSTR(TR,2,3)='DOU' THEN DO;
  BR=BOT(R); L=LG('-',L,BR);
  GOTO RET;
  END;
RETURN;
RET:  CAL=IRALST(INLSTR(L,MTLIST(STR)));
      END SUMINT;

```

```

SUMINT 35
SUMINT 36
SUMINT 37
SUMINT 38
SUMINT 39
SUMINT 40
SUMINT 41
SUMINT 42
SUMINT 43
SUMINT 44
SUMINT 45
SUMINT 46
SUMINT 47
SUMINT 48
SUMINT 49
SUMINT 50
SUMINT 51
SUMINT 52
SUMINT 53
SUMINT 54
SUMINT 55
SUMINT 56
SUMINT 57
SUMINT 58
SUMINT 59
SUMINT 60
SUMINT 61
SUMINT 62
SUMINT 63
SUMINT 64
SUMINT 65

```

SECTION 13 - PROCEDURES OF THE COLLECTING PROCESS

71

DCL (SYMB	1
TOP	SYMB	2
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));	SYMB	3
DCL (SYMB	4
STR	SYMB	5
) FIXED BIN(31,0);	SYMB	6
DCL CH CHAR(1);	SYMB	7
UNSPEC(CH)=SUBSTR(UNSPEC(TOP(STR)),25,6);	SYMB	8
RETURN(CH);	SYMB	9
END SYMB;	SYMB	10

WT:	PROC(CMXP)	FIXED BIN(31,0)	RECURSIVE;	WT	0
	DCL (WT	1
	CMXP,WT,WTB,WTM			WT	2
) FIXED BIN(31,0);			WT	3
	DCL (WT	4
	IDNT,GRPL			WT	5
)ENTRY(FIXED BIN(31,0)) RETURNS(BIT(1));			WT	6
	DCL (WT	7
	SYMB			WT	8
)ENTRY(FIXED BIN(31,0)) RETURNS(CHAR(1));			WT	9
	DCL (WT	10
	TOP,BOT,LOCT,CONT,WT			WT	11
)ENTRY(FIXED BIN(31,0)) RETURNS(FIXED BIN(31,0));			WT	12
	DCL (WT	13
	MADNTP			WT	14
)ENTRY(FIXED BIN(31,0),FIXED BIN(31,0))			WT	15
	RETURNS(FIXED BIN(31,0));			WT	16
	DCL (WT	17
	SYMBOL			WT	18
) CHAR(1);			WT	19
	CMXP=LOCT(CMXP);			WT	20

SECTION 13 - PROCEDURES OF THE COLLECTING PROCESS

72

SECTION 14 - SAMPLE CALCULATION

IF GRPL(CMXP) THEN RETURN(1);	/* GROUP GENERATOR */	WT	21
IF IDNT(CMXP) THEN RETURN('0'B (31)'1'B);	/* MAX POS NUMBER */	WT	22
SYMBOL=SYMB(CMXP); WTM=WT(CONT(MADNTP(CMXP,2)+1));		WT	23
IF SYMBOL=''' THEN RETURN(WTM);	/* INVERSE */	WT	24
IF SYMBOL=':' THEN RETURN(WTM);	/* EXP */	WT	25
WTB=WT(BOT(CMXP));		WT	26
IF SYMBOL='.' THEN		WT	27
IF WTM < WTB THEN RETURN(WTM); ELSE RETURN(WTB);	/* PROD */	WT	28
IF SYMBOL=', ' THEN RETURN(WTM+WTB);	/* COMMUTATOR */	WT	29
SIGNAL CONDITION(ALPHA);		WT	30
END WT;		WT	31

SECTION 14 - SAMPLE CALCULATION

1

(AB):4

ENTER ORDPROD

AB

AB

EXIT ORDPROD

(AB):4

ABABABAB

ENTER ORDPROD

ABABABAB

AAB(B,A)BABAB

AAAB(B,A)(B,A)(B,A,A)B(B,A)BAB

AAAAB(B,A)(B,A)(B,A,A)(B,A)(B,A,A)(B,A,A)(B,A,A,A)B(B,A)(B,A)(B,A,A)B(B,A)B

A:4B(B,A):2(B,A,A)(B,A)(B,A,A):2(B,A,A,A)B(B,A):2(B,A,A)B(B,A)B

ENTER LEO

(B,A):2,B

ENTER ORDPROD

(B,A,B):2(B,A,B,,B,A)

(B,A,B):2(B,A,B,,B,A)

SECTION 14 - SAMPLE CALCULATION

2

EXIT ORDPROD

(B, A, B):2(B, A, B,, B, A)

EXIT LEO

ENTER LEO

(B, A, A):2, B

(B, A, A, B):2

EXIT LEO

A:4BB(B, A):2(B, A, B):2(B, A, B,, B, A)(B, A, A)(B, A, A, B)(B, A)(B, A, B)(B, A, A):2(B, A, A, B):2(B, A, A, A, A)(B, A, A, A, B)(B, A):2(B, A, A)B(B, A)B

ENTER LEO

(B, A):2, B

ENTER ORDPROD

(E, A, B):2(B, A, B,, B, A)

(B, A, B):2(B, A, B,, B, A)

EXIT ORDPROD

(B, A, B):2(B, A, B,, B, A)

EXIT LEO

ENTER LEO

SECTION 14 - SAMPLE CALCULATION

3

(B,A,B):2,B

(B,A,B,B):2

EXIT LEO

ENTER LEO

(B,A,A):2,B

(B,A,A,B):2

EXIT LEO

ENTER LEO

(B,A,A,B):2,B

(B,A,A,B,B):2

EXIT LEO

ENTER LEO

(B,A):2,B

ENTER ORCPROD

(B,A,B):2(B,A,B,,B,A)

(B,A,B):2(B,A,B,,B,A)

EXIT ORCPROD

(B,A,B):2(B,A,B,,B,A)

EXIT LEO

A:4BBB(B,A):2(B,A,B):2(B,A,B,,B,A)(B,A,B):2(B,A,B,B):2(B,A,B,,B,A)(B,A,A)(B,A,A,B)(B,A,A,B)(B,A,A,B,B)(B,A)(B,A,B)(B,A,B)(B,A,B,B)(B,A,A):2(B,A,A,B):2(B,A,A,B):2(B,A,A,B,B):2(B,A,A,A)(B,A,A,A,B)(B,A,A,A,B)(B,A):2(B,A,B):2(B,A,B,,B,A)(B,A,A)(B,A,A,B)(B,A)B

ENTER LEO

(B,A):2,B

ENTER ORDPROD

(B,A,B):2(B,A,B,,B,A)

(B,A,B):2(B,A,B,,B,A)

EXIT ORDPROD

(B,A,B):2(B,A,B,,B,A)

EXIT LEO

(B,A,B):2,B

(B,A,B,B):2

EXIT LEO

ENTER LEO

(B,A,B):2,B

(B,A,B,B):2

SECTION 14 - SAMPLE CALCULATION

5

EXIT LEO

ENTER LEO

(B,A,B,B):2,B

(B,A,B,B,B):2

EXIT LEO

ENTER LEO

(B,A,A):2,B

(B,A,A,B):2

EXIT LEO

ENTER LEO

(B,A,A,B):2,B

(B,A,A,B,B):2

EXIT LEO

ENTER LEO

(B,A,A,B):2,B

(B,A,A,B,B):2

EXIT LEO

ENTER LEO

SECTION 14 - SAMPLE CALCULATION

6

(B,A):2,B

ENTER ORCPROD

(B,A,B):2(B,A,B,,B,A)

(B,A,B):2(B,A,B,,B,A)

EXIT ORCPROD

(B,A,B):2(B,A,B,,B,A)

EXIT LEO

ENTER LEO

(B,A,B):2,B

(B,A,B,B):2

EXIT LEO

A:4B:4(B,A):2(B,A,B):2(B,A,B,,B,A)(B,A,B):2(B,A,B,B):2(B,A,B,,B,A)(B,A,B):2(B,A,B,B):4(B,A,B,B,B):2(B,A,B,,B,A)(B,A,A)(B,A,A,B):2(B,A,A,B,B)(B,A,A,B)(B,A,A,B,B):2(B,A)(B,A,B):2(B,A,B,B)(B,A,B)(B,A,B,B):2(B,A,B,B,B)(B,A,A):2(B,A,A,B):4(B,A,A,B,B):2(B,A,A,B):2(B,A,A,B,B):4(B,A,A,A)(B,A,A,A,B):3(B,A):2(B,A,B):2(B,A,B,,B,A)(B,A,B):2(B,A,B,B):2(B,A,B,,B,A)(B,A,A)(B,A,A,B):2(B,A,A,B,B)(B,A)(B,A,B)

ENTER LEO

(B,A,B):2,,B,A

(B,A,B,,B,A):2

SECTION 14 - SAMPLE CALCULATION

7

EXIT LEO

ENTER LEO

(B,A,B):2,,B,A

(B,A,B,,B,A):2

EXIT LEO

ENTER LEO

(B,A,B):2,,B,A

(B,A,B,,B,A):2

EXIT LEO

A:4B:4(B,A):2(B,A)(B,A,B):2(B,A,B,,B,A):2(B,A,B,,B,A)(B,A,B):2(B,A,B,,B,A):2(B,A,B,B):2(B,A,B,,B,A)(B,A,B):2(B,A,B,,B,A):2(B,A,B,B):4(B,A,B,B,B):2(B,A,B,,B,A)(B,A,A)(B,A,A,,B,A)(B,A,A,B):2(B,A,A,B,B)(B,A,A,B)(B,A,A,B,B):2(B,A,B):2(B,A,B,B)(B,A,B)(B,A,B,B):2(B,A,B,B,B)(B,A,A):2(B,A,A,B):4(B,A,A,B,B):2(B,A,A,B):2(B,A,A,B,B):4(B,A,A,A)(B,A,A,A,B):3(B,A):2(B,A,B):2(B,A,B,,B,A)(B,A,B):2(B,A,B,B):2(B,A,B,,B,A)(B,A,A)(B,A,A,B):2(B,A,A,B,B)(B,A)(B,A,B)

ENTER REO

(B,A,B):2,(B,A):2

ENTER REO

B,A,(B,A,B):2

SECTION 14 - SAMPLE CALCULATION

8

(B,A,B,,B,A):-2

EXIT REO

(B,A,B,,B,A):4

EXIT REO

ENTER REO

(B,A,B):2,(B,A):2

ENTER REO

B,A,(B,A,B):2

(B,A,B,,B,A):-2

EXIT REO

(B,A,B,,B,A):4

EXIT REO

ENTER REO

(B,A,B):2,(B,A):2

ENTER REO

B,A,(B,A,B):2

(B,A,B,,B,A):-2

EXIT REO

SECTION 14 - SAMPLE CALCULATION

9

(B,A,B,,B,A):4

EXIT REO

ENTER REO

B,A,A,(B,A):2

(B,A,A,,B,A):2

EXIT REO

ENTER REO

(B,A,B):2,(B,A):2

ENTER REO

B,A,(B,A,B):2

(B,A,B,,B,A):-2

EXIT REO

(B,A,B,,B,A):4

EXIT REO

ENTER REO

B,A,B,(B,A):2

(B,A,B,,B,A):2

EXIT REO

SECTION 14 - SAMPLE CALCULATION

10

ENTER RED

(B,A,A):2,(B,A):2

ENTER RED

B,A,(B,A,A):2

(B,A,A,,B,A):-2

EXIT RED

(B,A,A,,B,A):4

EXIT RED

A:4B:4(B,A):2(B,A)(B,A):2(B,A,B):2(B,A,B,,B,A):4(B,A,B,,B,A):2(B,A,B,,B,A)(B,A,B):2(B,A,
B,,B,A):4(B,A,B,,B,A):2(B,A,B,B):2(B,A,B,,B,A)(B,A,B):2(B,A,B,,B,A):4(B,A,B,,B,A):2(B,A,
B,B):4(B,A,B,B,B):2(B,A,B,,B,A)(B,A,A)(B,A,A,,B,A):2(B,A,A,,B,A)(B,A,A,B):2(B,A,A,B,B)(
B,A,A,B)(B,A,A,B,B):2(B,A,B):2(B,A,B,,B,A):4(B,A,B,B)(B,A,B)(B,A,B,,B,A):2(B,A,B,B):2(B,
A,B,B,B)(B,A,A):2(B,A,A,,B,A):4(B,A,A,B):4(B,A,A,B,B):2(B,A,A,B):2(B,A,A,B,B):4(B,A,A,A)
(B,A,A,A,B):3(B,A,B):2(B,A,B,,B,A)(B,A,B):2(B,A,B,B):2(B,A,B,,B,A)(B,A,A)(B,A,A,B):2(B,
A,A,B,B)(B,A)(B,A,B)

ENTER LEU

(B,A,B):2,,B,A

(B,A,B,,B,A):2

EXIT LEU

SECTION 14 - SAMPLE CALCULATION

11

ENTER LEO

(B,A,B):2,,B,A

(B,A,B,,B,A):2

EXIT LEO

ENTER LEO

(B,A,B):2,,B,A

(B,A,B,,B,A):2

EXIT LEO

ENTER LEO

(B,A,B):2,,B,A

(B,A,B,,B,A):2

EXIT LEO

ENTER LEO

(B,A,A):2,,B,A

(B,A,A,,B,A):2

EXIT LEO

ENTER LEO

(B,A,B):2,,B,A

SECTION 14 - SAMPLE CALCULATION

12

(B,A,B,,B,A):2

EXIT LEO

ENTER LEO

(B,A,B):2,,B,A

(B,A,B,,B,A):2

EXIT LEO

A:4B:4(B,A):2(B,A)(B,A):2(B,A)(B,A,B):2(B,A,B,,B,A):2(B,A,B,,B,A):4(B,A,B,,B,A):2(B,A,B,,
B,A)(B,A,B):2(B,A,B,,B,A):2(B,A,B,,B,A):4(B,A,B,,B,A):2(B,A,B,B):2(B,A,B,,B,A)(B,A,B):2
(B,A,B,,B,A):2(B,A,B,,B,A):4(B,A,B,,B,A):2(B,A,B,B):4(B,A,B,B,B):2(B,A,B,,B,A)(B,A,A)(B,
A,A,,B,A)(B,A,A,,B,A):2(B,A,A,,B,A)(B,A,A,B):2(B,A,A,B,B)(B,A,A,B)(B,A,A,B,B):2(B,A,B):
2(B,A,B,,B,A):2(B,A,B,,B,A):4(B,A,B,B)(B,A,B)(B,A,B,,B,A)(B,A,B,,B,A):2(B,A,B,B):2(B,A,
B,B,B)(B,A,A):2(B,A,A,,B,A):2(B,A,A,,B,A):4(B,A,A,B):4(B,A,A,B,B):2(B,A,A,B):2(B,A,A,B,
B):4(B,A,A,A)(B,A,A,A,B):3(B,A,B):2(B,A,B,,B,A):2(B,A,B,,B,A)(B,A,B):2(B,A,B,,B,A):2(B,
A,B,B):2(B,A,B,,B,A)(B,A,A)(B,A,A,,B,A)(B,A,A,B):2(B,A,A,B,B)(B,A,B)

A:4B:4(B,A):6(B,A,B):2(B,A,B,,B,A):9(B,A,B):2(B,A,B,,B,A):8(B,A,B,B):2(B,A,B,,B,A)(B,A,
B):2(B,A,B,,B,A):8(B,A,B,B):4(B,A,B,B,B):2(B,A,B,,B,A)(B,A,A)(B,A,A,,B,A):4(B,A,A,B):2(
B,A,A,B,B)(B,A,A,B)(B,A,A,B,B):2(B,A,B):2(B,A,B,,B,A):6(B,A,B,B)(B,A,B)(B,A,B,,B,A):3(B,
A,B,B):2(B,A,B,B,B)(B,A,A):2(B,A,A,,B,A):6(B,A,A,B):4(B,A,A,B,B):2(B,A,A,B):2(B,A,A,B,B):
4(B,A,A,A)(B,A,A,A,B):3(B,A,B):2(B,A,B,,B,A):3(B,A,B):2(B,A,B,,B,A):2(B,A,B,B):2(B,A,B,,
B,A)(B,A,A)(B,A,A,,B,A)(B,A,A,B):2(B,A,A,B,B)(B,A,B)

SECTION 14 - SAMPLE CALCULATION

13

A:4B:4(B,A):6(B,A,A):4(B,A,B):2(B,A,B,,B,A):9(B,A,B):2(B,A,B,,B,A):8(B,A,B,B):2(B,A,B,,
B,A)(B,A,B):2(B,A,B,,B,A):8(B,A,B,B):4(B,A,B,B,B):2(B,A,B,,B,A)(B,A,A,,B,A):4(B,A,A,B):
2(B,A,A,B,B)(B,A,A,B)(B,A,A,B,B):2(B,A,B):2(B,A,B,,B,A):6(B,A,B,B)(B,A,B)(B,A,B,,B,A):3
(B,A,B,B):2(B,A,B,B,B)(B,A,A,,B,A):6(B,A,A,B):4(B,A,A,B,B):2(B,A,A,B):2(B,A,A,B,B):4(B,
A,A,A)(B,A,A,A,B):3(B,A,B):2(B,A,B,,B,A):3(B,A,B):2(B,A,B,,B,A):2(B,A,B,B):2(B,A,B,,B,A)(
B,A,A,,B,A)(B,A,A,B):2(B,A,A,B,B)(B,A,B)

A:4B:4(B,A):6(B,A,A):4(B,A,B):14(B,A,B,,B,A):17(B,A,B,B):2(B,A,B,,B,A):9(B,A,B,B):4(B,A,
B,B,B):2(B,A,B,,B,A)(B,A,A,,B,A):4(B,A,A,B):2(B,A,A,B,B)(B,A,A,B)(B,A,A,B,B):2(B,A,B,,B,
A):6(B,A,B,B)(B,A,B,,B,A):3(B,A,B,B):2(B,A,B,B,B)(B,A,A,,B,A):6(B,A,A,B):4(B,A,A,B,B):2
(B,A,A,B):2(B,A,A,B,B):4(B,A,A,A)(B,A,A,A,B):3(B,A,B,,B,A):5(B,A,B,B):2(B,A,B,,B,A)(B,A,
A,,B,A)(B,A,A,B):2(B,A,A,B,B)

A:4B:4(B,A):6(B,A,A):4(B,A,B):14(B,A,A,A)(B,A,A,B):11(B,A,B,B):11(B,A,B,,B,A):26(B,A,B,
B,B):2(B,A,B,,B,A)(B,A,A,,B,A):4(B,A,A,B,B):3(B,A,B,,B,A):9(B,A,B,B,B)(B,A,A,,B,A):6(B,
A,A,B,B):6(B,A,A,A,B):3(B,A,B,,B,A):6(B,A,A,,B,A)(B,A,A,B,B)

A:4B:4(B,A):6(B,A,A):4(B,A,B):14(B,A,A,A)(B,A,A,B):11(B,A,B,B):11(B,A,A,,B,A):11(B,A,B,,
B,A):26(B,A,B,B,B):2(B,A,B,,B,A)(B,A,A,B,B):3(B,A,B,,B,A):9(B,A,B,B,B)(B,A,A,B,B):6(B,A,
A,A,B):3(B,A,B,,B,A):6(B,A,A,B,B)

A:4B:4(B,A):6(B,A,A):4(B,A,B):14(B,A,A,A)(B,A,A,B):11(B,A,B,B):11(B,A,A,,B,A):11(B,A,B,,
B,A):42(B,A,B,B,B):2(B,A,A,B,B):3(B,A,B,B,B)(B,A,A,B,B):6(B,A,A,A,B):3(B,A,A,B,B)

SECTION 14 - SAMPLE CALCULATION

14

A:4B:4(B,A):6(B,A,A):4(B,A,B):14(B,A,A,A)(B,A,A,B):11(B,A,B,B):11(B,A,A,,B,A):11(B,A,B,,
B,A):42(B,A,A,A,B):3(B,A,B,B,B):2(B,A,A,B,B):3(B,A,B,B,B)(B,A,A,B,B):7

A:4B:4(B,A):6(B,A,A):4(B,A,B):14(B,A,A,A)(B,A,A,B):11(B,A,B,B):11(B,A,A,,B,A):11(B,A,B,,
B,A):42(B,A,A,A,B):3(B,A,A,B,B):10(B,A,B,B,B):3

A:4B:4(B,A):6(B,A,A):4(B,A,B):14(B,A,A,A)(B,A,A,B):11(B,A,B,B):11(B,A,A,,B,A):11(B,A,B,,
B,A):42(B,A,A,A,B):3(B,A,A,B,B):10(B,A,B,B,B):3

EXIT URDPROD

A:4B:4(B,A):6(B,A,A):4(B,A,B):14(B,A,A,A)(B,A,A,B):11(B,A,B,B):11(B,A,A,,B,A):11(B,A,B,,
B,A):42(B,A,A,A,B):3(B,A,A,B,B):10(B,A,B,B,B):3

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